

Computer Science 70
Discrete Mathematics and Probability Theory
Hashing
Lecture 23



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1 Handout: notes



CS70 L23 Hashing (1)

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Big Idea: memoization

- General principle: store rather than recompute.
- Context is a tree-recursive algorithm with lots of repeated computation, e.g. Fibonacci:

```
int Fib (int n) {
  if (n==0 || n==1) {
    return n;
  } else if (we've computed n's value already) {
    return that value;
  } else {
    int value = Fib(n-1) + Fib(n-2);
    store (n, value);
    return value;
  }
}
```

- Pairs (n, value of Fib(n)) are stored in the table.



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Hash Function

- If what we want to memoize isn't a simple number, how do we convert it to a number to easily store it into a table?
- We need something that can help us map this data into an integer, to serve as an index into an array (used to store the table).
- This mapping function is called a hash function



http://en.wikipedia.org/wiki/Hash_function

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Writing hash functions - TTT (1)

- Let's consider Tic-Tac-Toe:
 - One player chooses X, the other chooses O
 - They take turns placing their piece on the board
 - Assume X goes first
 - Once a piece is placed, it isn't moved
 - The player who first gets 3-in-a-row wins
 - If the board gets filled up and nobody wins, it's a tie



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Writing hash functions - TTT (2)

- Writing a Tic-Tac-Toe hash function:

$$h \begin{pmatrix} \times & & \\ & \circ & \\ & & \times \end{pmatrix} = 13,205$$

- One idea is to ignore the 2D nature of the game and make it a 1D array of slots



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Writing hash functions - TTT (3)

- Think of each of the 9 slots as 1 of 3 values
 - Blank, O and X
 - Let's assign values 0, 1 and 2 to these



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Writing hash functions - TTT (4)

- Analysis of ternary polynomial hashcode:

- What's the smallest #? 0
- What's the biggest #? $3^9 - 1$
- Is this as optimal (I.e., tightly-packed) as possible? No!
- Any suggestions for making this more optimal?



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Writing hash functions - TTT (5)

- Optimizing the Tic-Tac-Toe hash function
 - This involves understanding the rules of placement
 - » The players take turns & X goes first!
 - Let's consider some small 1D boards ($S = \#$ of slots)
 - » $S=1$: 2 boards (- | X) We'll use "|" to separate groups
 - » $S=2$: 5 boards (-- | -X, X- | XO, OX)
 - » $S=3$: 13 boards (--- | --X, -X-, X-- | -OX, -XO, O-X, OX-, X-O, XO- | OXX, XOX, XXO) = $(1 + 3 + 6 + 3)$
 - » $S=4$: ___ boards (___ + ___ + ___ + ___)
 - » ...pattern?



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Remember your Combinatorics!

- Let's figure out $\text{numBoards}(s)$, $s = \# \text{ slots}$
- For $n=5$, we had:

ways to rearrange 0 Xs, 0 Os in 4 slots +
 # ways to rearrange 1 Xs, 0 Os in 4 slots +
 # ways to rearrange 1 Xs, 1 Os in 4 slots +
 # ways to rearrange 2 Xs, 1 Os in 4 slots +
 # ways to rearrange 2 Xs, 2 Os in 4 slots

- Generalizing from this example ($p = \# \text{ pieces}$):

$$\text{numBoards}(s) = \sum_{p=0}^{\lfloor s/2 \rfloor} \text{rearrange}(p, p, s) + \sum_{p=1}^{\lfloor s/2 \rfloor} \text{rearrange}(p, p-1, s)$$

- But what is $\text{rearrange}(x, o, s)$?

- # of ways to rearrange x Xs, o Os in s slots?



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Recall Pascal's Triangle $\binom{5}{2} = 10$

		0	1	2	3	4	5	6
0	1							
1	1	1						
2	1	2	1					
3	1	3	3	1				
4	1	4	6	4	1			
5	1	5	10	10	5	1		
6	1	6	15	20	15	6	1	

This table describes how to calculate combinations. I.e., "N choose K".

$$\binom{N}{K} = \frac{N!}{K!(N-K)!}$$

That is, the number of ways to rearrange 2 pieces in 5 slots is "5 choose 2", which is the expression at the top. 10 ways.



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$$\text{rearrange}(x, o, s) = r(x, o, s)$$

- How many ways to rearrange x Xs, o Os in s slots?



- Blur method

- First, blur eyes, how many ways to rearrange ALL $(x+o)$ pieces in s slots? [stop blurring now] $\binom{s}{x+o}$
 - For EACH, how many ways to rearrange Xs in pieces? $\binom{x+o}{x}$
 - Answer is product of these

$$\frac{s!}{(s-x-o)!(x+o)!} \cdot \frac{(x+o)!}{o!x!} = \binom{s}{x+o} \binom{x+o}{x}$$

- Overcount method

- Think of permuting all the elements; how many? $s!$
 - How many were overcounted? Xs, Os, spaces $o!x!(s-x-o)!$
 - Answer is quotient of these

$$\frac{s!}{o!x!(s-x-o)!}$$



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Now we know our Hash Table size

- Now we know $\text{numBoards}(s)$

- $\text{numBoards}(4) \Rightarrow (1 + 4 + 12 + 12 + 6) = 35$

- $\text{numBoards}(9) \Rightarrow (1 + 9 + 72 + 252 + 756 + 1260 + 1680 + 1260 + 630 + 126) = 6,046 < 19,683 = 3^9$

- Plotting $\text{rearrange}(x, o, 4)$

		2		6
o	1	12	12	
	0	1	4	
s=4	0	1	2	x

Note zig-zag pattern as a result of the alternating moves of each player! numBoards just sums 'em!



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But what about the hash function?

- How do we write the **combinatorially optimal** hash ()?
 - This take our board and generates a # between 0 and (numBoards - 1)
- **Two steps (sum the following numbers)**
 1. Finding out how many numbers there were in the zigzag up to our box (this is the **BIAS**, or **OFFSET**)
 2. Finding out our number **REARRANGEMENT** within our box
 - » Exactly same idea as the ternary polynomial hash code:
 - X counts as 2, i.e., $2 \cdot 3^i$, O counts as 1, i.e., $1 \cdot 3^i$, - = 0
 - » Here, we consider the leftmost slot & how much it's worth
 - X counts for all ways to rearrange if it were O & -
 - O counts for all ways to rearrange if it were -
 - - counts for 0
 - (Shortcut when a board has all the same piece, counts for 0)



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Example TicTacToe hash function

- Let's hash **XO-X** = $X_3 O_2 -_1 X_0$
 - Must be a # between 0 and (numBoards (4) - 1) = 34
- **Two steps: BIAS + REARRANGEMENT #**
 - **BIAS:** X=2,O=1,S=4; Count buckets up to us: $1+4+12=17$
 - **REARRANGEMENT #:** [R(X,O,S)]
 - » $X_3 = r(2,1,3) + r(2,0,3) = 3 + 3$
 - » $O_2 = r(1,1,2) = 2$
 - » $-_1 = 0$
 - » $X_0 = 0$ (from shortcut)
 - » **REARRANGEMENT # = $3 + 3 + 2 = 8$**
- Thus, **combinatorially optimal** hash(**XO-X**) = $17 + 8 = 25$

	2			6
O	1		12	12
	0	1	4	
S=4	0	1	2	



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Summary

- We showed how to calculate **combinatorially optimal** hash functions for a game
 - In real-world applications, we often find this useful
 - If it's too expensive, we usu. settle for sub-optimal
- A good hash function spreads out values evenly
- Sometimes hard to write good hash function
 - In 8 real applications, 2 had written poor hash funs
- Java has a **great** hash function for Strings
 - Strings are commonly used as the keys (the things you hash upon for a data structure)



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