



Shortest-Path Routing: Link-State & Distance-Vector

EE 122: Intro to Communication Networks

Fall 2007 (WF 4-5:30 in Cory 277)

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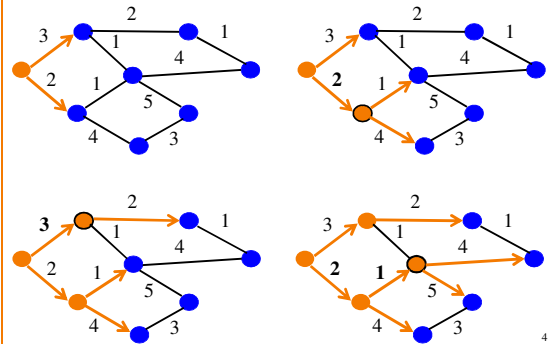
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<http://inst.eecs.berkeley.edu/~ee122/>

Materials with thanks to Jennifer Rexford

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Dijkstra's Algorithm Example



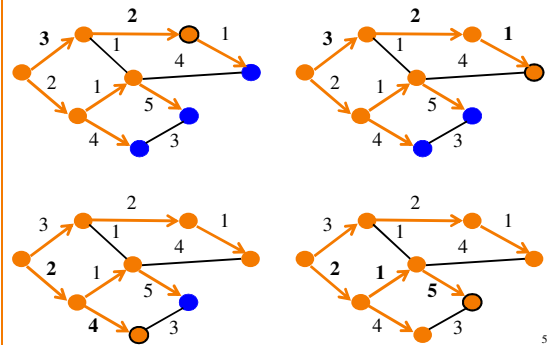
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Dijkstra's Shortest-Path Algorithm

- Iterative algorithm
 - After k iterations, know least-cost path to k nodes
- **S**: nodes whose least-cost path definitively known
 - Initially, $S = \{u\}$ where u is the source node
 - Add one node to S in each iteration
- **D(v)**: current cost of path from source to node v
 - Initially, $D(v) = c(u,v)$ for all nodes v adjacent to u
 - ... and $D(v) = \infty$ for all other nodes v
 - Continually update $D(v)$ as shorter paths are learned

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Dijkstra's Algorithm Example



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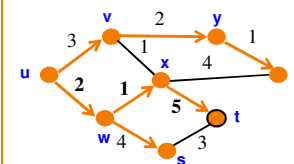
Dijkstra's Algorithm

- 1 **Initialization:**
- 2 $S = \{u\}$
- 3 for all nodes v
- 4 if v adjacent to u {
- 5 $D(v) = c(u,v)$
- 6 else $D(v) = \infty$
- 7 }
- 8 **Loop**
- 9 find w not in S with the smallest $D(w)$
- 10 add w to S
- 11 update $D(v)$ for all v adjacent to w and not in S :
- 12 $D(v) = \min\{D(v), D(w) + c(w,v)\}$
- 13 **until all nodes in S**

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Shortest-Path Tree

- Shortest-path tree from u
- Forwarding table at u



	link
v	(u,v)
w	(u,w)
x	(u,w)
y	(u,v)
z	(u,v)
s	(u,w)
t	(u,w)

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Distance Vector Algorithm

- $c(x,v)$ = cost for direct link from x to v
 - Node x maintains costs of direct links $c(x,v)$
- $D_x(y)$ = estimate of least cost from x to y
 - Node x maintains distance vector $D_x = [D_x(y) : y \in N]$
- Node x maintains its neighbors' distance vectors
 - For each neighbor v , x maintains $D_v = [D_v(y) : y \in N]$
- Each node v periodically sends D_v to its neighbors
 - And neighbors update their own distance vectors
 - $D_x(y) \leftarrow \min_v \{c(x,v) + D_v(y)\}$ for each node $y \in N$
- Over time, the distance vector D_x converges

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Distance Vector Example: Step 3

Optimum 3-hop paths

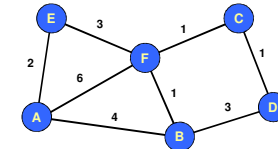


Table for A			Table for B		
Dst	Cst	Hop	Dst	Cst	Hop
A	0	A	A	4	A
B	4	B	B	0	B
C	6	E	C	2	F
D	7	B	D	3	D
E	2	E	E	4	F
F	5	E	F	1	F

Table for C			Table for D			Table for E			Table for F		
Dst	Cst	Hop	Dst	Cst	Hop	Dst	Cst	Hop	Dst	Cst	Hop
A	6	F	A	7	B	A	2	A	A	5	B
B	2	F	B	3	B	B	4	F	B	1	B
C	0	C	C	1	C	C	4	F	C	1	C
D	1	D	D	0	D	D	5	F	D	2	C
E	4	F	E	5	C	E	0	E	E	3	E
F	1	F	F	2	C	F	3	F	F	0	F

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Distance Vector Example: Step 1

Optimum 1-hop paths

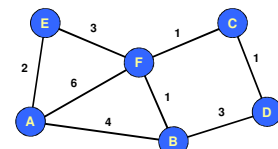


Table for A			Table for B		
Dst	Cst	Hop	Dst	Cst	Hop
A	0	A	A	4	A
B	4	B	B	0	B
C	∞	–	C	∞	–
D	∞	–	D	3	D
E	2	E	E	∞	–
F	6	F	F	1	F

Table for C			Table for D			Table for E			Table for F		
Dst	Cst	Hop	Dst	Cst	Hop	Dst	Cst	Hop	Dst	Cst	Hop
A	∞	–	A	∞	–	A	2	A	A	6	A
B	∞	–	B	3	B	B	∞	–	B	1	B
C	0	C	C	1	C	C	∞	–	C	1	C
D	1	D	D	0	D	D	∞	–	D	∞	–
E	∞	–	E	∞	–	E	0	E	E	3	E
F	1	F	F	∞	–	F	3	F	F	0	F

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Distance Vector Example: Step 2

Optimum 2-hop paths

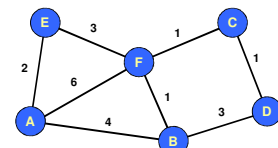


Table for A			Table for B		
Dst	Cst	Hop	Dst	Cst	Hop
A	0	A	A	4	A
B	4	B	B	0	B
C	7	F	C	2	F
D	7	B	D	3	D
E	2	E	E	4	F
F	5	E	F	1	F

Table for C			Table for D			Table for E			Table for F		
Dst	Cst	Hop	Dst	Cst	Hop	Dst	Cst	Hop	Dst	Cst	Hop
A	7	F	A	7	B	A	2	A	A	5	B
B	2	F	B	3	B	B	4	F	B	1	B
C	0	C	C	1	C	C	4	F	C	1	C
D	1	D	D	0	D	D	∞	–	D	2	C
E	4	F	E	∞	–	E	0	E	E	3	E
F	1	F	F	2	C	F	3	F	F	0	F

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