

EECS 219C: Computer-Aided Verification

Syntax-Guided Synthesis

(selected/adapted slides from
FMCAD'13 tutorial by R. Alur)

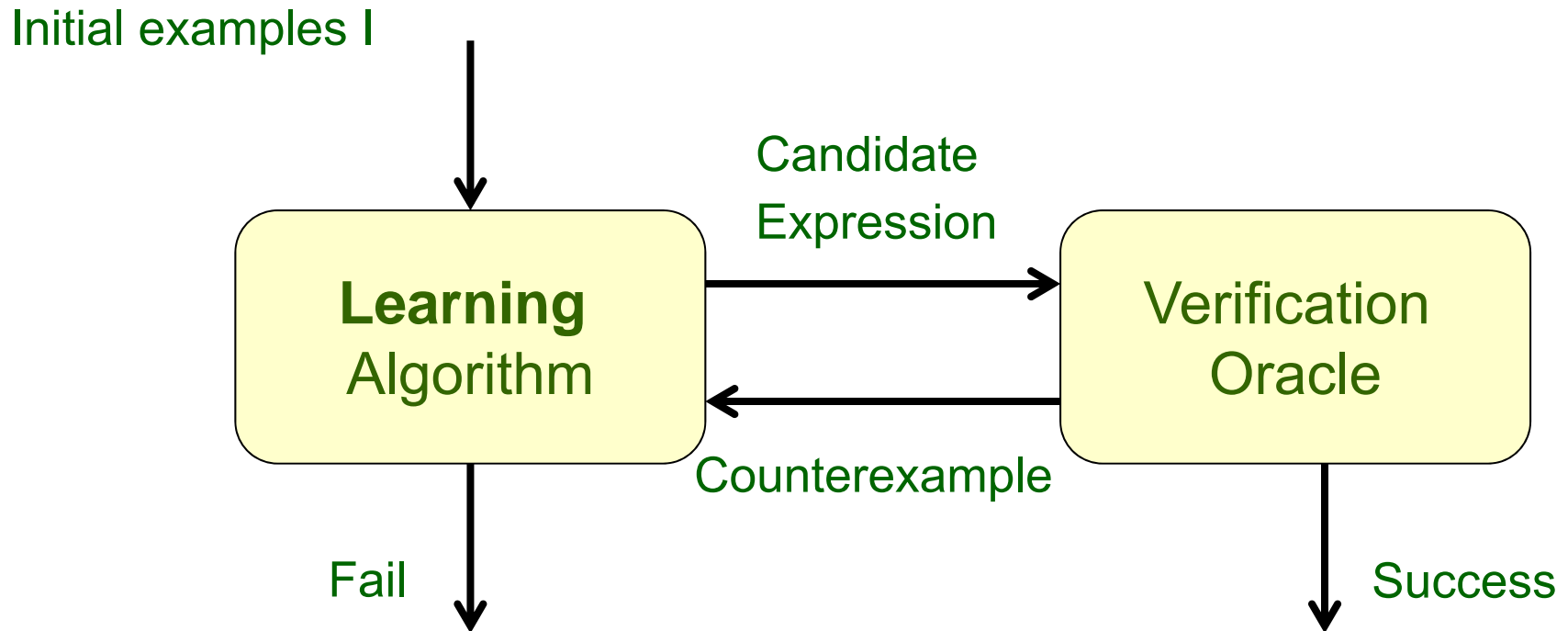
Sanjit A. Seshia

EECS, UC Berkeley

Solving SyGuS

- ❑ Is SyGuS same as solving SMT formulas with quantifier alternation?
- ❑ SyGuS can sometimes be reduced to Quantified-SMT, but not always
 - ◆ Set E is all linear expressions over input vars x, y
SyGuS reduces to $\exists a, b, c. \forall X. \varphi [f / ax+by+c]$
 - ◆ Set E is all conditional expressions
SyGuS cannot be reduced to deciding a formula in LIA
- ❑ Syntactic structure of the set E of candidate implementations can be used effectively by a solver
- ❑ Existing work on solving Quantified-SMT formulas suggests solution strategies for SyGuS

SyGuS as Active Learning

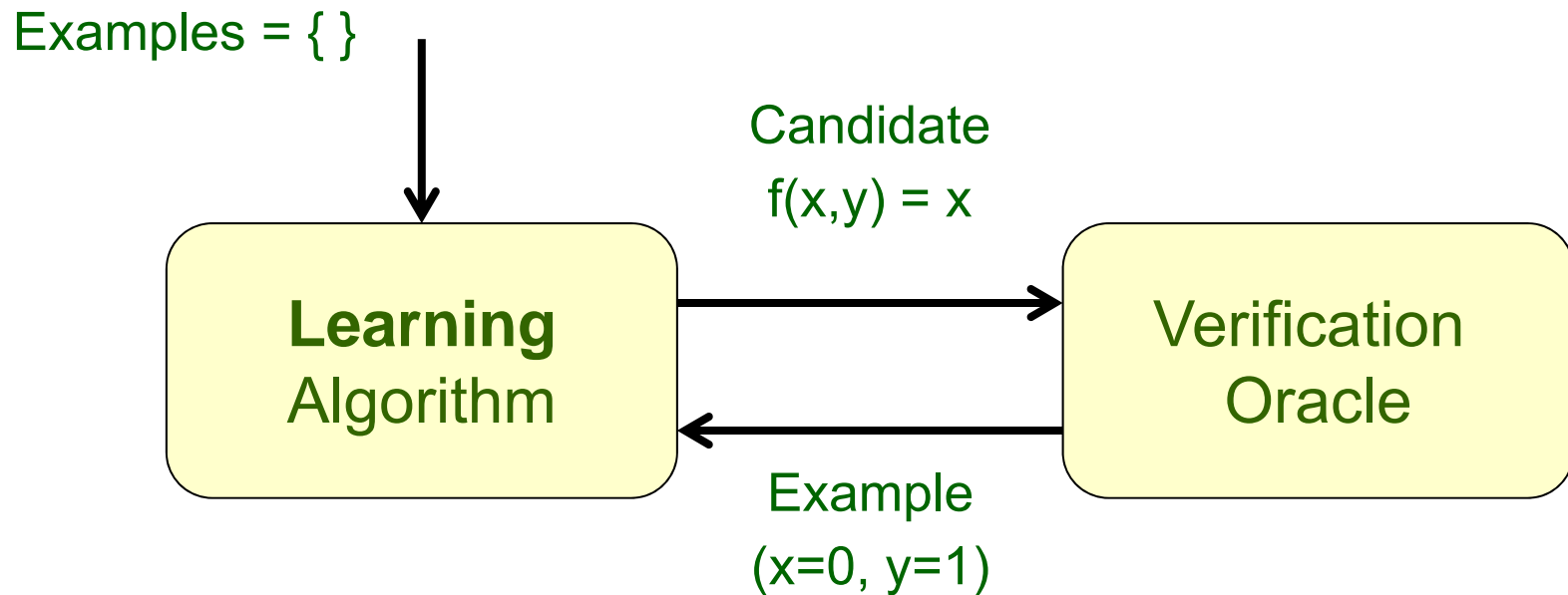


Concept class: Set E of expressions

Examples: Concrete input values

CEGIS Example

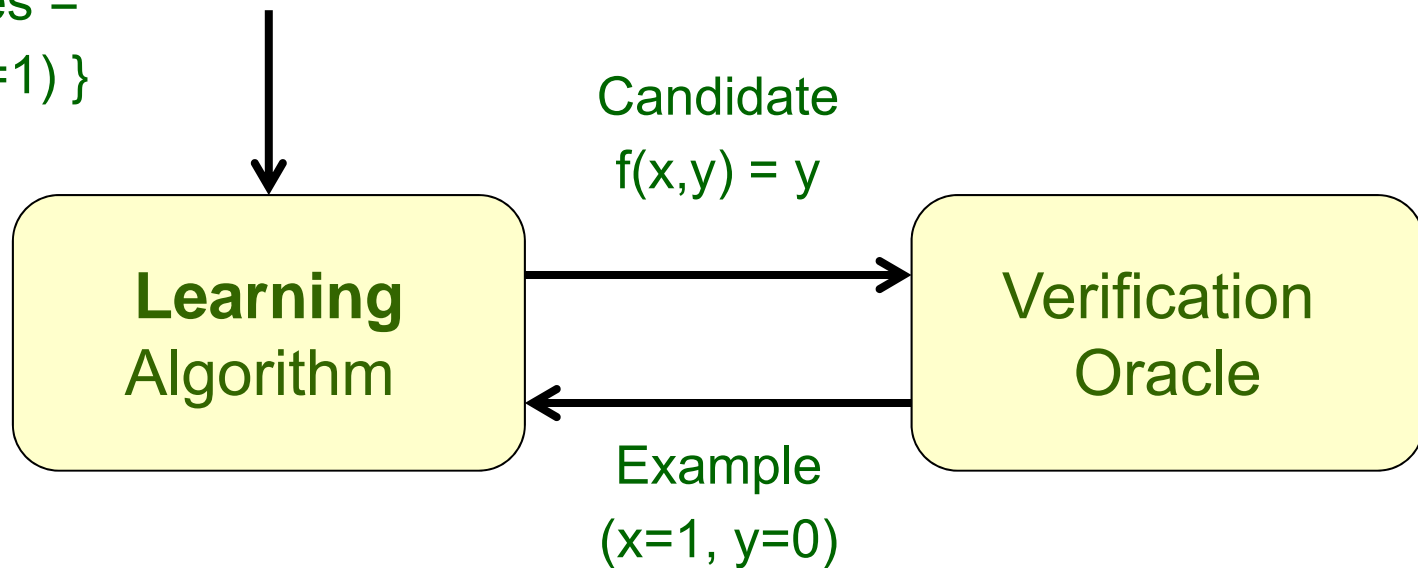
- ❑ Specification: $(x \leq f(x,y)) \ \& \ (y \leq f(x,y)) \ \& \ (f(x,y) = x \mid f(x,y) = y)$
- ❑ Set E: All expressions built from $x,y,0,1$, Comparison, $+$, If-Then-Else



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Examples =
 $\{(x=0, y=1)\}$



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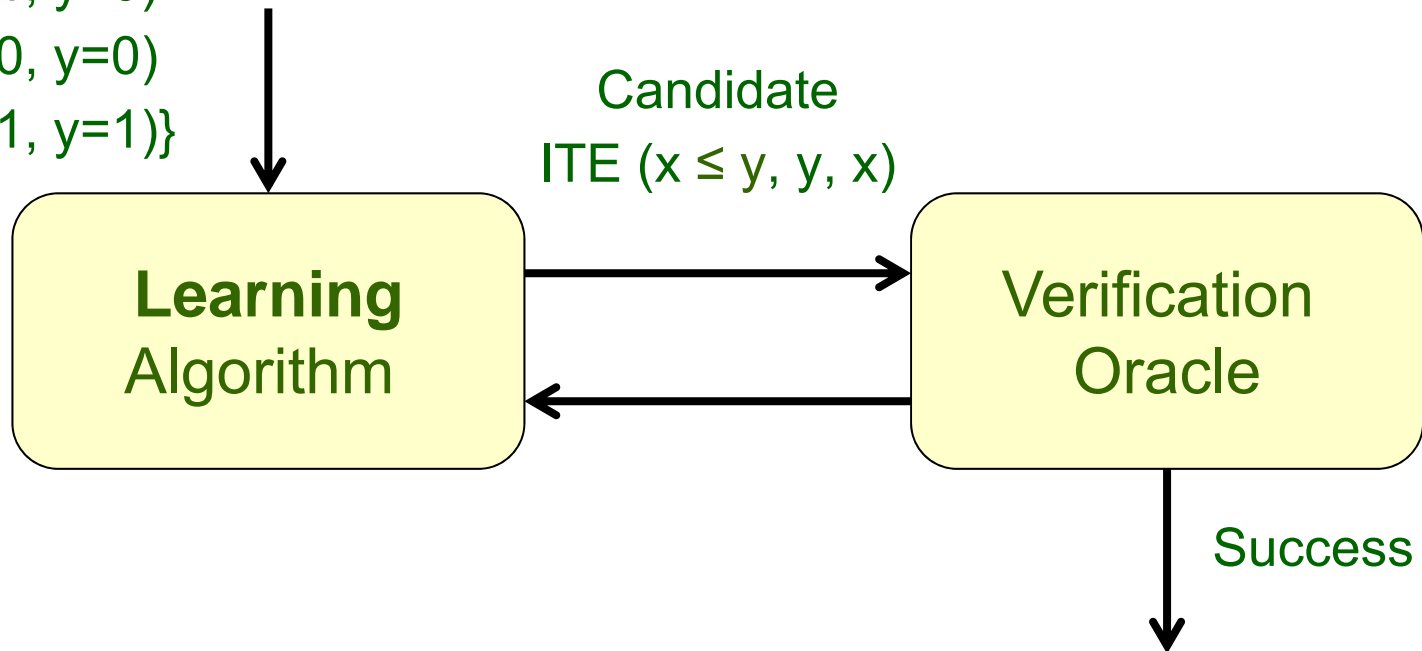
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$\{(x=0, y=1)$

$(x=1, y=0)$

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SyGuS Solutions

- CEGIS approach (Solar-Lezama, Seshia et al)

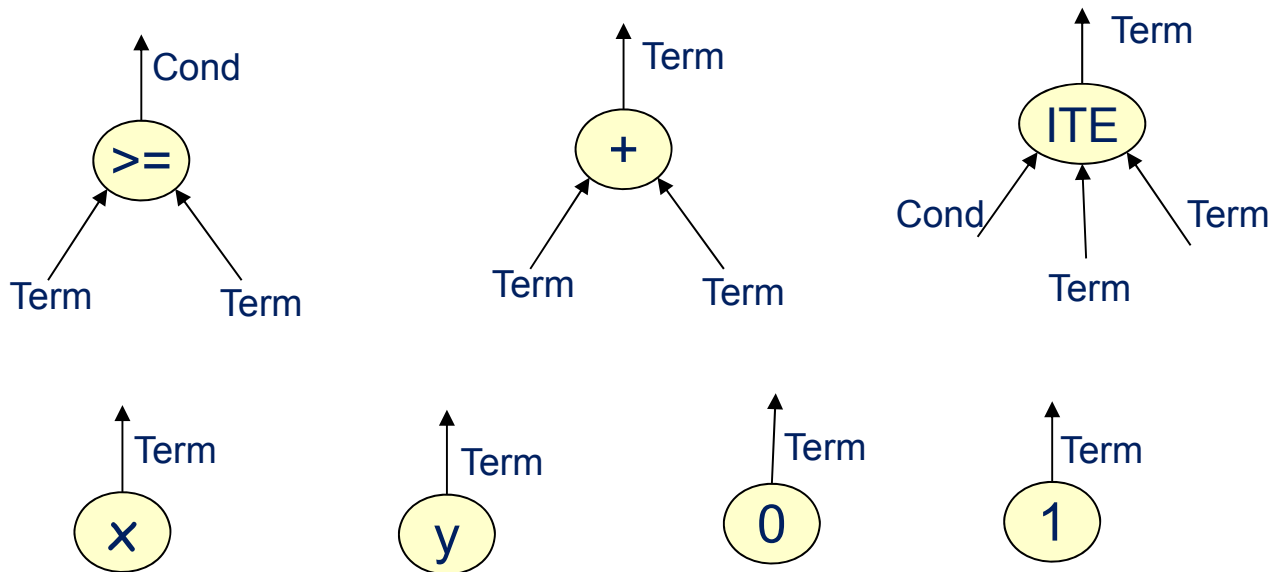
- Coming up: Learning strategies based on:
 - ◆ Enumerative (search with pruning): Udupa et al (PLDI'13)
 - ◆ Symbolic (solving constraints): Jha et al (ICSE'10,PLDI'11)
 - ◆ Stochastic (probabilistic walk): Schkufza et al (ASPLOS'13)

Enumerative Learning

- ❑ Find an expression consistent with a given set of concrete examples
- ❑ Enumerate expressions in increasing size, and evaluate each expression on all concrete inputs to check consistency
- ❑ Key optimization for efficient pruning of search space:
 - ◆ Expressions e_1 and e_2 are equivalent if $e_1(a,b)=e_2(a,b)$ on all concrete values $(x=a,y=b)$ in Examples
 - ◆ $(x+y)$ and $(y+x)$ always considered equivalent
 - ◆ If-Then-Else $(0 \leq x, e_1, e_2)$ considered equivalent to e_1 if in current set of Examples x has only non-negative values
 - ◆ Only one representative among equivalent subexpressions needs to be considered for building larger expressions
- ❑ Fast and robust for learning expressions with ~ 15 nodes

Symbolic Learning

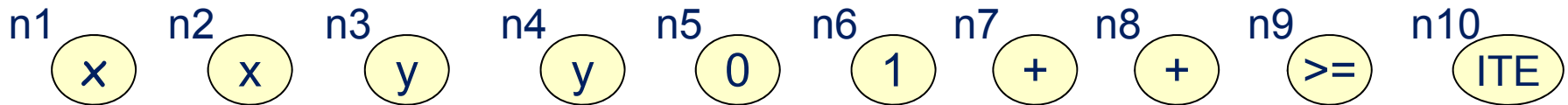
- Use a constraint solver for both the synthesis and verification steps
- Each production in the grammar is thought of as a component.
Input and Output ports of every component are typed.



- A well-typed loop-free program comprising these components corresponds to an expression DAG from the grammar.

Symbolic Learning

- Start with a library consisting of some number of occurrences of each component.



- Synthesis Constraints:

- ◆ Shape is a DAG, Types are consistent
- ◆ Spec $\varphi[f/e]$ is satisfied on every concrete input values in Examples

- Use an SMT solver (Z3) to find a satisfying solution.

- If synthesis fails, try increasing the number of occurrences of components in the library in an outer loop

Stochastic Learning

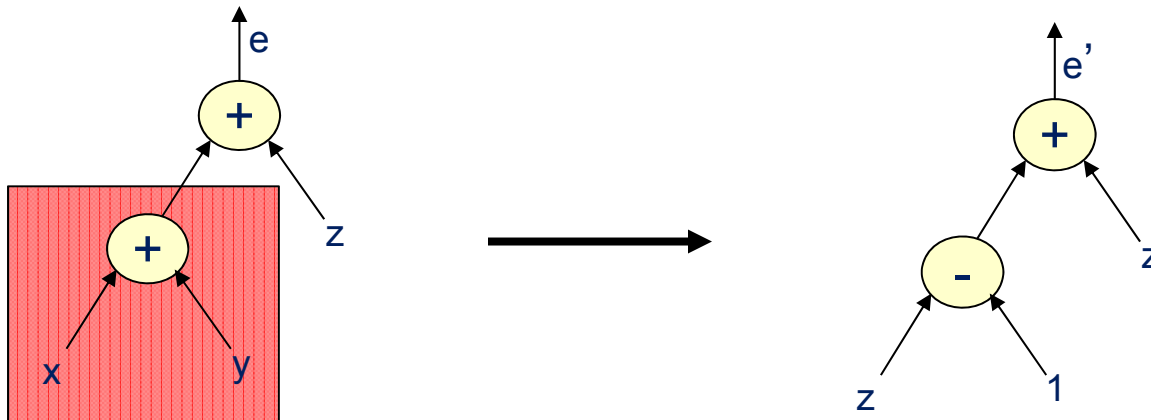
- ❑ Idea: Find desired expression e by probabilistic walk on graph where nodes are expressions and edges capture single-edits

- ❑ Metropolis-Hastings Algorithm: Given a probability distribution P over domain X , and an ergodic Markov chain over X , samples from X

- ❑ Fix expression size n .
 - ◆ X is the set of expressions E_n of size n .
 - ◆ $P(e) \propto \text{Score}(e)$ (“Extent to which e meets the spec φ ”)

Stochastic Learning

- ❑ Initial candidate expression e sampled uniformly from E_n
- ❑ If e works on all examples, return e
- ❑ Pick node v in parse tree of e uniformly at random. Replace subtree rooted at e with subtree of same size, sampled uniformly



- ❑ With probability $\min\{ 1, \text{Score}(e')/\text{Score}(e) \}$, replace e with e'
- ❑ Outer loop responsible for updating expression size n

Benchmarks and Implementation

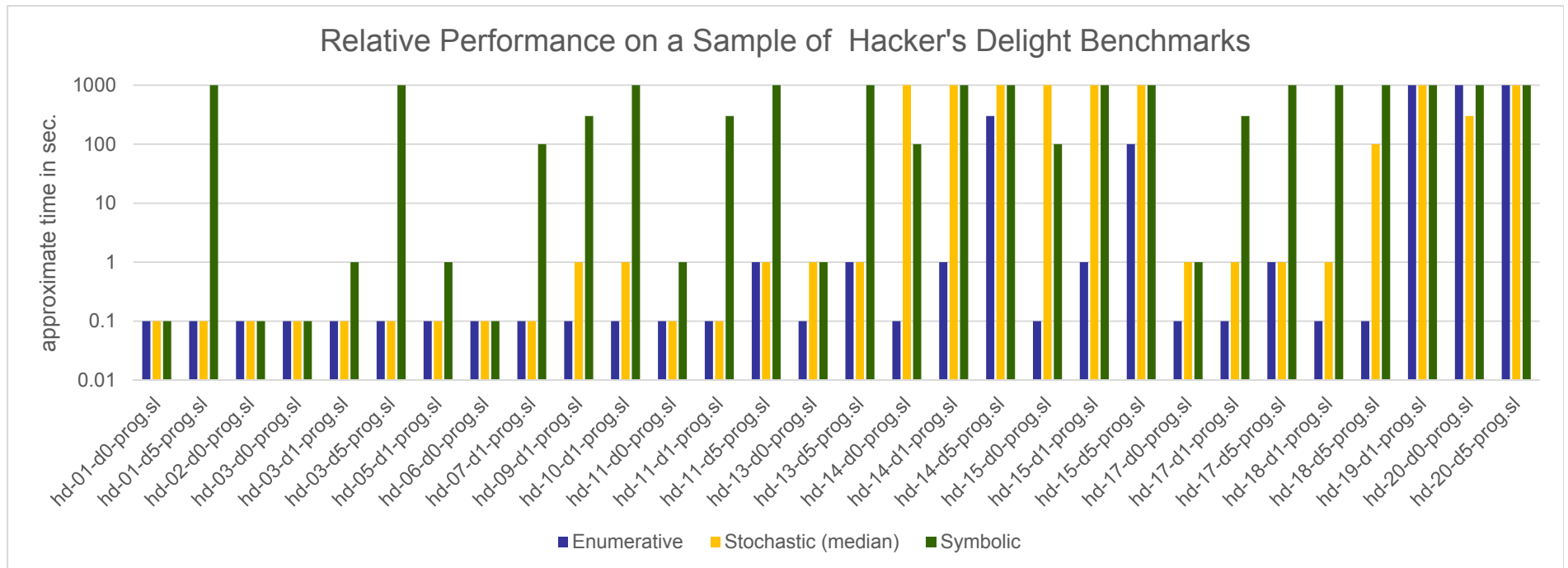
- ❑ Prototype implementation of Enumerative/Symbolic/Stochastic CEGIS

- ❑ Benchmarks:
 - ◆ Bit-manipulation programs from Hacker's delight
 - ◆ Integer arithmetic: Find max, search in sorted array
 - ◆ Challenge problems such as computing Morton's number

- ❑ Multiple variants of each benchmark by varying grammar

- ❑ Results are not conclusive as implementations are unoptimized, but offers first opportunity to compare solution strategies

Evaluation: Hacker's Delight Benchmarks



Evaluation Summary

- ❑ Enumerative CEGIS has best performance, and solves many benchmarks within seconds
 - Potential problem: Synthesis of complex constants
- ❑ Symbolic CEGIS is unable to find answers on most benchmarks
 - Caveat: Sketch succeeds on many of these
- ❑ Choice of grammar has impact on synthesis time
 - When E is set of all possible expressions, solvers struggle
- ❑ None of the solvers succeed on some benchmarks
 - Morton constants, Search in integer arrays of size > 4
- ❑ Bottomline: Improving solvers is a great opportunity for research !