## CS 70 Discrete Mathematics for CS Spring 2008 David Wagner

# Due Thursday, March 6th

### 1. (10 pts.) Break the generator, become a poker champion

A *pseudorandom number generator* is a way of generating a large quantity of random-looking numbers, if all we have is a little bit of randomness (known as the *seed*). One simple scheme is the *linear congruential generator*, where we let  $x_0$  denote the seed and define

 $x_{t+1} = (ax_t + b) \mod m$ 

for some modulus *m* and some constants *a*,*b*. (Notice that  $0 \le x_t < m$  holds for every *t*.)

You've discovered that a popular web site uses a linear congruential generator to generate poker hands for its players. For instance, it uses  $x_0$  to pseudo-randomly pick the first card to go into your hand,  $x_1$  to pseudo-randomly pick the second card to go into your hand, and so on. For extra security, the poker site has kept the parameters *a* and *b* secret, but you do know that the modulus is  $m = 2^{31} - 1$  (which is prime).

Suppose that you can observe the values  $x_0$ ,  $x_1$ ,  $x_2$ ,  $x_3$ , and  $x_4$  from the information available to you, and that the values  $x_5, \ldots, x_9$  will be used to pseudo-randomly pick the cards for the next person's hand. Describe how to efficiently predict the values  $x_5, \ldots, x_9$ , given the values known to you.

#### 2. (5 pts.) Interpolation practice

Find a polynomial  $h(x) = ax^2 + bx + c$  of degree at most 2 such that  $h(0) \equiv 3 \pmod{7}$ ,  $h(1) \equiv 6 \pmod{7}$ , and  $h(2) \equiv 6 \pmod{7}$ .

*Hint:* You could write a system of linear equations in the unknowns a, b, c, or use Question 4 below.

#### **3.** (10 pts.) Rational interpolation

Find a non-zero polynomial  $g(x) = ax^2 + bx + c$  of degree at most 2 and a non-zero polynomial h(x) = dx + e of degree at most 1 such that  $g(0) \equiv 3 \cdot h(0) \pmod{7}$ ,  $g(1) \equiv 3 \cdot h(1) \pmod{7}$ ,  $g(2) \equiv 2 \cdot h(2) \pmod{7}$ , and  $g(3) \equiv 4 \cdot h(3) \pmod{7}$ .

Hint: You might try setting up a system of linear equations and solving it.

*Important note:* The solution  $a \equiv b \equiv c \equiv d \equiv e \equiv 0 \pmod{7}$  doesn't count. Find any other solution. (Revised 3/2/2008 to add this requirement.)

#### 4. (25 pts.) Polynomial interpolation

In this problem, you will develop an algorithm for polynomial interpolation: given a prime p and n values  $a_0, a_1, \ldots, a_n \pmod{p}$  (where n < p), you will show how to construct a polynomial h(x) of degree at most n such that  $h(i) \equiv a_i \pmod{p}$  for  $i = 0, \ldots, n$ .

- (a) Prove the following: If *p* is a prime and  $y_1, \ldots, y_n \in \mathbb{N}$  are all different from 0 modulo *p*, then  $y_1 \times \cdots \times y_n$  is also different from 0 modulo *p*.
- (b) Prove the following: Given a prime p and two integers a, b, it is always possible to find a polynomial f(x) of degree at most 1 such that  $f(0) \equiv a \pmod{p}$  and  $f(1) \equiv b \pmod{p}$ .

(c) You are given a prime *p* and a positive number n < p. Show how to find a polynomial f(x) of degree at most *n* satisfying  $f(0) \equiv f(1) \equiv \cdots \equiv f(n-1) \equiv 0 \pmod{p}$  and  $f(n) \equiv 1 \pmod{p}$ . In other words, the polynomial *f* should be congruent to zero at the points  $x = 0, \dots, n-1$ ; at x = n the polynomial should be 1 mod *p*.

*Hint:* Consider  $F(x) = (x-0)(x-1)(x-2)\cdots(x-(n-1))$ ; what can you say about it?

(d) You are given *p* and *n* as before, but now you are also given an index *j* with  $0 \le j \le n$ . Show how to find a polynomial  $g_j(x)$  of degree at most *n* satisfying

$$g_j(i) \equiv \begin{cases} 0 & \text{if } i \neq j \\ 1 & \text{if } i = j \end{cases} \pmod{p} \quad \text{for each } i = 0, 1, \dots, n.$$

In other words, the polynomial  $g_j$  should be congruent to zero at the points x = 0, ..., n, except that at x = j it should be congruent to 1 mod p.

(e) You are given a prime p, a number n with 0 < n < p, and a sequence of values a<sub>0</sub>, a<sub>1</sub>,..., a<sub>n</sub> (mod p). Describe an efficient algorithm to find a polynomial h(x) of degree at most n satisfying h(0) ≡ a<sub>0</sub> (mod p), h(1) ≡ a<sub>1</sub> (mod p), ..., h(n) ≡ a<sub>n</sub> (mod p).

*Hint:* What can you say about the polynomial  $3g_0(x) + 7g_1(x)$ , where  $g_0(x), g_1(x)$  are as defined in part (d)? Does this give you any ideas?