

Concurrency Control and Performance

Agrawal/Carey/Livny: Locking vs. Optimistic

Previous work had conflicting results:

- Carey & Stonebraker (VLDB84), Agrawal & DeWitt (TODS85): blocking beats restarts
- Tay (Harvard PhD) & Balter (PODC82): restarts beat blocking
- Franaszek & Robinson (TODS85): optimistic beats locking

Goal of this paper:

- Do a good job modeling the problem and its variants
- Capture causes of previous conflicting results
- Make recommendations based on variables of problem

Methodology

- simulation study, compare Blocking (i.e. 2PL), Immediate Restart (restart when denied a lock), and Optimistic (a la Kung & Robinson)
- pay attention to model of system:
 - database system model: hardware and software model (CPUs, disks, size & granule of DB, load control mechanism, CC algorithm)
 - user model: arrival of user tasks, nature of tasks (e.g. batch vs. interactive)
 - transaction model: logical reference string (i.e. CC schedule), physical reference string (i.e. disk block requests, CPU processing bursts).
 - Probabilistic modeling of each. They argue this is key to a performance study of a DBMS.
- logical queueing model
- physical queueing model

Measurements

- measure throughput, mostly
- pay attention to variance of response time, too
- pick a DB size so that there are noticeable conflicts (else you get comparable performance)

Experiment 1: Infinite Resources

- as many disks and CPUs as you want
- blocking thrashes due to transactions blocking numerous times
- restart plateaus: adaptive wait period (avg response time) before restart
 - serves as a primitive load control!
- optimistic scales logarithmically
- standard deviation of response time under locking much lower

Experiment 2: Limited Resources (1 CPU, 2 disks)

- Everybody thrashes
- blocking throughput peaks at mpl 25
- optimistic peaks at 10
- restart peaks at 10, plateaus at 50 – as good or better than optimistic
- at super-high mpl (200), restart beats both blocking and optimistic

- but total throughput worse than blocking @ mpl 25
- effectively, restart is achieving mpl 60
- load control is the answer here – adding it to blocking & optimistic makes them handle higher mpls better

Experiment 3: Multiple Resources (5, 10, 25, 50 CPUs, 2 disks each)

- optimistic starts to win at 25 CPUs
 - when useful disk utilization is only about 30%, system begins to behave like infinite resources
- even better at 50

Experiment 4: Interactive Workloads

Add user think time.

- makes the system appear to have more resources
- so optimistic wins with think times 5 & 10 secs. Blocking still wins for 1 second think time.

Questioning 2 assumptions:

- fake restart – biases for optimistic
 - fake restarts result in less conflict.
 - cost of conflict in optimistic is higher
 - issue of $k > 2$ transactions contending for one item
 - will have to punish $k-1$ of them with real restart
- write-lock acquisition
 - recall our discussion of lock upgrades and deadlock
 - blind write biases for restart (optimistic not an issue here), particularly with infinite resources (blocking holds write locks for a long time; waste of deadlock restart not an issue here).
 - with finite resources, blind write restarts transactions earlier (making restart look better)

Conclusions

- blocking beats restarting, unless resource utilization is low
- possible in situations of high think time
- mpl control important. admission control the typical scheme.
 - Restart's adaptive load control is too clumsy, though.
- false assumptions made blocking look relatively worse

Final quote by Wulf!