

# Computer Science 70

## Discrete Mathematics and Probability Theory

### Hashing

### Lecture 23



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**1 Handout:** notes



# Big Idea: memoization

- General principle: store rather than recompute.
- Context is a tree-recursive algorithm with lots of repeated computation, e.g. Fibonacci:

```
int Fib (int n) {  
    if (n==0 || n==1) {  
        return n;  
    } else if (we've computed n's value already) {  
        return that value;  
    } else {  
        int value = Fib(n-1) + Fib(n-2);  
        store (n, value);  
        return value;  
    }  
}
```

- Pairs (n, value of Fib(n)) are stored in the table.



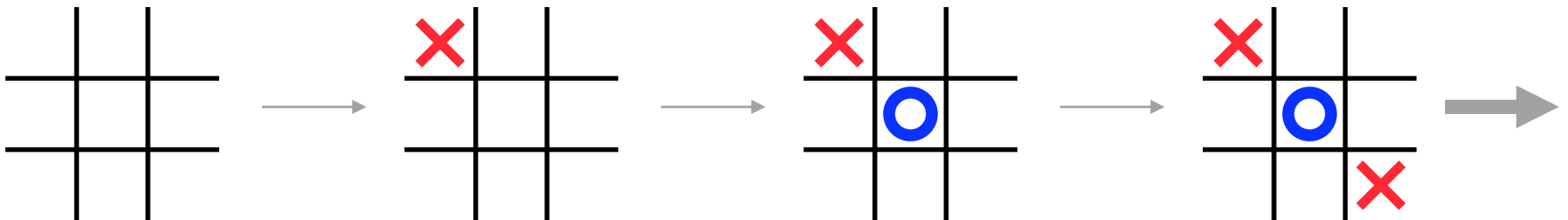
# Hash Function

- If what we want to memoize isn't a simple number, how do we convert it to a number to easily store it into a table?
- We need something that can help us **map** this data into an integer, to serve as an index into an array (used to store the table).
- This mapping function is called a **hash function**



# Writing hash functions - TTT (1)

- Let's consider Tic-Tac-Toe:
  - One player chooses **X**, the other chooses **O**
  - They take turns placing their piece on the board
  - Assume **X** goes first
  - Once a piece is placed, it isn't moved
  - The player who first gets 3-in-a-row wins
  - If the board gets filled up and nobody wins, it's a tie

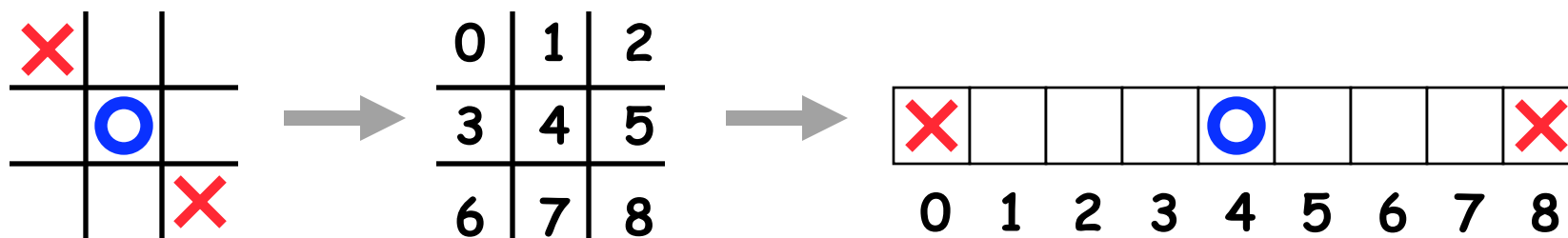


# Writing hash functions - TTT (2)

- Writing a Tic-Tac-Toe hash function:

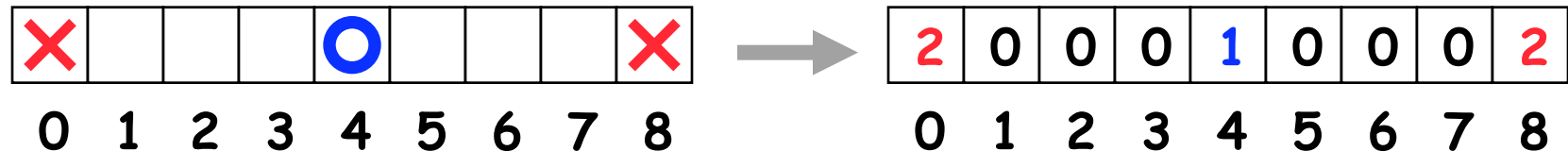
$$h \left( \begin{array}{|c|c|c|} \hline \times & & \\ \hline & \circ & \\ \hline & & \times \\ \hline \end{array} \right) = 13,205$$

- One idea is to ignore the 2D nature of the game and make it a 1D array of slots



# Writing hash functions - TTT (3)

- Think of each of the 9 slots as 1 of 3 values
  - Blank, O and X
  - Let's assign values 0, 1 and 2 to these



# Writing hash functions - TTT (4)

- Analysis of ternary polynomial hashcode:

- What's the smallest #?

0

- What's the biggest #?

$3^9-1$

- Is this as optimal (I.e., tightly-packed) as possible?

No!

- Any suggestions for making this more optimal?



# Writing hash functions - TTT (5)

- **Optimizing** the Tic-Tac-Toe hash function
  - This involves understanding the rules of placement
    - » The players take turns & X goes first!
  - Let's consider some small 1D boards ( $S = \#$  of slots)
    - »  $S=1$ : 2 boards (- | X) We'll use "|" to separate groups
    - »  $S=2$ : 5 boards (-- | -X, X- | XO, OX)
    - »  $S=3$ : 13 boards (--- | --X, -X-, X-- | -OX, -XO, O-X, OX-, X-O, XO- | OXX, XOX, XXO) =  $(1 + 3 + 6 + 3)$
    - »  $S=4$ : \_\_\_ boards (\_\_\_ + \_\_\_ + \_\_\_ + \_\_\_ + \_\_\_)
    - » ...pattern?





# Remember your Combinatorics!

- Let's figure out  $\text{numBoards}(s)$ ,  $s = \# \text{ slots}$
- For  $n=5$ , we had:

# ways to rearrange 0 Xs, 0 Os in 4 slots +  
# ways to rearrange 1 Xs, 0 Os in 4 slots +  
# ways to rearrange 1 Xs, 1 Os in 4 slots +  
# ways to rearrange 2 Xs, 1 Os in 4 slots +  
# ways to rearrange 2 Xs, 2 Os in 4 slots

- Generalizing from this example ( $p = \# \text{ pieces}$ ):

$$\text{numBoards}(s) = \sum_{p=0}^{\lceil s/2 \rceil} \text{rearrange}(p, p, s) + \sum_{p=1}^{\lfloor s/2 \rfloor} \text{rearrange}(p, p-1, s)$$

- But what is  $\text{rearrange}(x, o, s)$ ?
  - # of ways to rearrange  $x$  Xs,  $o$  Os in  $s$  slots?



# Recall Pascal's Triangle $\binom{5}{2}=10$

	0	1	2	3	4	5	6
0	1						
1	1	1					
2	1	2	1				
3	1	3	3	1			
4	1	4	6	4	1		
5	1	5	10	10	5	1	
6	1	6	15	20	15	6	1

This table describes how to calculate combinations. I.e., "N choose K".

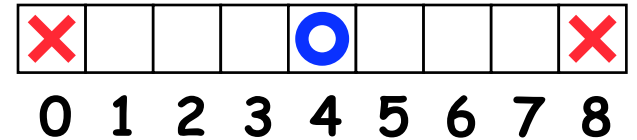
$$\binom{N}{K} = \frac{N!}{K!(N-K)!}$$

That is, the number of ways to rearrange 2 pieces in 5 slots is "5 choose 2", which is the expression at the top. 10 ways.



$$\text{rearrange}(x, o, s) = r(x, o, s)$$

- How many ways to rearrange  $x$  **Xs**,  $o$  **Os** in  $s$  **slots**?



- **Blur method**

- First, blur eyes, how many ways to rearrange ALL  $(x+o)$  pieces in  $s$  slots? [stop blurring now]  $\binom{s}{x+o}$
- For EACH, how many ways to rearrange Xs in pieces?  $\binom{x+o}{x}$
- Answer is product of these

$$\frac{s!}{(s-x-o)!(x+o)!} \cdot \frac{(x+o)!}{o! x!} = \binom{s}{x+o} \binom{x+o}{x}$$

- **Overcount method**

- Think of permuting **all** the elements; how many?  $s!$
- How many were overcounted? Xs, Os, spaces  $o! x! (s-x-o)!$
- Answer is quotient of these

$$\frac{s!}{o! x! (s-x-o)!}$$



# Now we know our Hash Table size

- Now we know numBoards (s)
  - numBoards (4)  $\Rightarrow (1 + 4 + 12 + 12 + 6) = 35$
  - numBoards (9)  $\Rightarrow (1+9+72+252+756+1260+1680+1260+630+126) = 6,046 < 19,683 = 3^9$
- Plotting rearrange (x, o, 4)

Note zig-zag pattern as a result of the alternating moves of each player! numBoards just sums 'em!

o

2			6
1		12	12
0	1	4	
s=4	0	1	2

x



# But what about the hash function?

- How do we write **the combinatorially optimal** hash ()?
  - This take our board and generates a # between 0 and (numBoards - 1)
- **Two steps (sum the following numbers)**
  1. Finding out how many numbers there were in the zigzag **up to our box** (this is the **BIAS**, or **OFFSET**)
  2. Finding out our number **REARRANGEMENT** **within our box**
    - » Exactly same idea as the ternary polynomial hash code:
      - **X** counts as 2, i.e.,  $2 \cdot 3^i$ , **O** counts as 1, I.e.,  $1 \cdot 3^i$ , **-** = 0
    - » Here, we consider the leftmost slot & how much it's worth
      - **X** counts for all ways to rearrange **if it were O** & **-**
      - **O** counts for all ways to rearrange **if it were -**
      - **-** counts for 0
      - (Shortcut when a board has all the same piece, counts for 0)



# Example TicTacToe hash function

- Let's hash  $XO-X = X_3O_2^{-1}X_0$ 
  - Must be a # between 0 and  $(\text{numBoards } (4) - 1) = 34$
- Two steps: **BIAS + REARRANGEMENT #**
  - BIAS:  $X=2, O=1, S=4$ ; Count buckets up to us:  $1+4+12=17$
  - REARRANGEMENT #:  $[R(X, O, S)]$ 
    - »  $X_3 = r(2, 1, 3) + r(2, 0, 3) = 3 + 3$
    - »  $O_2 = r(1, 1, 2) = 2$
    - »  $^{-1} = 0$
    - »  $X_0 = 0$  (from shortcut)
    - » REARRANGEMENT # =  $3 + 3 + 2 = 8$
- Thus, combinatorially optimal  
 $\text{hash}(XO-X) = 17 + 8 = 25$

	2		6
O	1	12	12
	0	1	4
S=4	0	1	2
		X	



# Summary

- We showed how to calculate **combinatorially optimal** hash functions for a game
  - In real-world applications, we often find this useful
  - If it's too expensive, we usu. settle for sub-optimal
- A good hash function spreads out values evenly
- Sometimes hard to write good hash function
  - In 8 real applications, 2 had written poor hash funs
- Java has a **great** hash function for Strings
  - Strings are commonly used as the keys (the things you hash upon for a data structure)

