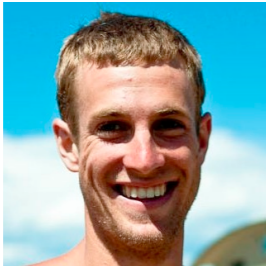


Goal-Directed Program Analysis with Jumping



Sam Blackshear
University of Colorado Boulder



Bor-Yuh Evan Chang
University of Colorado Boulder



Manu Sridharan
Samsung Research America

Google
July 24, 2015

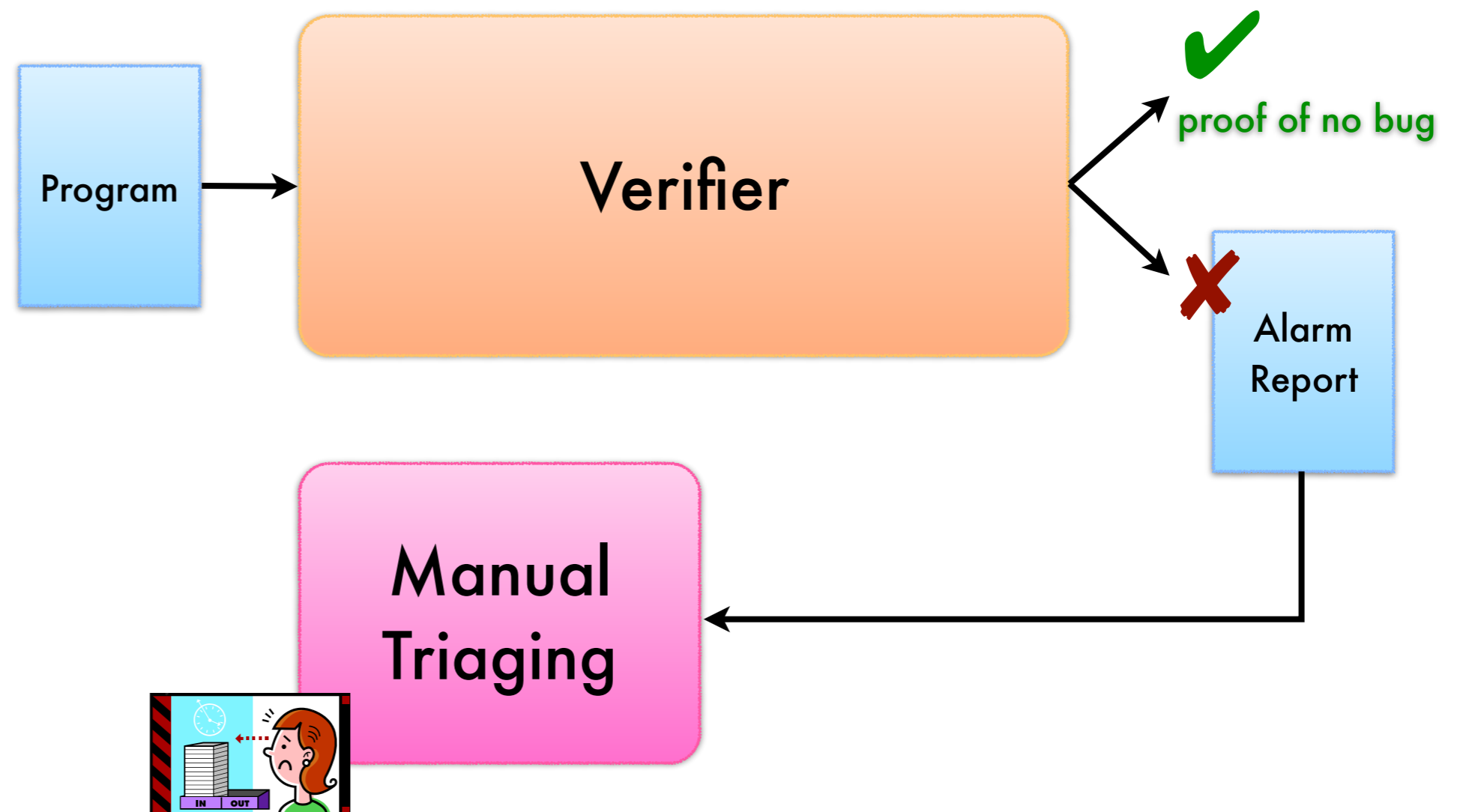


Lab: **Program analysis** in the **whole** bug mitigation **process**

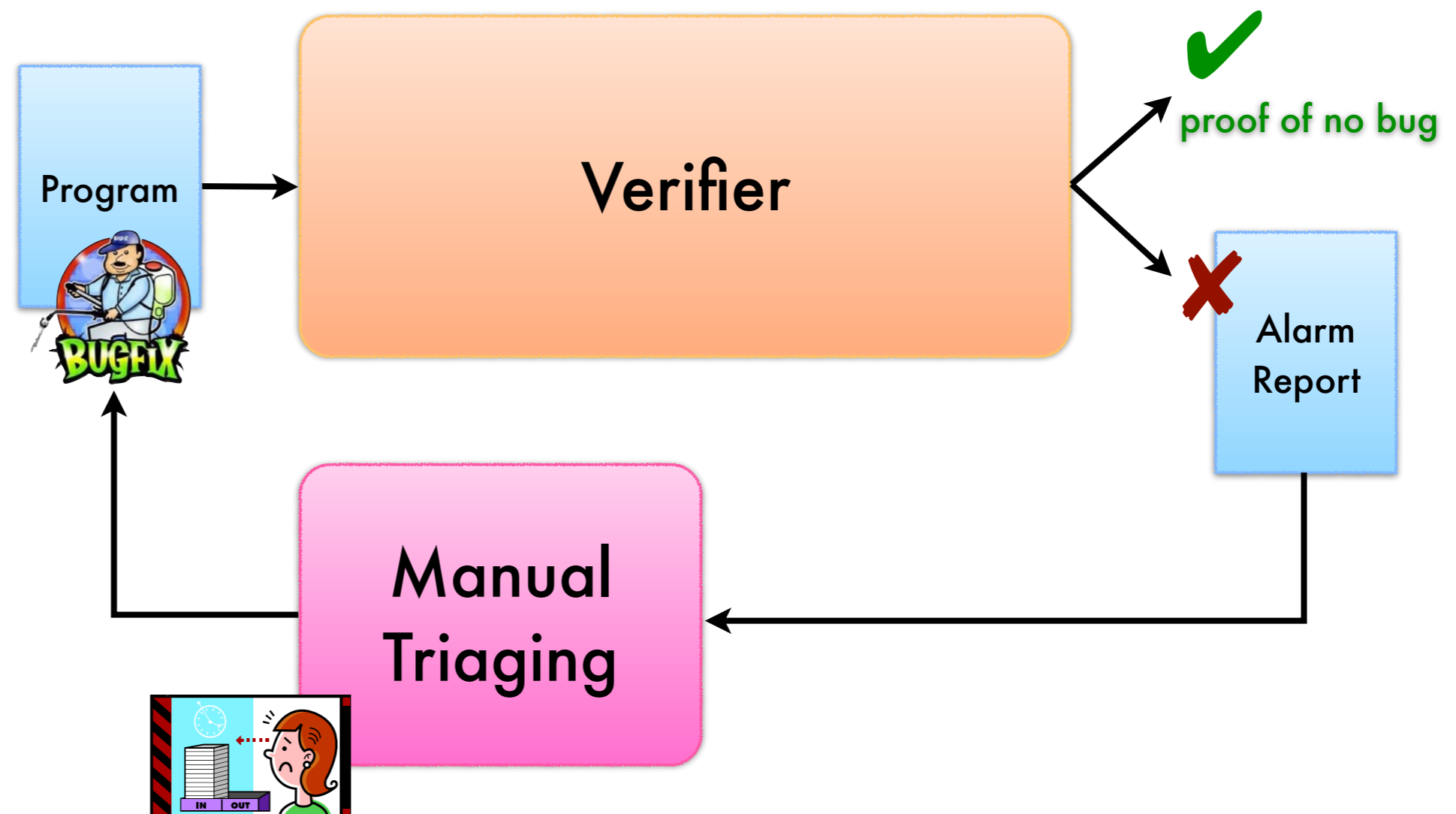
Lab: Program analysis in the whole bug mitigation process



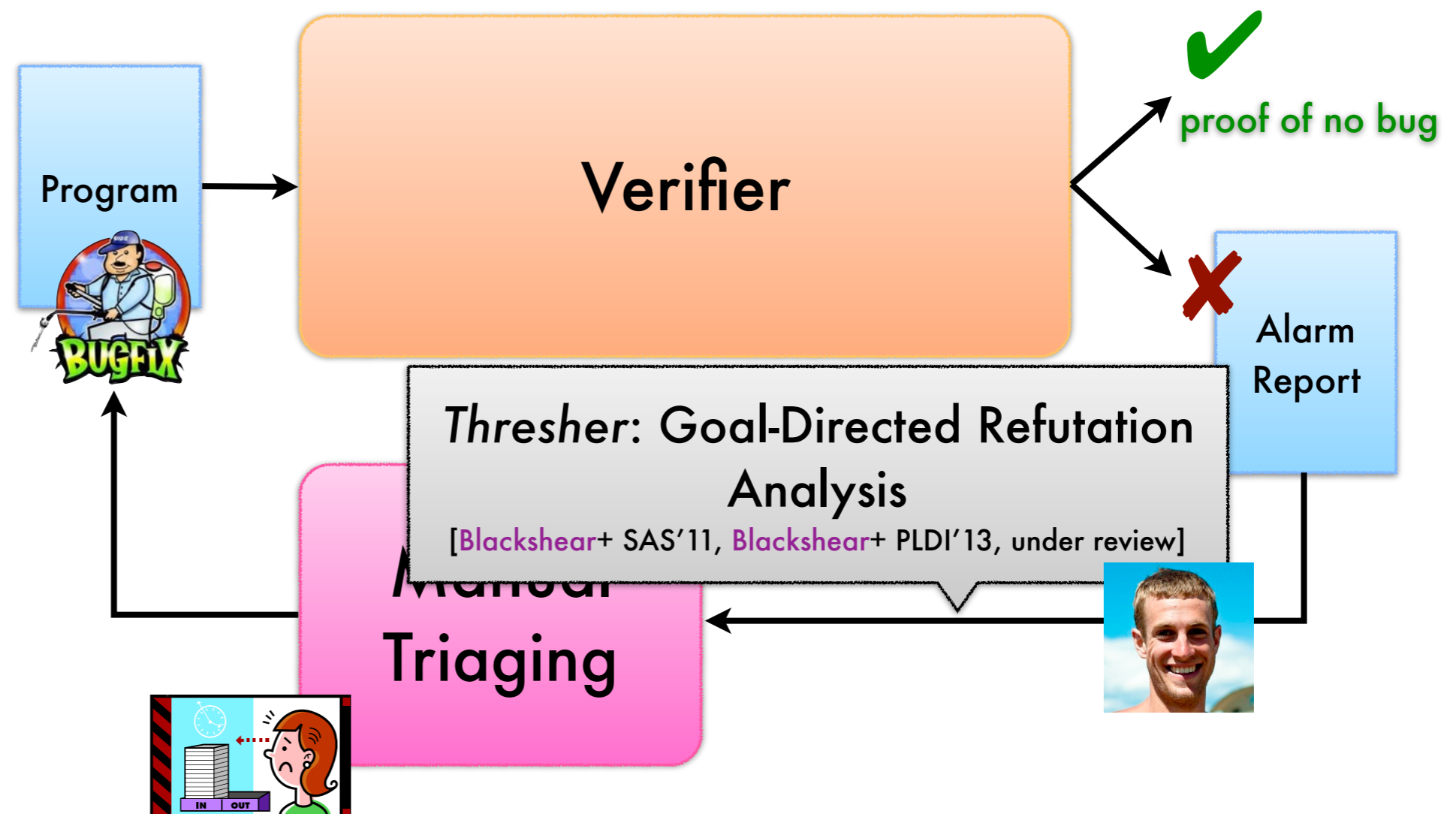
Lab: Program analysis in the whole bug mitigation process



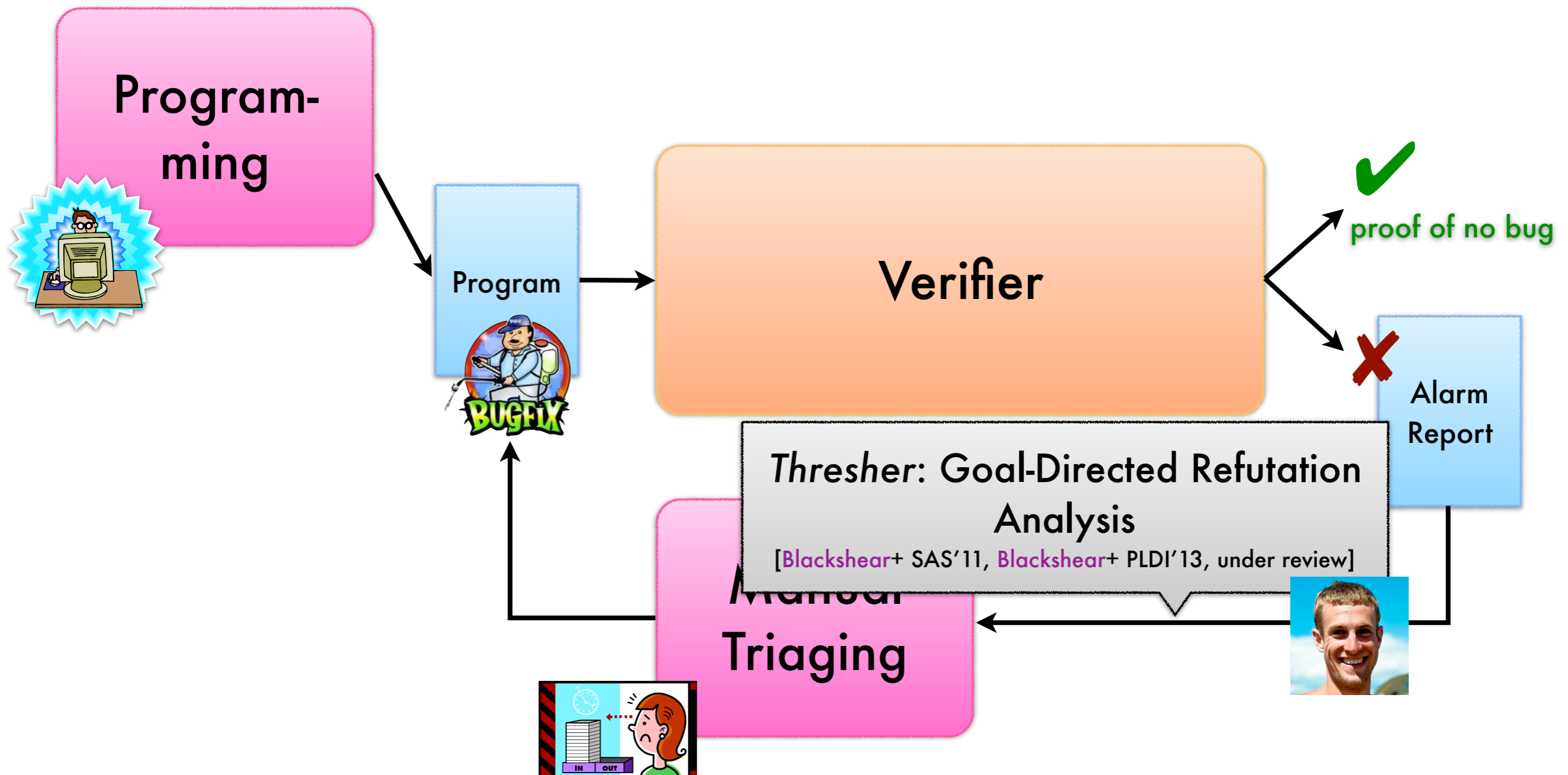
Lab: Program analysis in the whole bug mitigation process



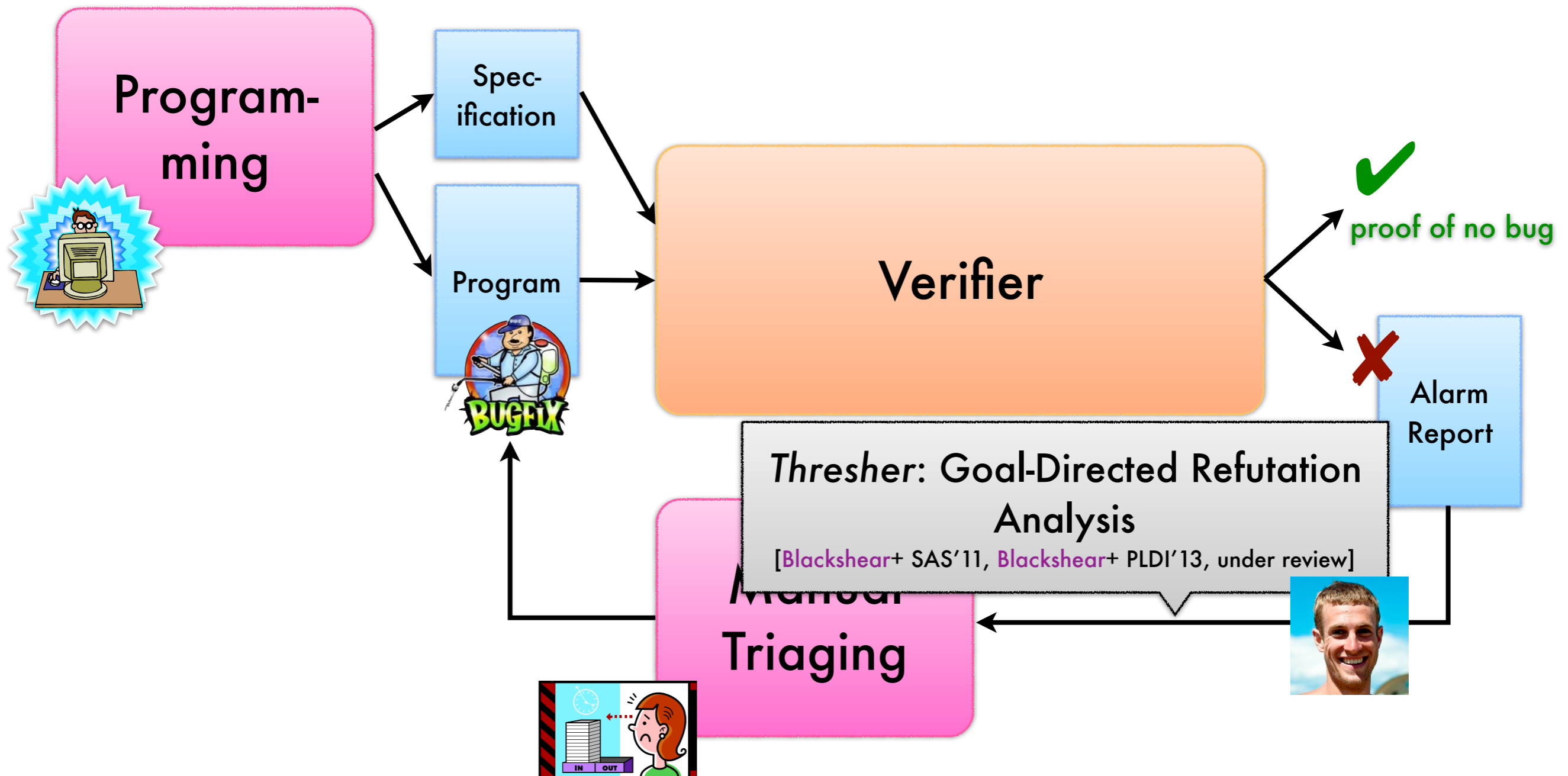
Lab: Program analysis in the whole bug mitigation process



Lab: Program analysis in the whole bug mitigation process



Lab: Program analysis in the whole bug mitigation process



Analysis in the whole bug mitigation process



**Fissile Types:
Checking Almost
Everywhere
Invariants**
[Coughlin+ POPL'14, in prep]

**Program-
ming**

Spec-
ification

Program



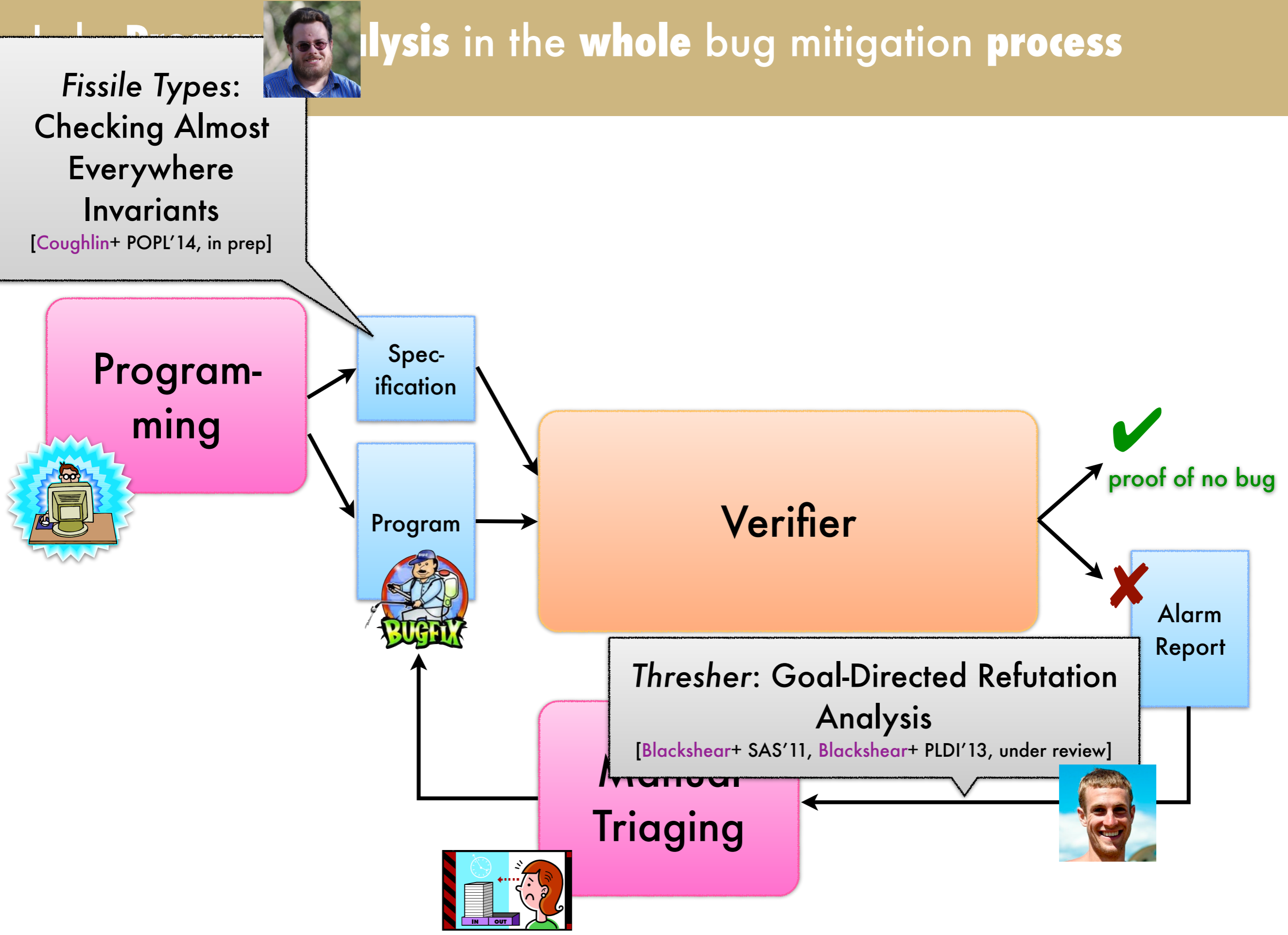
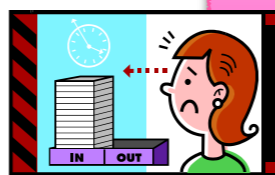
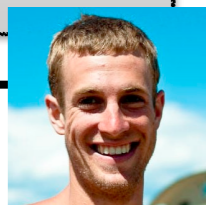
Verifier

✓
proof of no bug

✗
Alarm
Report

**Thresher: Goal-Directed Refutation
Analysis**
[Blackshear+ SAS'11, Blackshear+ PLDI'13, under review]

**Manual
Triaging**





Analysis in the whole bug mitigation process

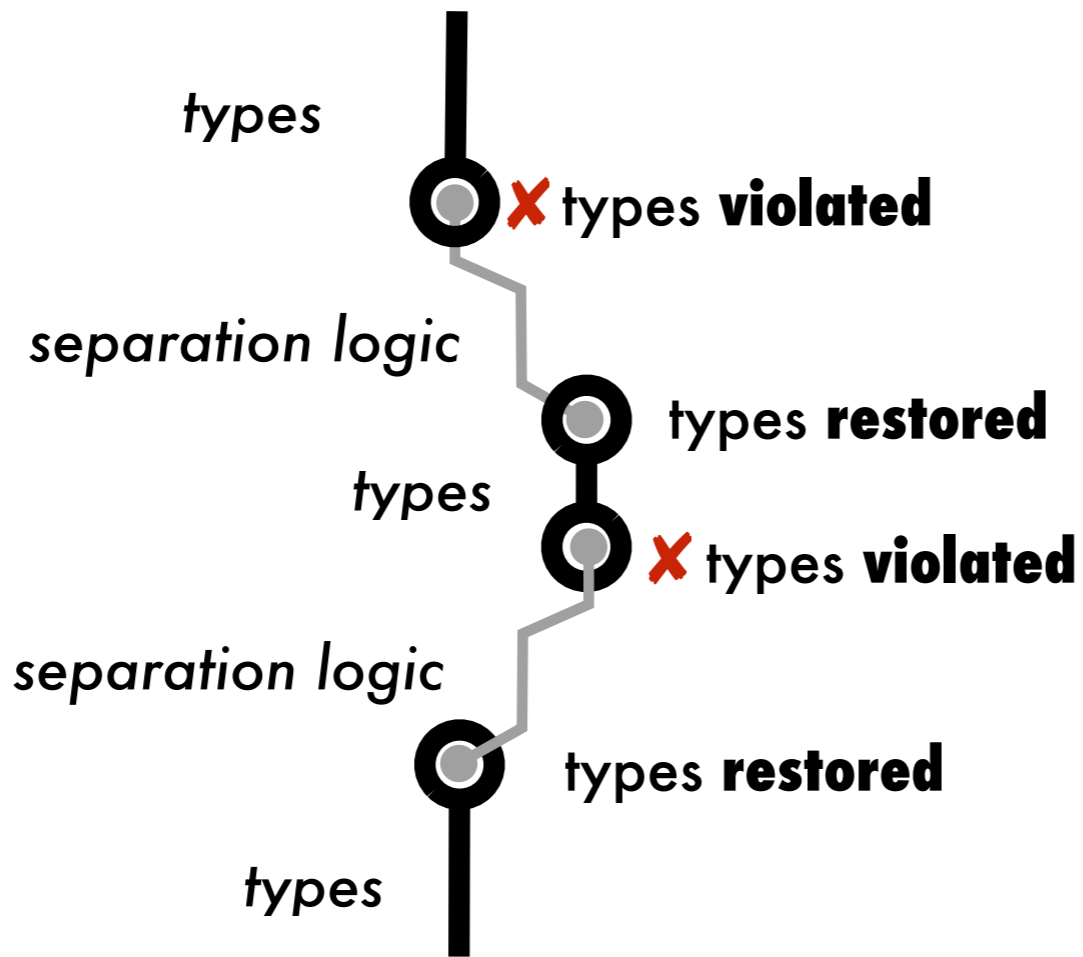
**Fissile Types:
Checking Almost
Everywhere
Invariants**

[Coughlin+ POPL'14, in prep]

Program



analysis time



type-intertwined separation logic

Intertwine type-based and separation-based analysis to verify **almost-everywhere** invariants



Objective-C

no bug

m
ort

Analysis in the whole bug mitigation process



**Fissile Types:
Checking Almost
Everywhere
Invariants**
[Coughlin+ POPL'14, in prep]

**Program-
ming**



Spec-
ification

Program



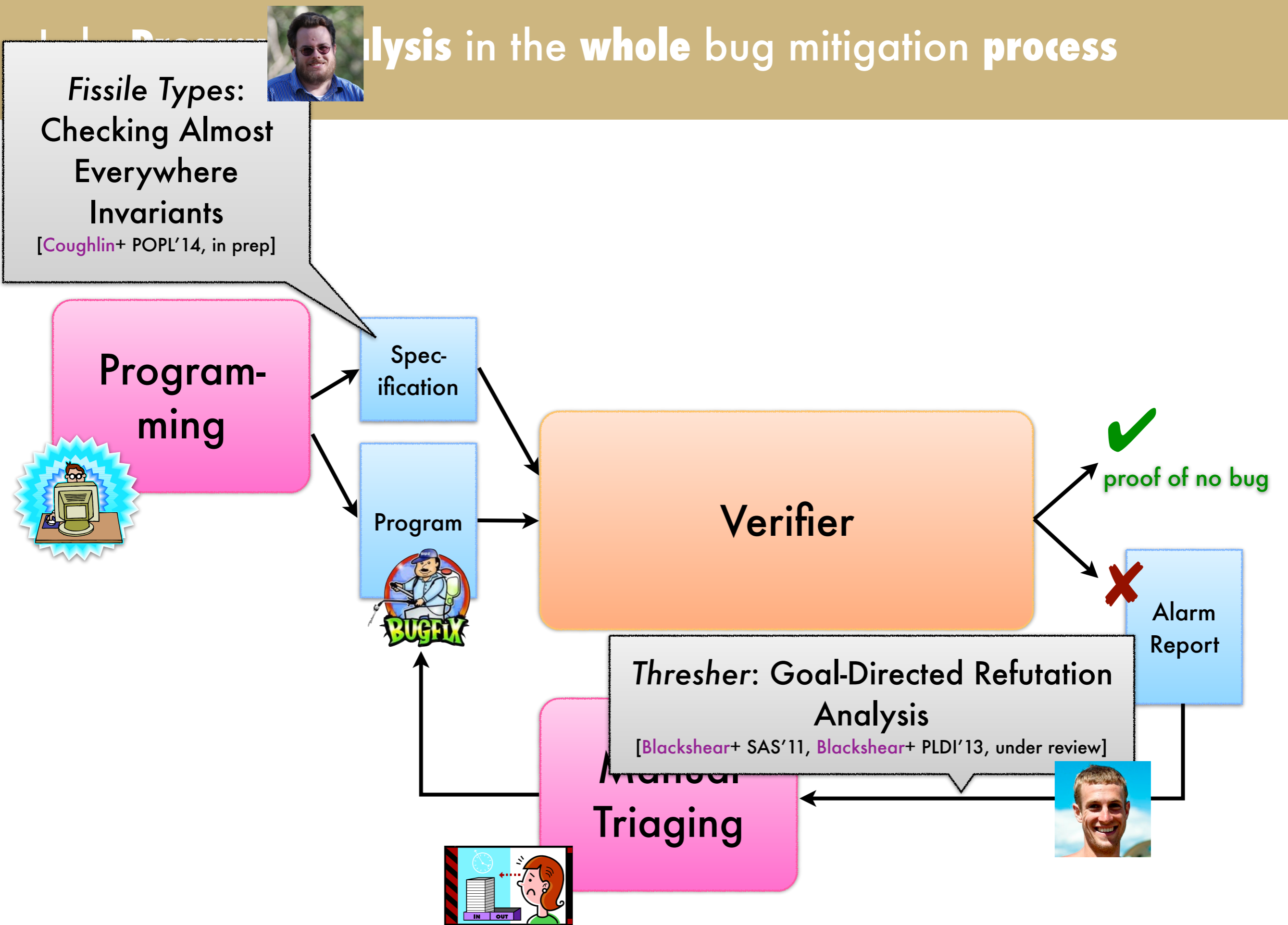
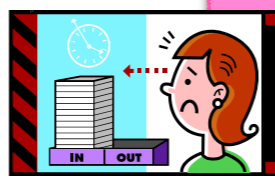
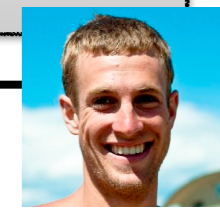
Verifier

✓
proof of no bug

✗
Alarm
Report

**Thresher: Goal-Directed Refutation
Analysis**
[Blackshear+ SAS'11, Blackshear+ PLDI'13, under review]

**Manual
Triaging**



Analysis in the whole bug mitigation process



**Fissile Types:
Checking Almost
Everywhere
Invariants**
[Coughlin+ POPL'14, in prep]

**Program-
ming**



Test
Input

Spec-
ification

Program



Runner

Test
Output

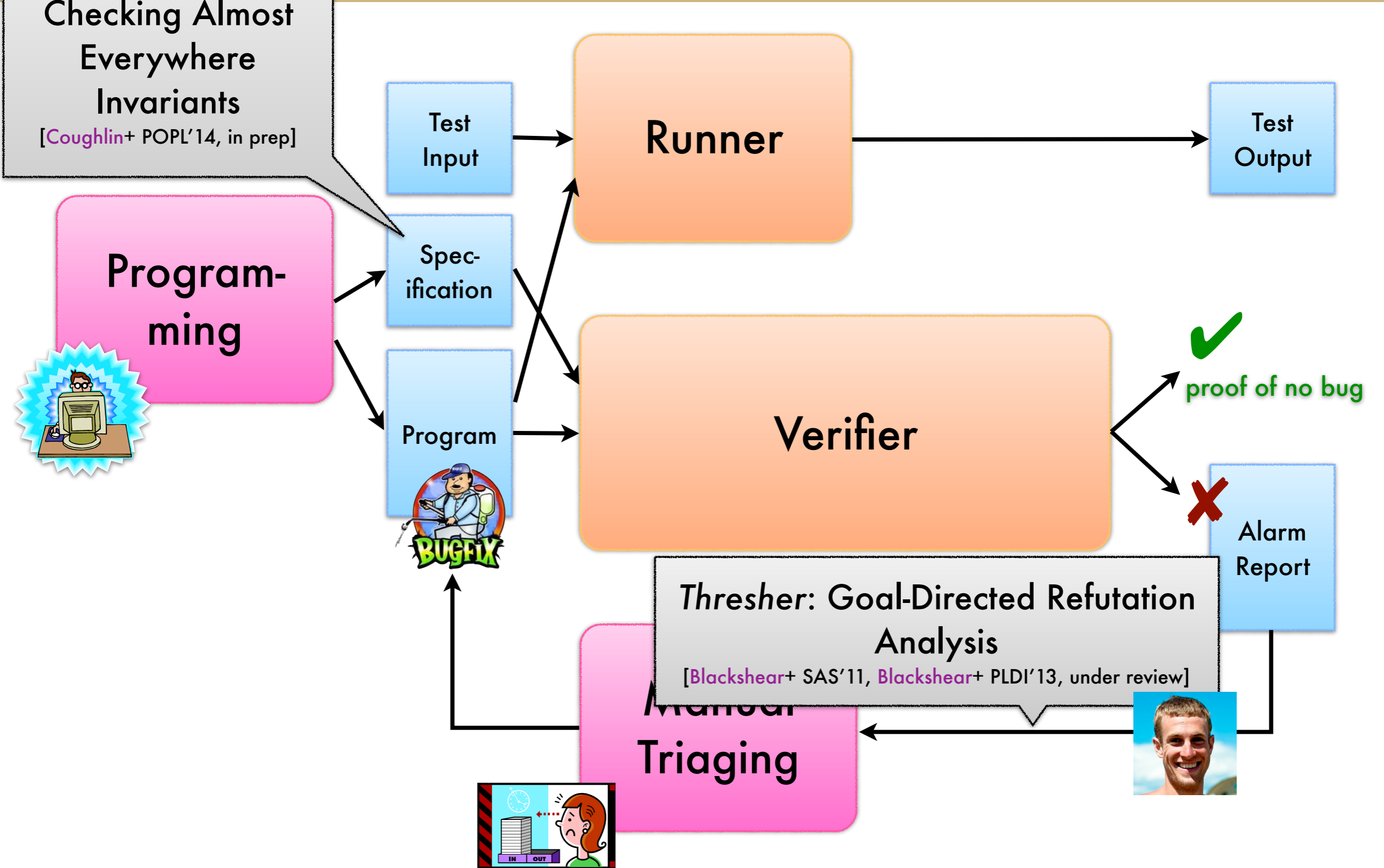
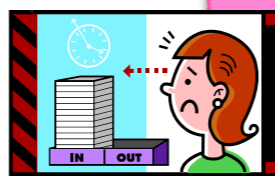
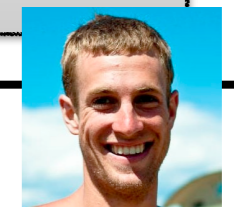
Verifier

✓
proof of no bug

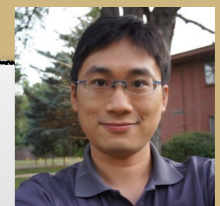
✗
Alarm
Report

**Thresher: Goal-Directed Refutation
Analysis**
[Blackshear+ SAS'11, Blackshear+ PLDI'13, under review]

**Manual
Triaging**

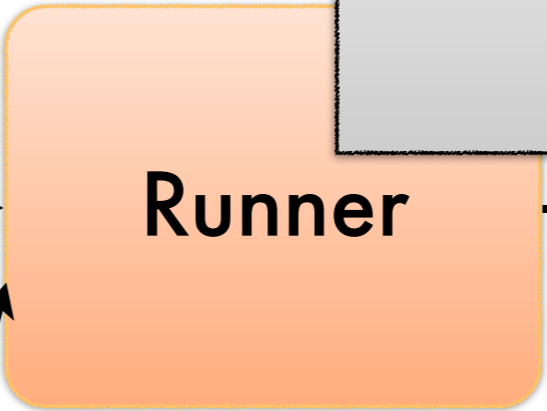


Analysis in the whole

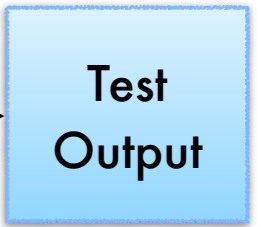
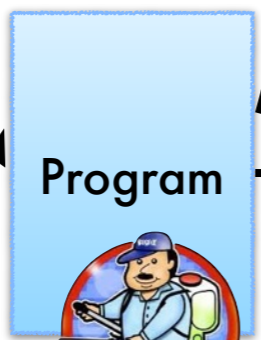
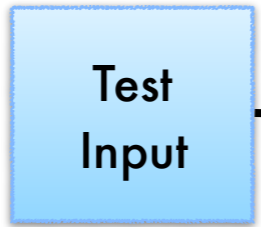
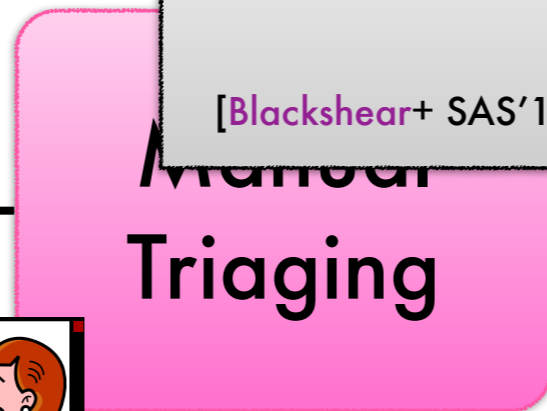


Fissile Types: Checking Almost Everywhere Invariants
[Coughlin+ POPL'14, in prep]

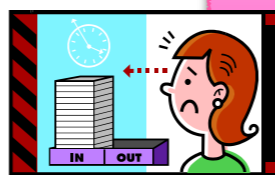
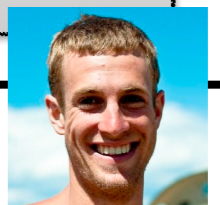
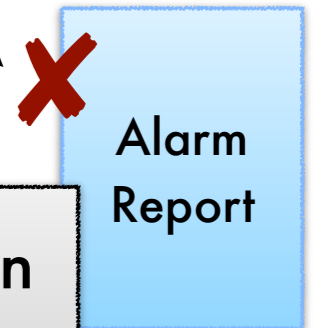
Divva: Synthesizing Short-Circuiting Data Structure Checks
[under review]



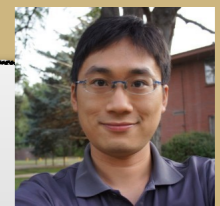
Thresher: Goal-Directed Refutation Analysis
[Blackshear+ SAS'11, Blackshear+ PLDI'13, under review]



✓
proof of no bug



Analysis in the whole



Fissile Types: Checking Almost Everywhere Invariants
[Coughlin+ POPL'14, in prep]

Divva: Synthesizing Short-Circuiting Data Structure Checks
[under review]

**Program-
ming**



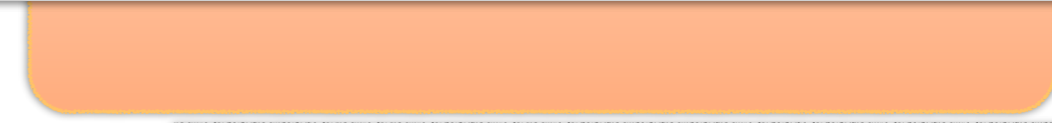
Test Input

Spec-
ification

Program



Use **static shape analysis** to **synthesize short-circuiting dynamic validation** of data structure invariants
uninterpreted inductive separation logic predicates

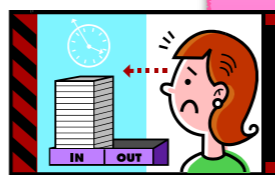
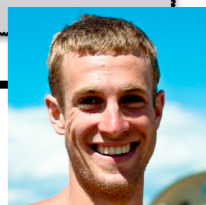


Test

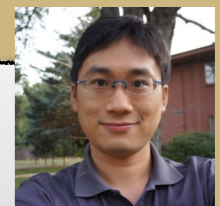
Alarm Report

Thresher: Goal-Directed Refutation Analysis
[Blackshear+ SAS'11, Blackshear+ PLDI'13, under review]

**Manual
Triaging**

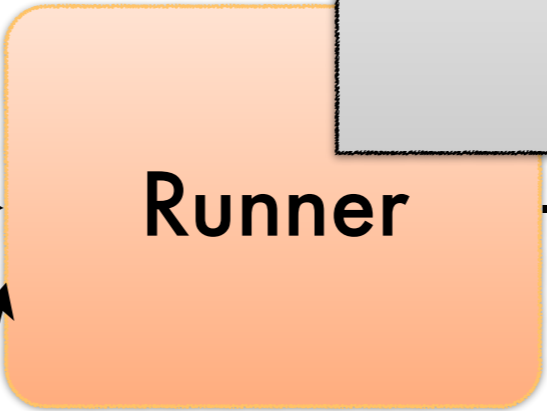


Analysis in the whole

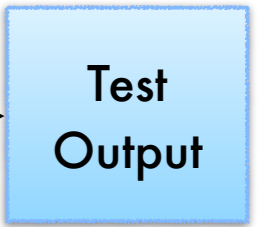
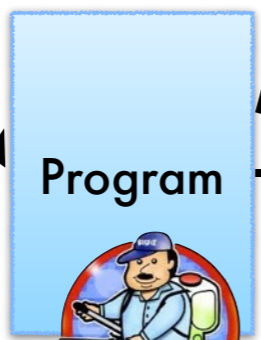
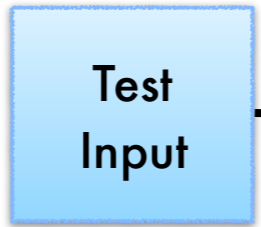
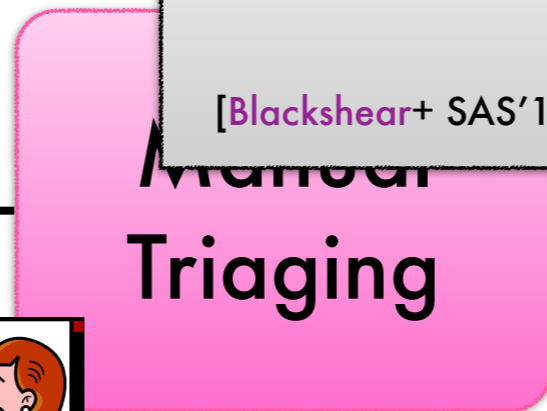


Fissile Types: Checking Almost Everywhere Invariants
[Coughlin+ POPL'14, in prep]

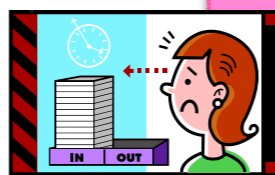
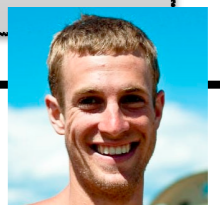
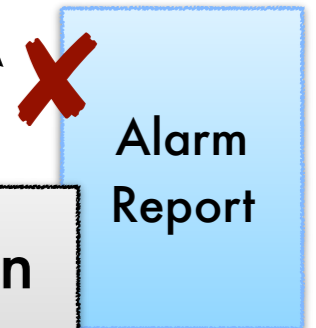
Divva: Synthesizing Short-Circuiting Data Structure Checks
[under review]



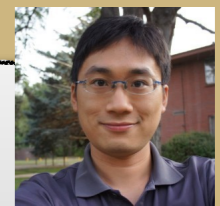
Thresher: Goal-Directed Refutation Analysis
[Blackshear+ SAS'11, Blackshear+ PLDI'13, under review]



✓
proof of no bug

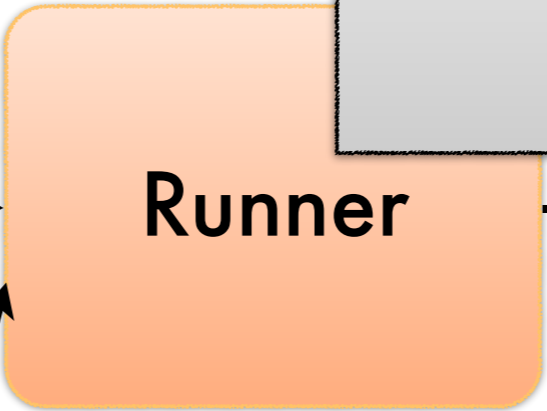


Analysis in the whole

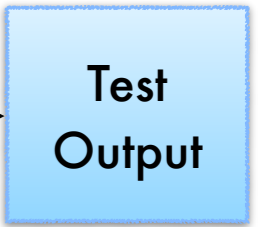
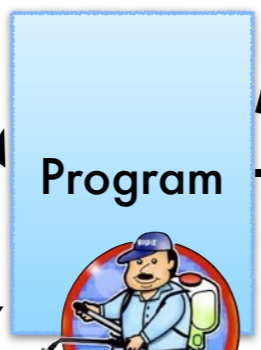
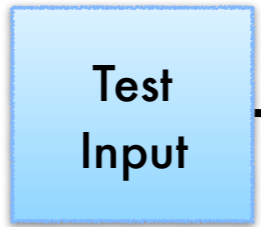
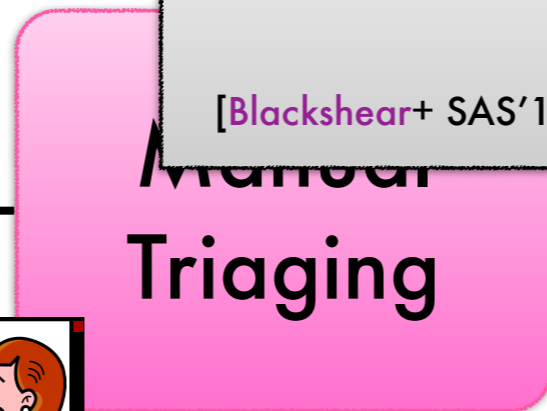


Fissile Types: Checking Almost Everywhere Invariants
[Coughlin+ POPL'14, in prep]

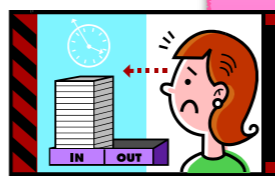
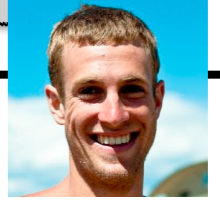
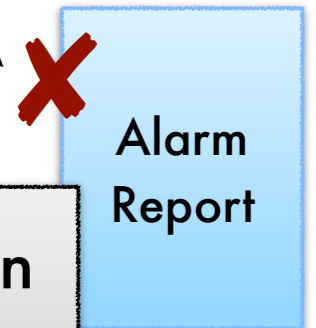
Divva: Synthesizing Short-Circuiting Data Structure Checks
[under review]



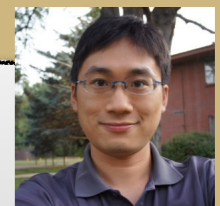
Thresher: Goal-Directed Refutation Analysis
[Blackshear+ SAS'11, Blackshear+ PLDI'13, under review]



✓
proof of no bug



Analysis in the whole



Fissile Types: Checking Almost Everywhere Invariants
[Coughlin+ POPL'14, in prep]

Divva: Synthesizing Short-Circuiting Data Structure Checks
[under review]

Test Input

Runner

Test Output

Specification

Verifier

✓
proof of no bug

Program

✗
Alarm Report

Programming



Github

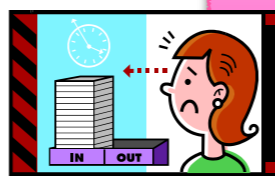
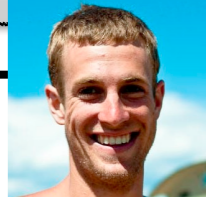


Thresher: Goal-Directed Refutation Analysis
[Blackshear+ SAS'11, Blackshear+ PLDI'13, under review]

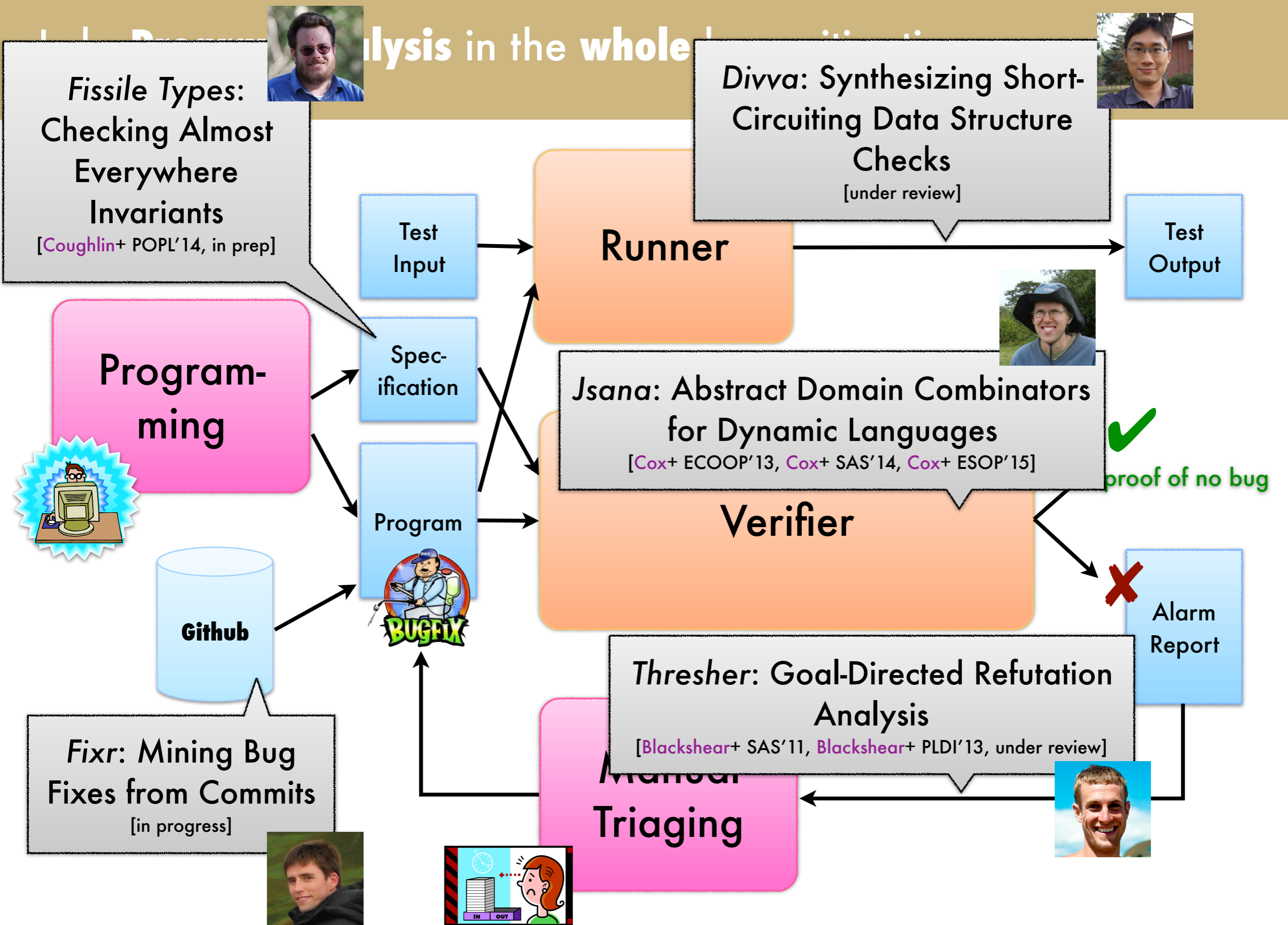
Fixr: Mining Bug Fixes from Commits
[in progress]



Manual Triaging



Analysis in the whole



Fissile Types: Checking Almost Everywhere Invariants
[Coughlin+ POPL'14, in prep]

Divva: Synthesizing Short-Circuiting Data Structure Checks
[under review]

**Program-
ming**

Test Input

Runner

Test Output

Spec-
ification

Jsana: Abstract Domain Combinators for Dynamic Languages
[Cox+ ECOOP'13, Cox+ SAS'14, Cox+ ESOP'15]

proof of no bug ✓

Program

Verifier

Alarm Report ✗

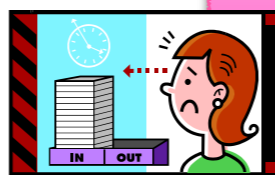
Github



Thresher: Goal-Directed Refutation Analysis
[Blackshear+ SAS'11, Blackshear+ PLDI'13, under review]

Fixr: Mining Bug Fixes from Commits
[in progress]

**Manual
Triaging**



Analysis in the whole



Fissile Types: Checking Almost Everywhere Invariants
[Coughlin+ POPL'14, in prep]

Divva: Synthesizing Short-Circuiting Data Structure Checks
[under review]

Test Input

Runner

Test Output

Specification

Jsana: Abstract Domain Combinators for Dynamic Languages
[Cox+ ECOOP'13, Cox+ SAS'14, Cox+ ESOP'15]



proof of no bug ✓

Program-
ming

Program



Github



Fixr: Mining Bug Fixes from Commits
[in progress]



X
↓

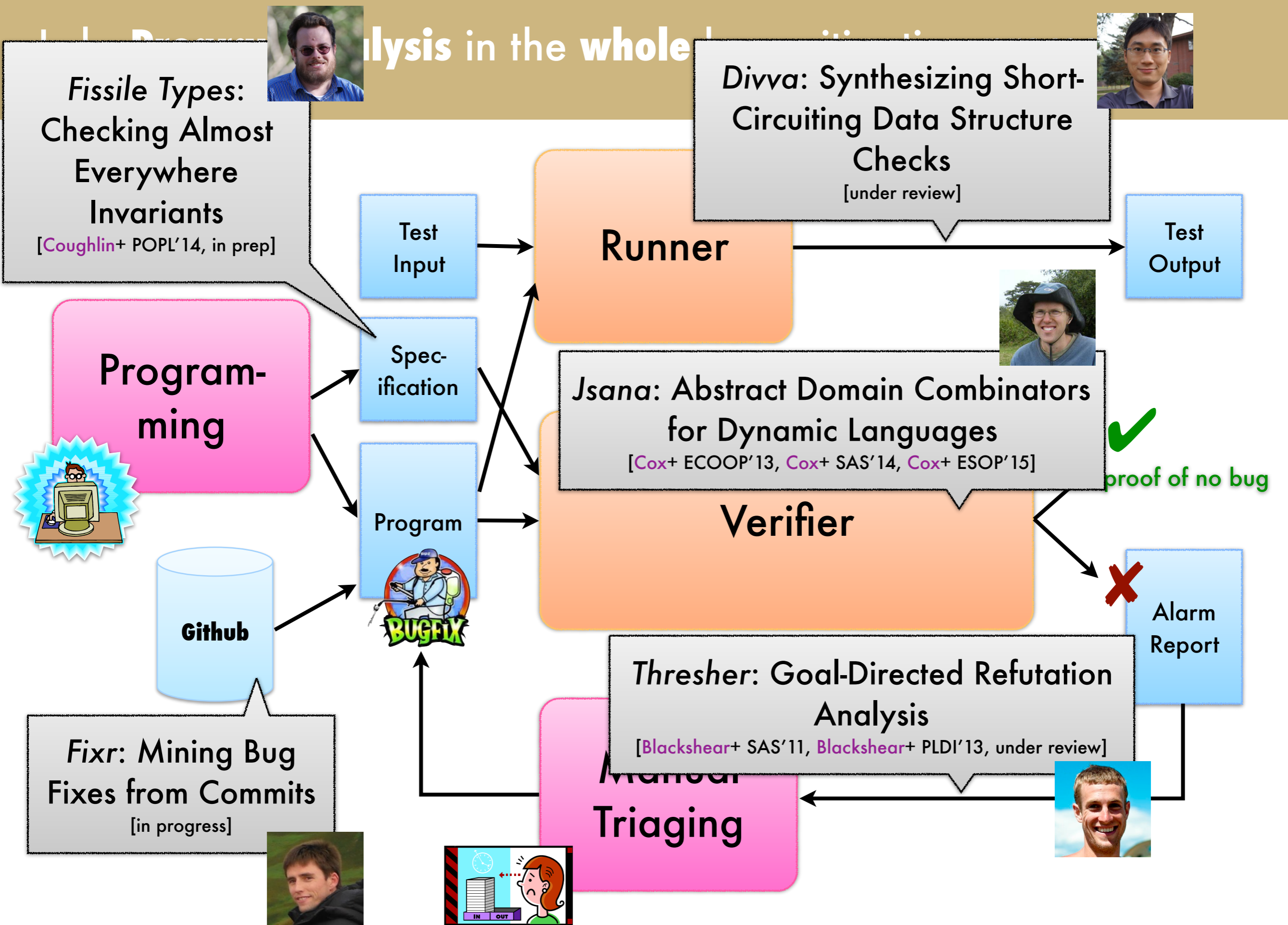
A ₁	
F ₁	V ₁
F ₂	V ₂
F ₃	A ₂

Heap with **set symbols** partitioning "open objects"

separation logic with open object predicates and desynchronized separation

arm
port

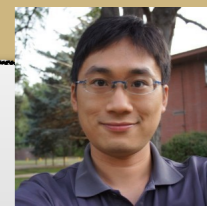
Analysis in the whole



Fissile Types: Checking Almost Everywhere Invariants
[Coughlin+ POPL'14, in prep]



Divva: Synthesizing Short-Circuiting Data Structure Checks
[under review]



**Program-
ming**



Test Input

Spec-
ification

Program

Runner

Jsana: Abstract Domain Combinators for Dynamic Languages
[Cox+ ECOOP'13, Cox+ SAS'14, Cox+ ESOP'15]



Verifier

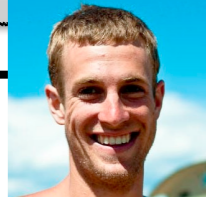
proof of no bug ✓

Github



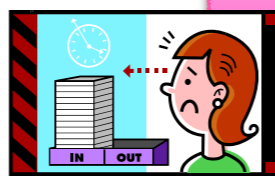
Alarm Report

Thresher: Goal-Directed Refutation Analysis
[Blackshear+ SAS'11, Blackshear+ PLDI'13, under review]



**Manual
Triaging**

Fixr: Mining Bug Fixes from Commits
[in progress]



Lab: Program analysis in the whole bug mitigation process

Fissile Types:

Checking Almost Everywhere Invariants
[Coughlin+ POPL'14, in prep]

Divva: Synthesizing Short-Circuiting Data Structure Checks
[under review]

Program-
ming



This Talk

Fixr: Mining Bug Fixes from Commits
[in progress]

Runner

Jsana: Abstract Domain Combinators for Dynamic Languages
[Cox+ ECOOP'13, Cox+ SAS'14, Cox+ ESOP'15]

Verifier

Thresher: Goal-Directed Refutation Analysis
[Blackshear+ SAS'11, Blackshear+ PLDI'13, under review]

Triaging

Test Input

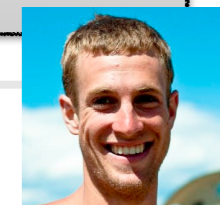
Spec-
ification

Program

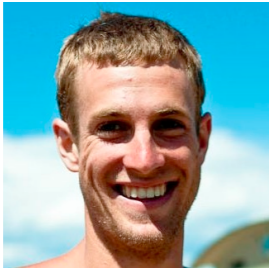
Test Output

✓
proof of no bug

✗
Alarm Report



Goal-Directed Program Analysis with Jumping



Sam Blackshear
University of Colorado Boulder



Bor-Yuh Evan Chang
University of Colorado Boulder



Manu Sridharan
Samsung Research America

Google
July 24, 2015





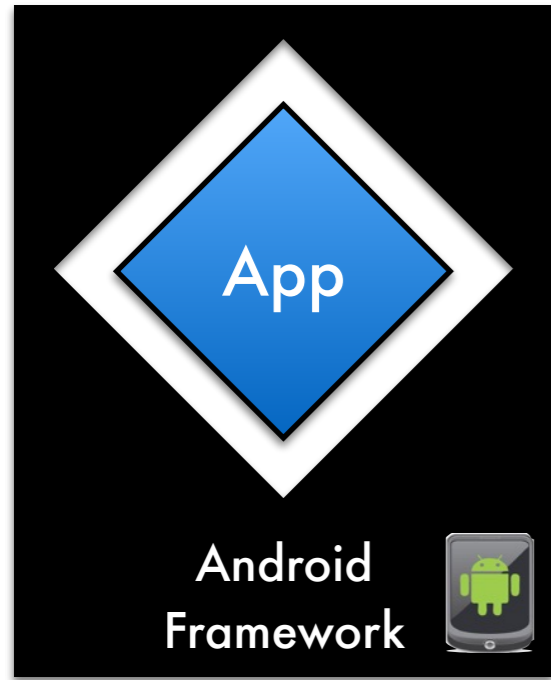


**Roughly, 3% of all commmits fix
NullPointerExceptions**

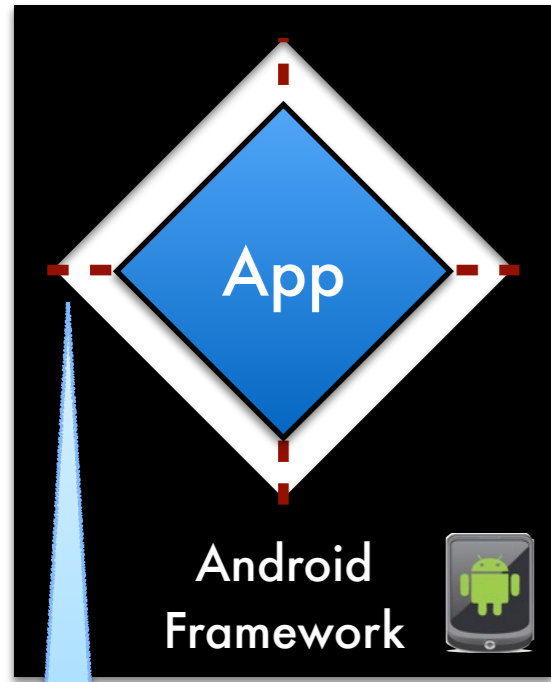
Callback-oriented programming



Callback-oriented programming

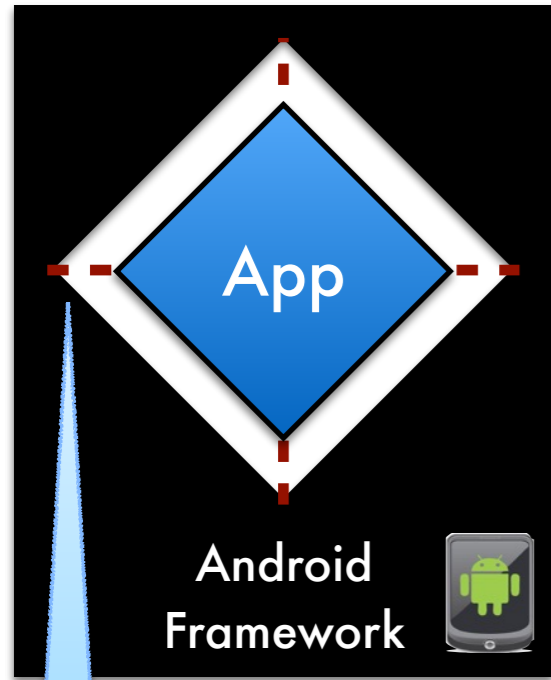


Callback-oriented programming

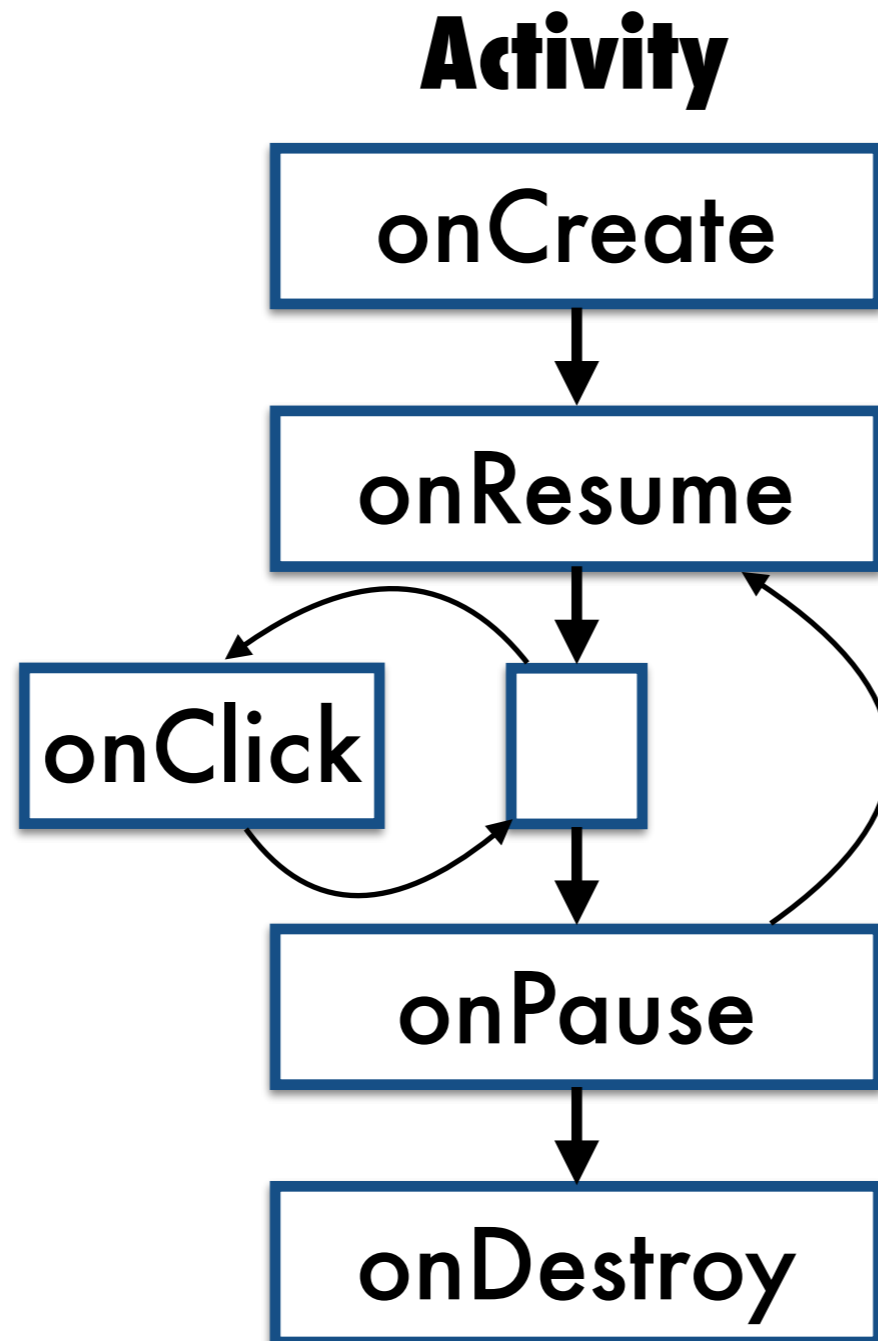


callbacks (e.g.,
Activity.onCreate)

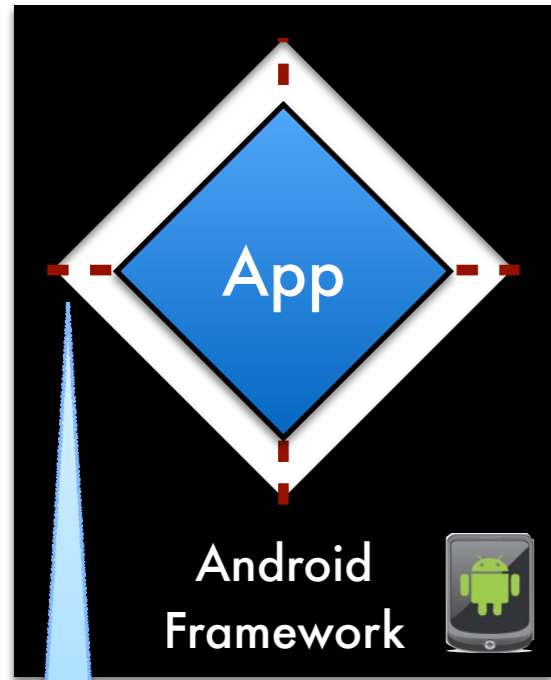
Callback-oriented programming



callbacks (e.g.,
Activity.onCreate)

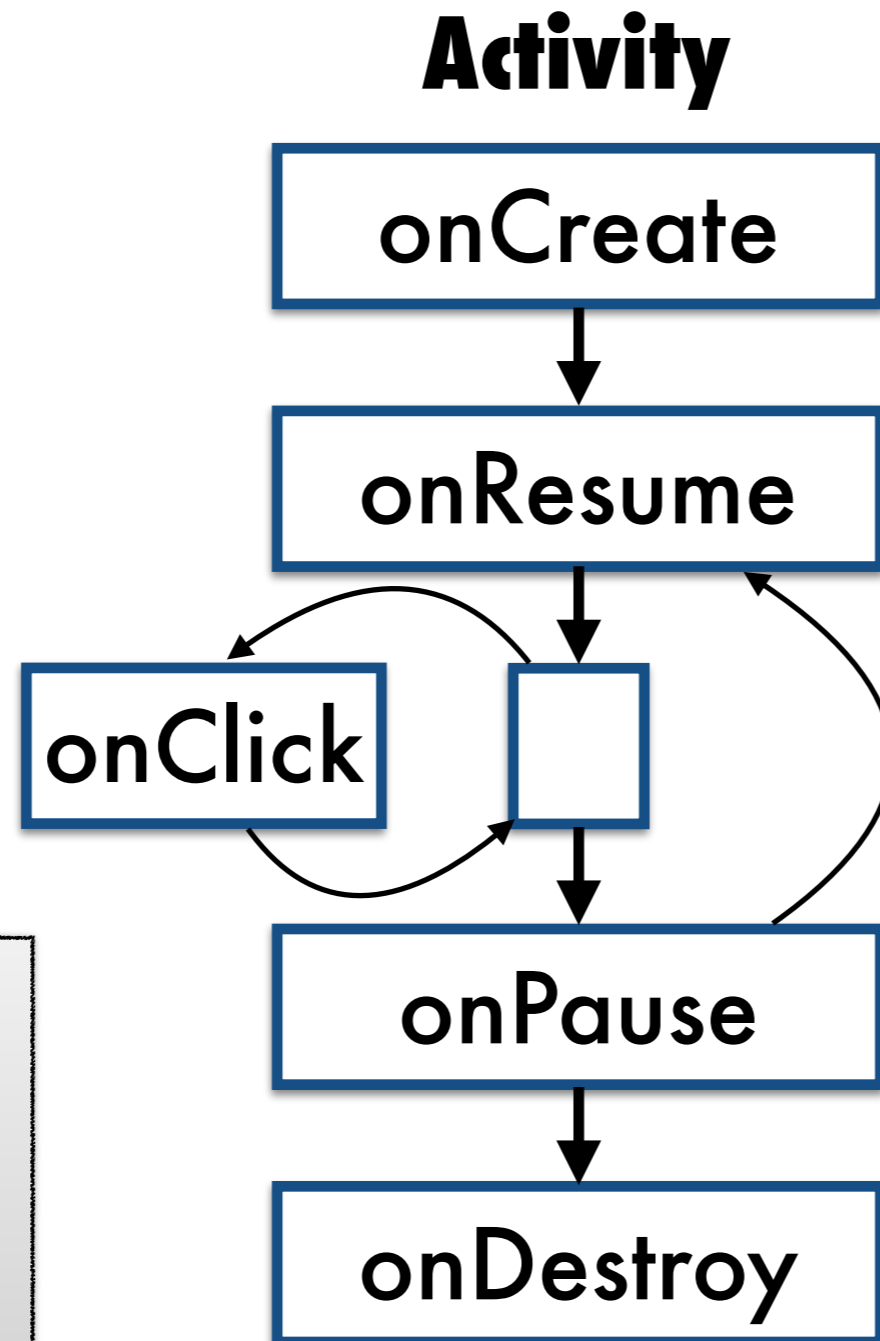


Callback-oriented programming

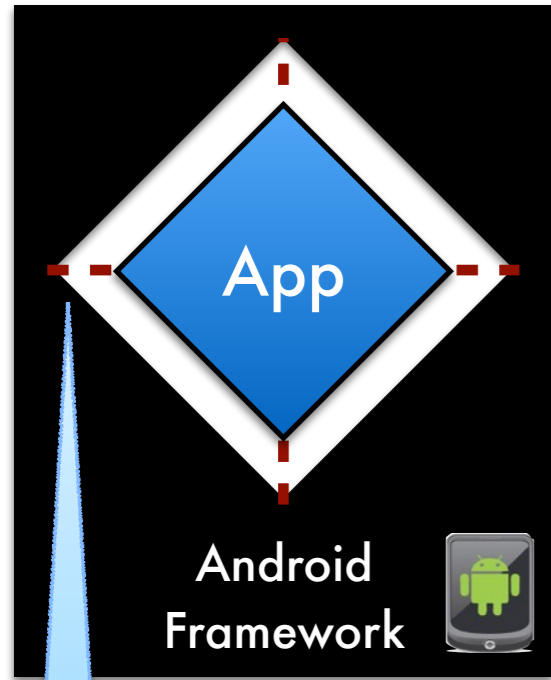


callbacks (e.g.,
Activity.onCreate)

Android
components
have an
ordered, event-
driven lifecycle

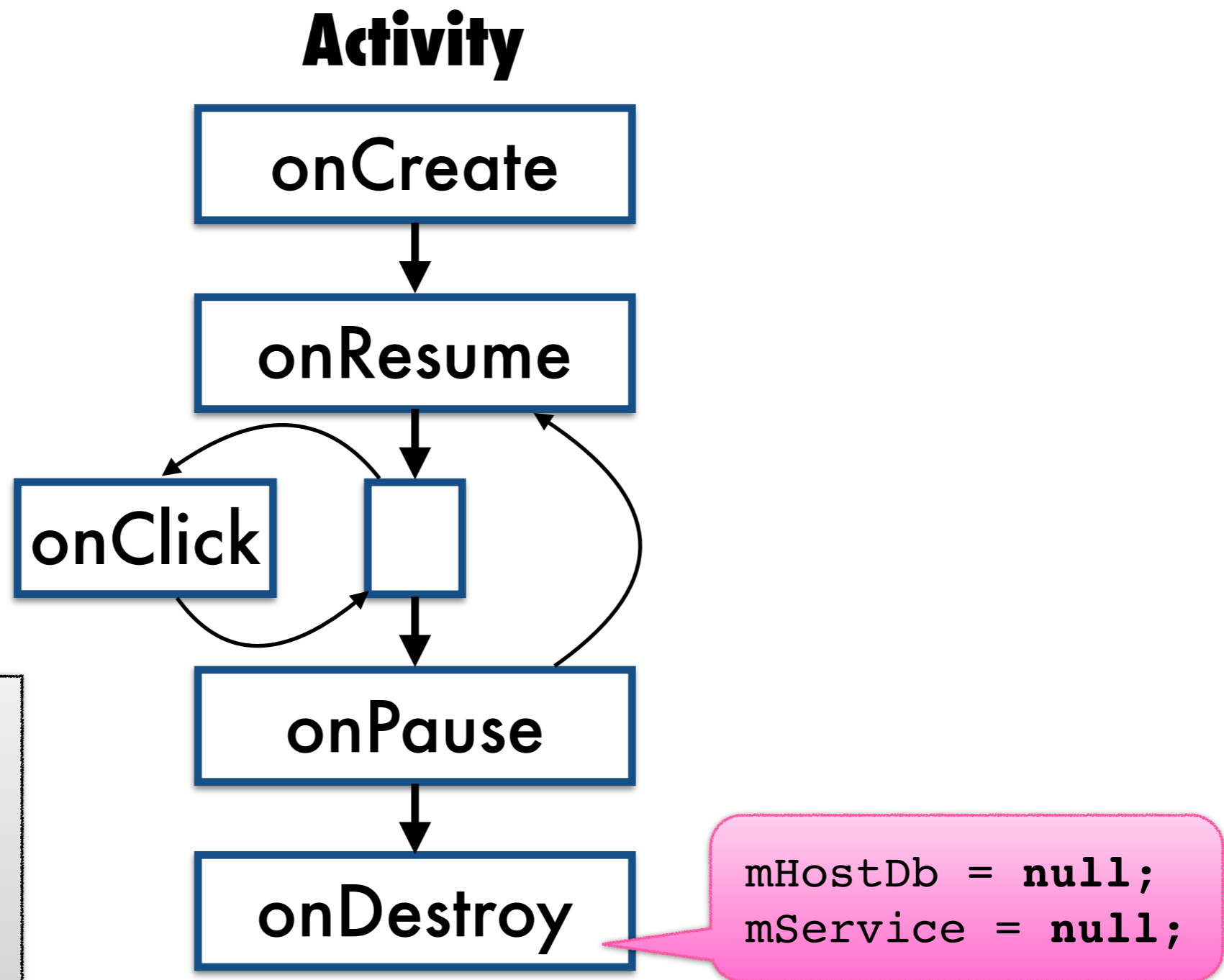


Callback-oriented programming

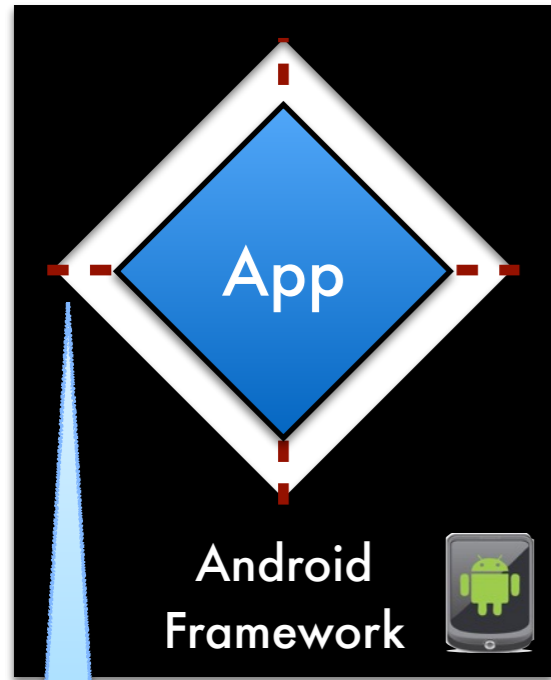


callbacks (e.g.,
Activity.onCreate)

Android
components
have an
ordered, event-
driven lifecycle

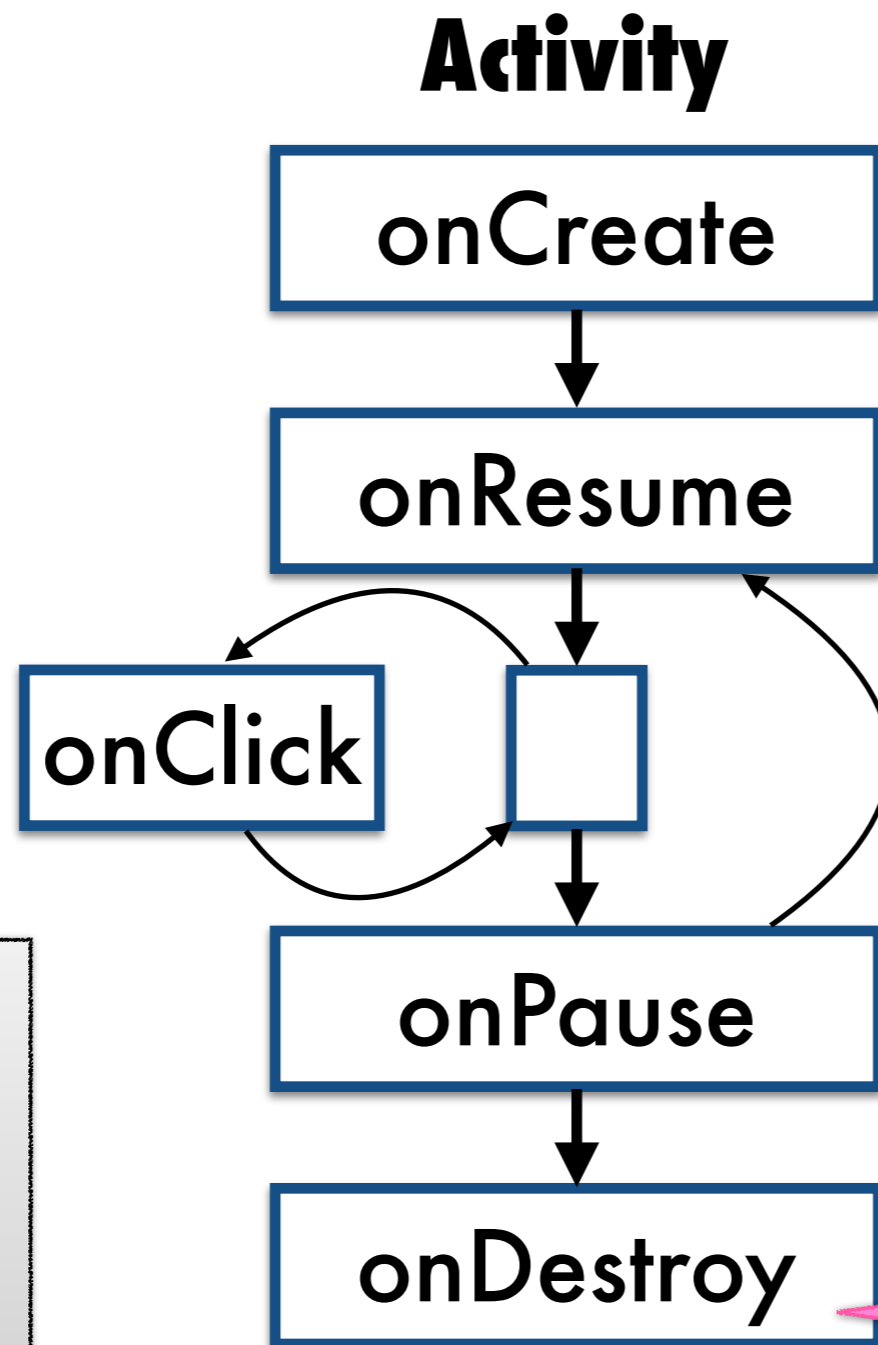


Callback-oriented programming



callbacks (e.g.,
Activity.onCreate)

Android
components
have an
ordered, event-
driven lifecycle



But, lifecycles of
different components
and other callbacks
can interleave ...

```
mHostDb = null;  
mService = null;
```

Callback-oriented programming

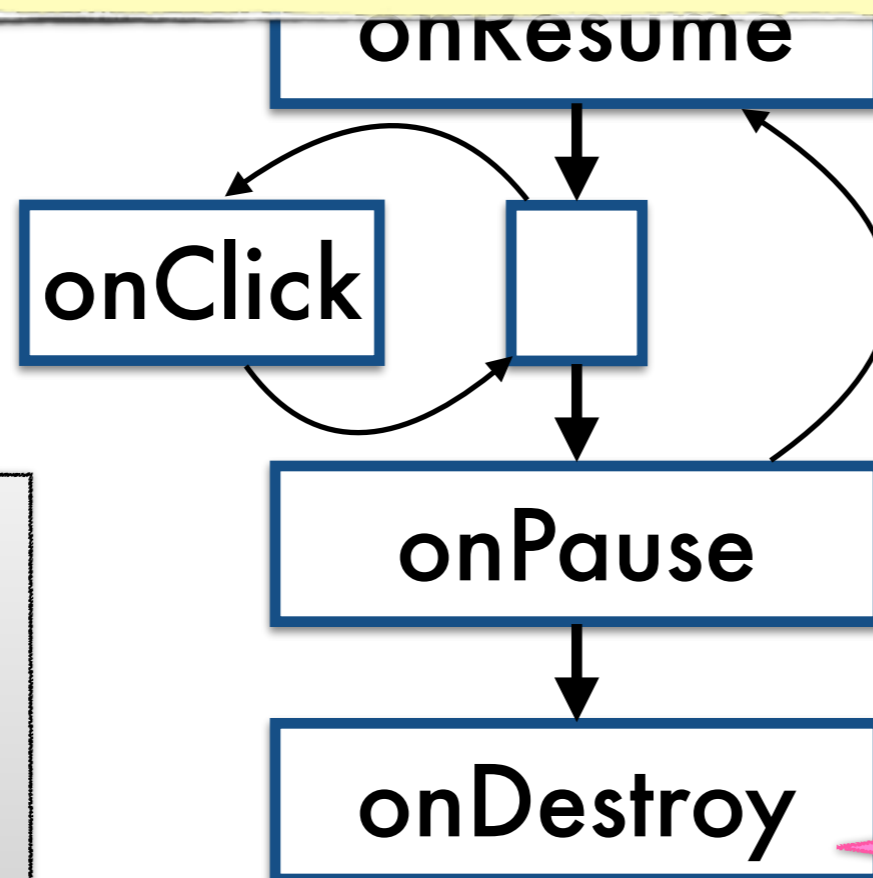
Challenge: Verifying safety (of dereferences)
depends on **callback interleaving**

Android
Framework



callbacks (e.g.,
Activity.onCreate)

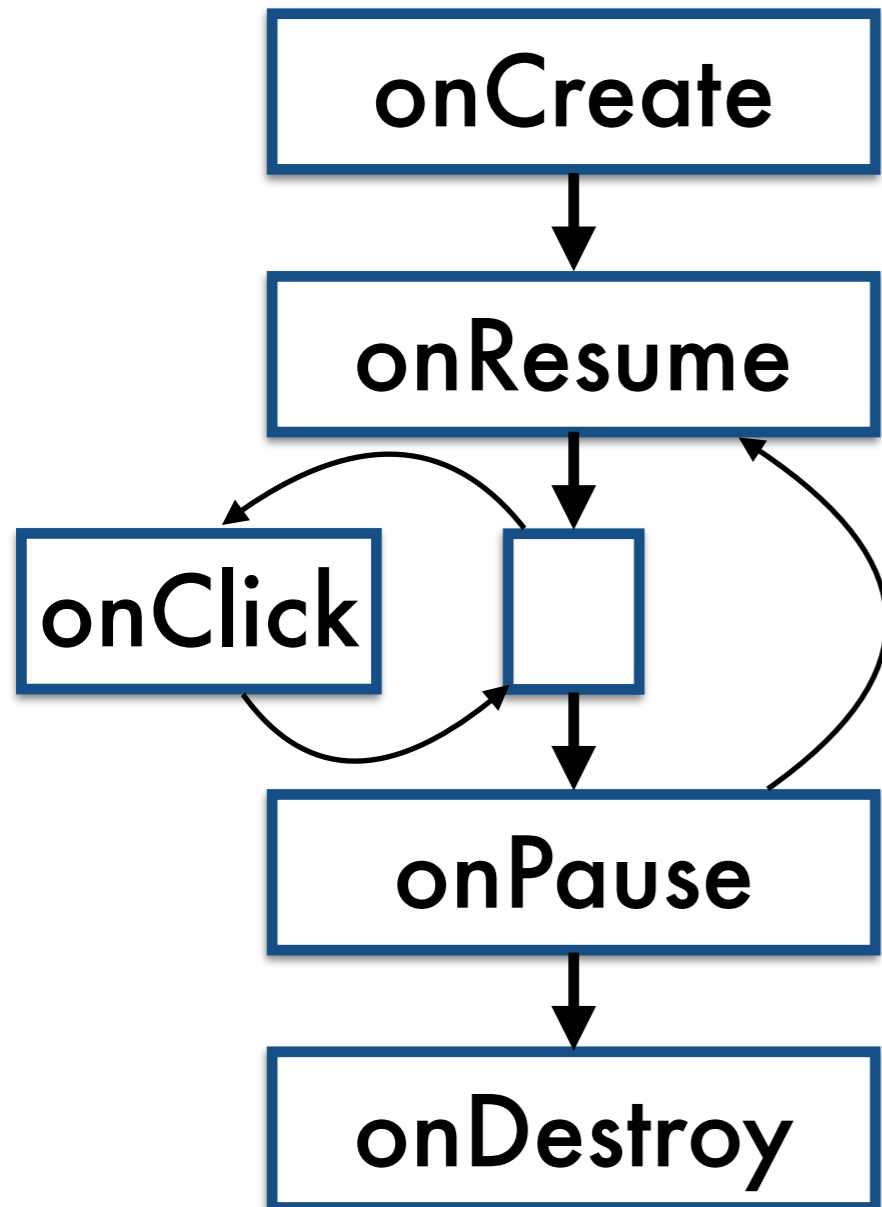
Android
components
have an
ordered, event-
driven lifecycle



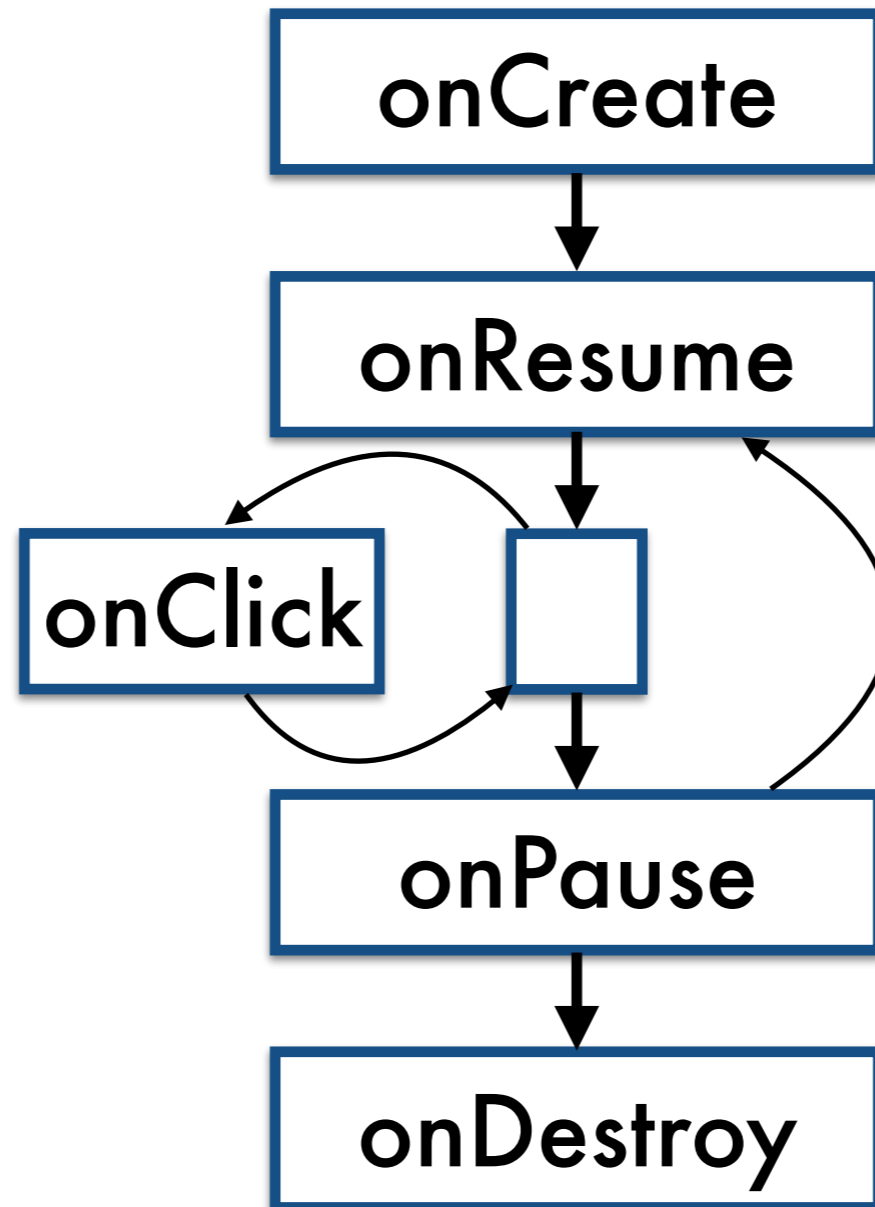
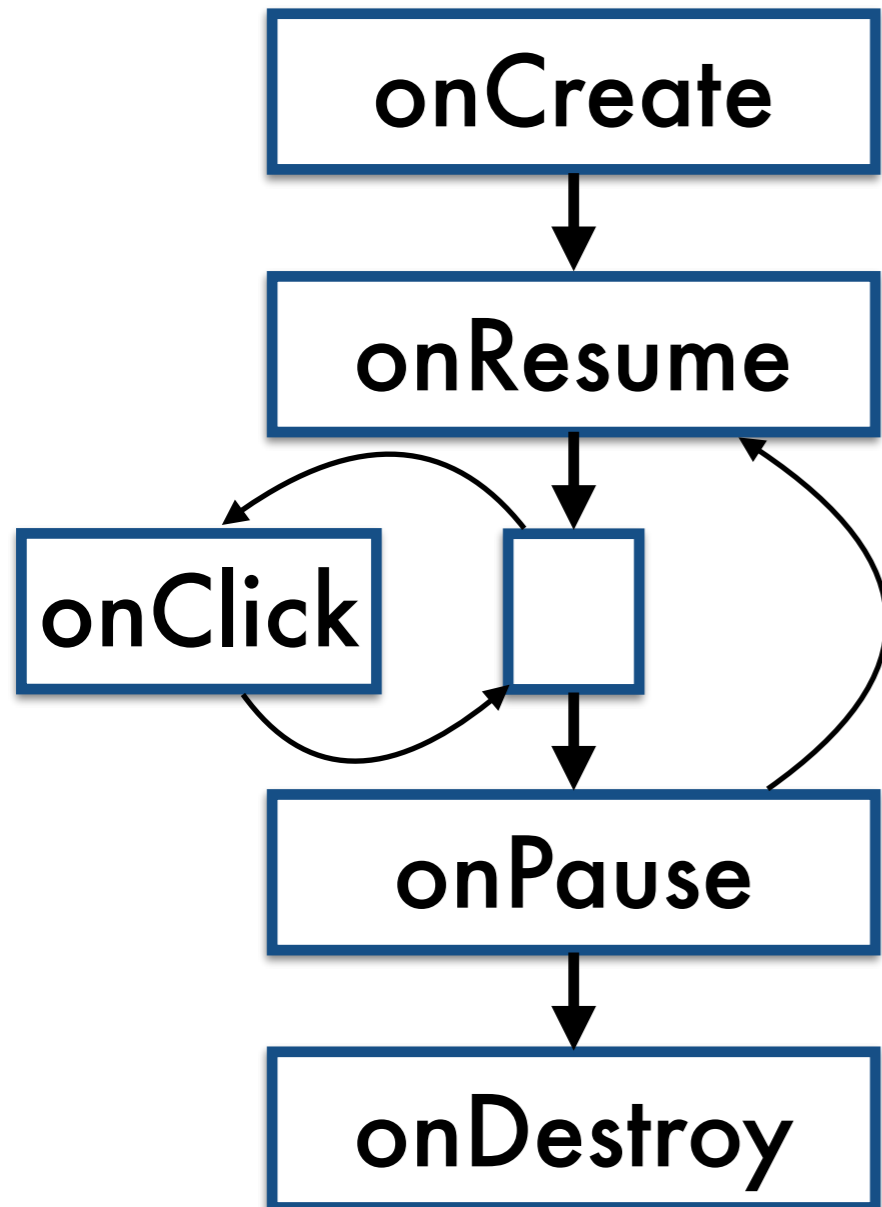
But, lifecycles of
different components
and other callbacks
can interleave ...

```
mHostDb = null;  
mService = null;
```


Callback-oriented programming

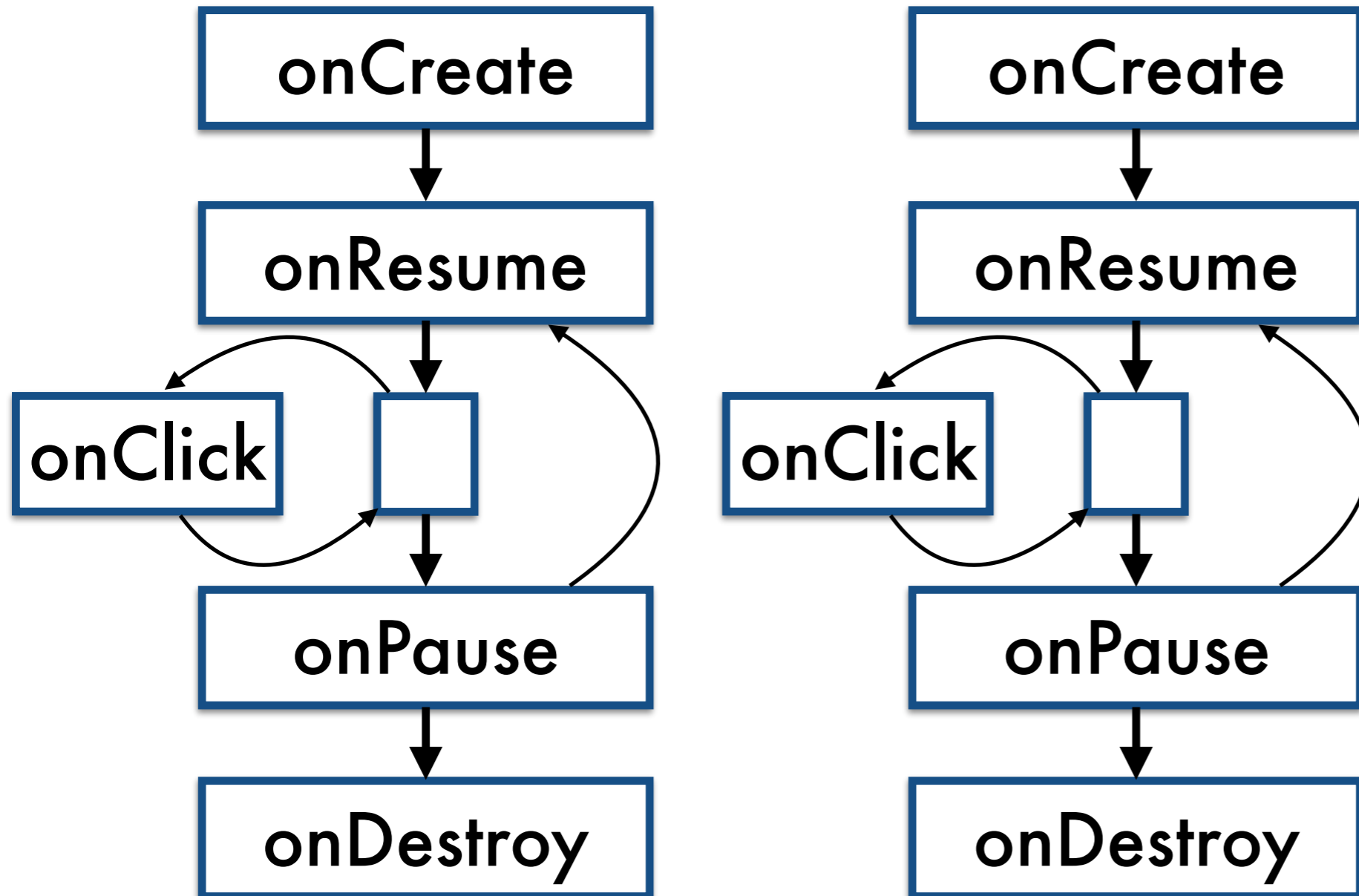


Callback-oriented programming

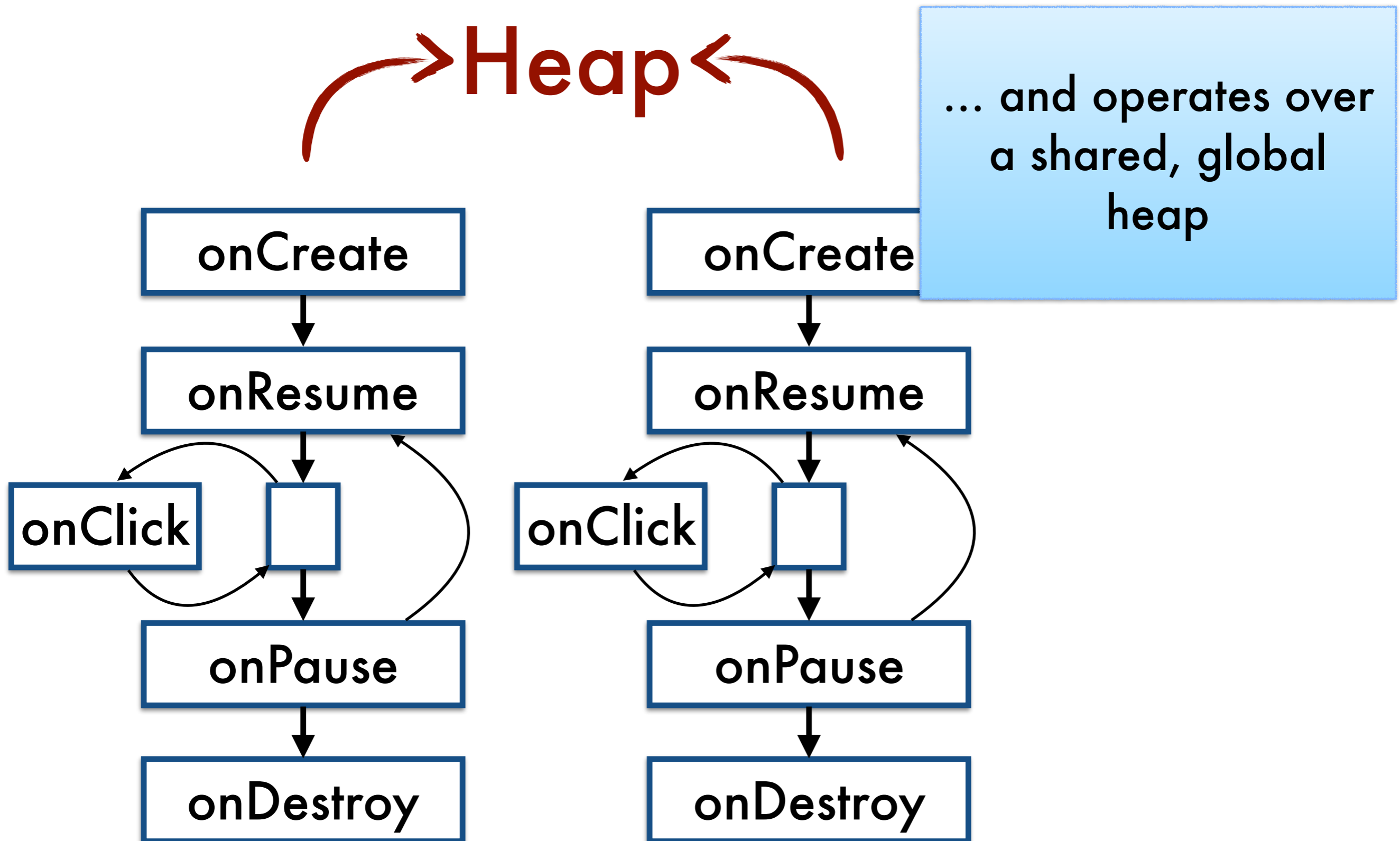


Callback-oriented programming

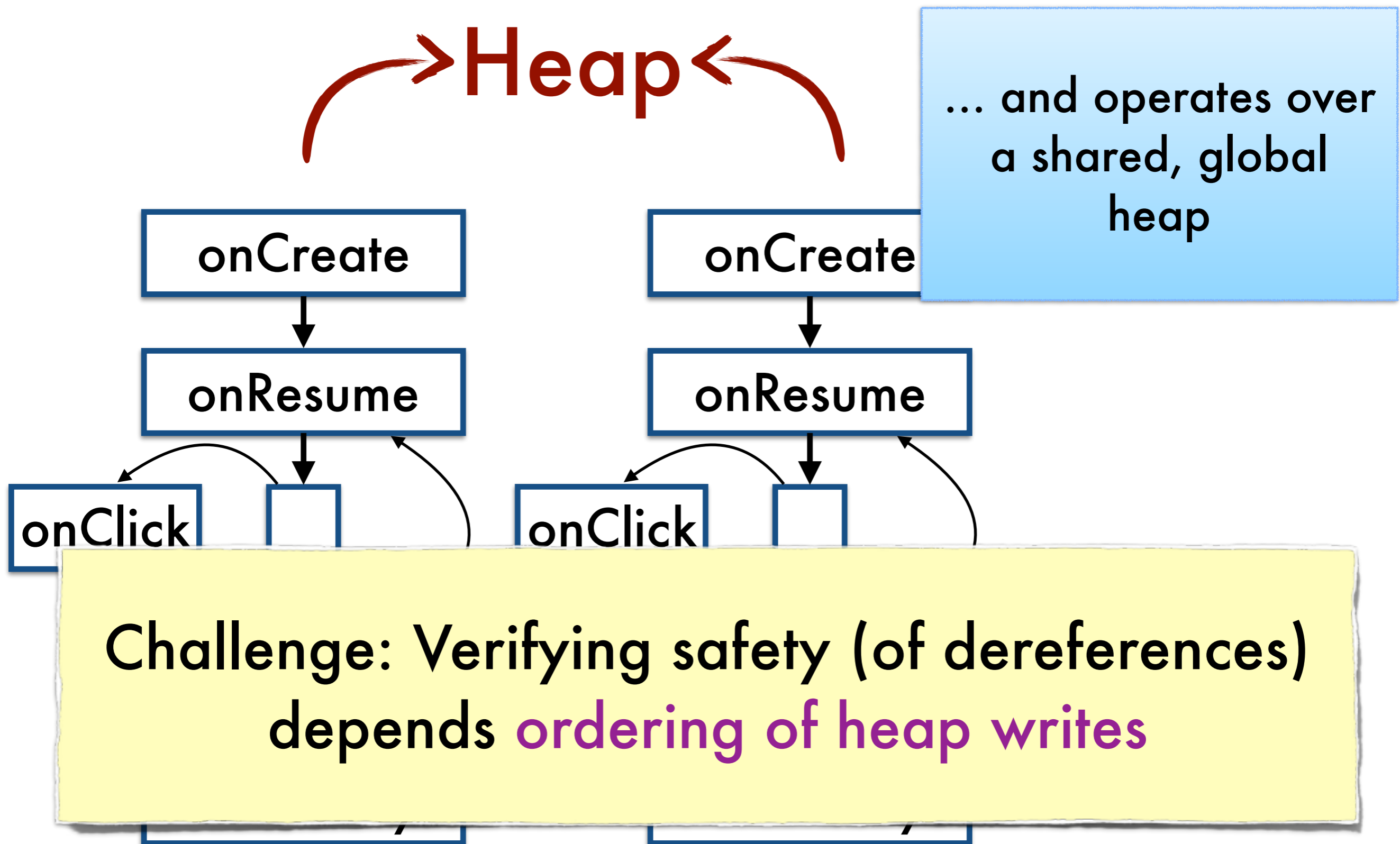
→ **Heap** ←



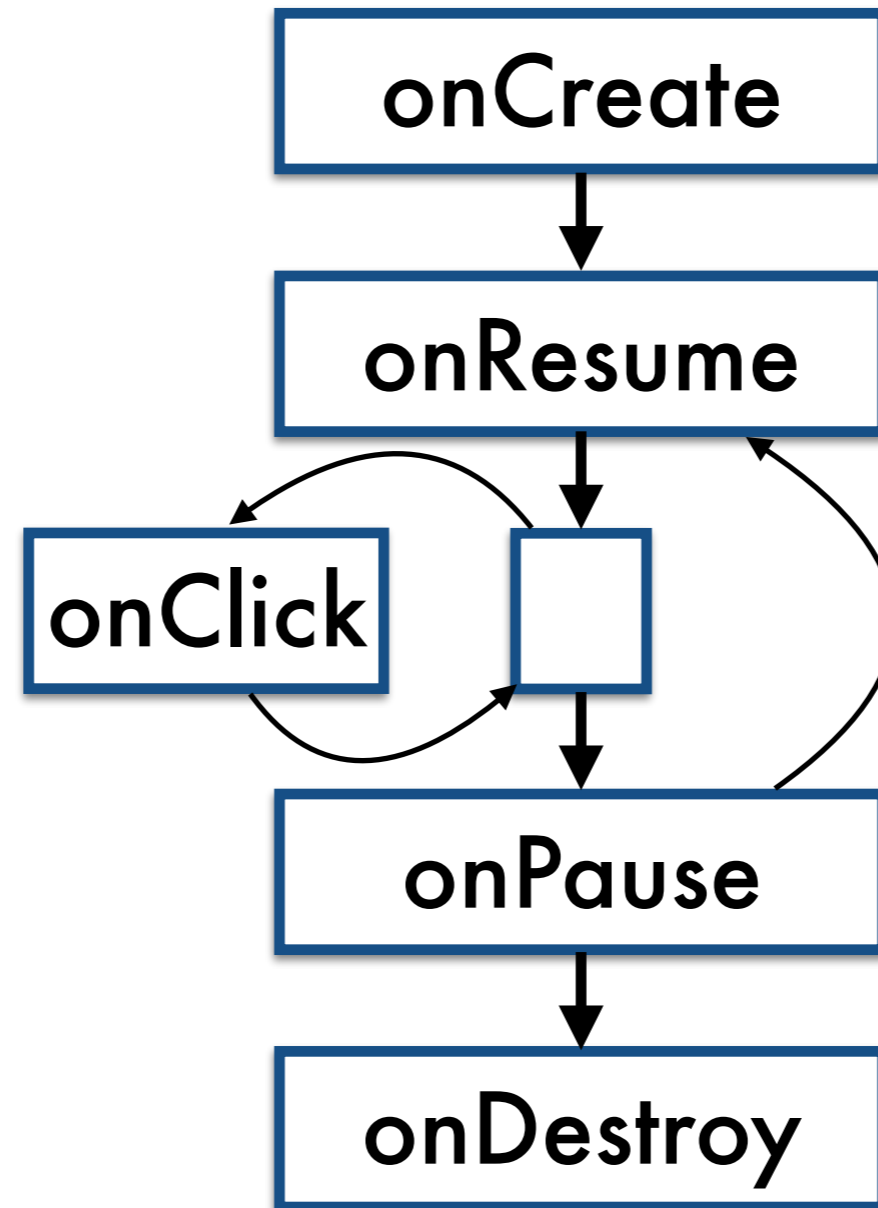
Callback-oriented programming



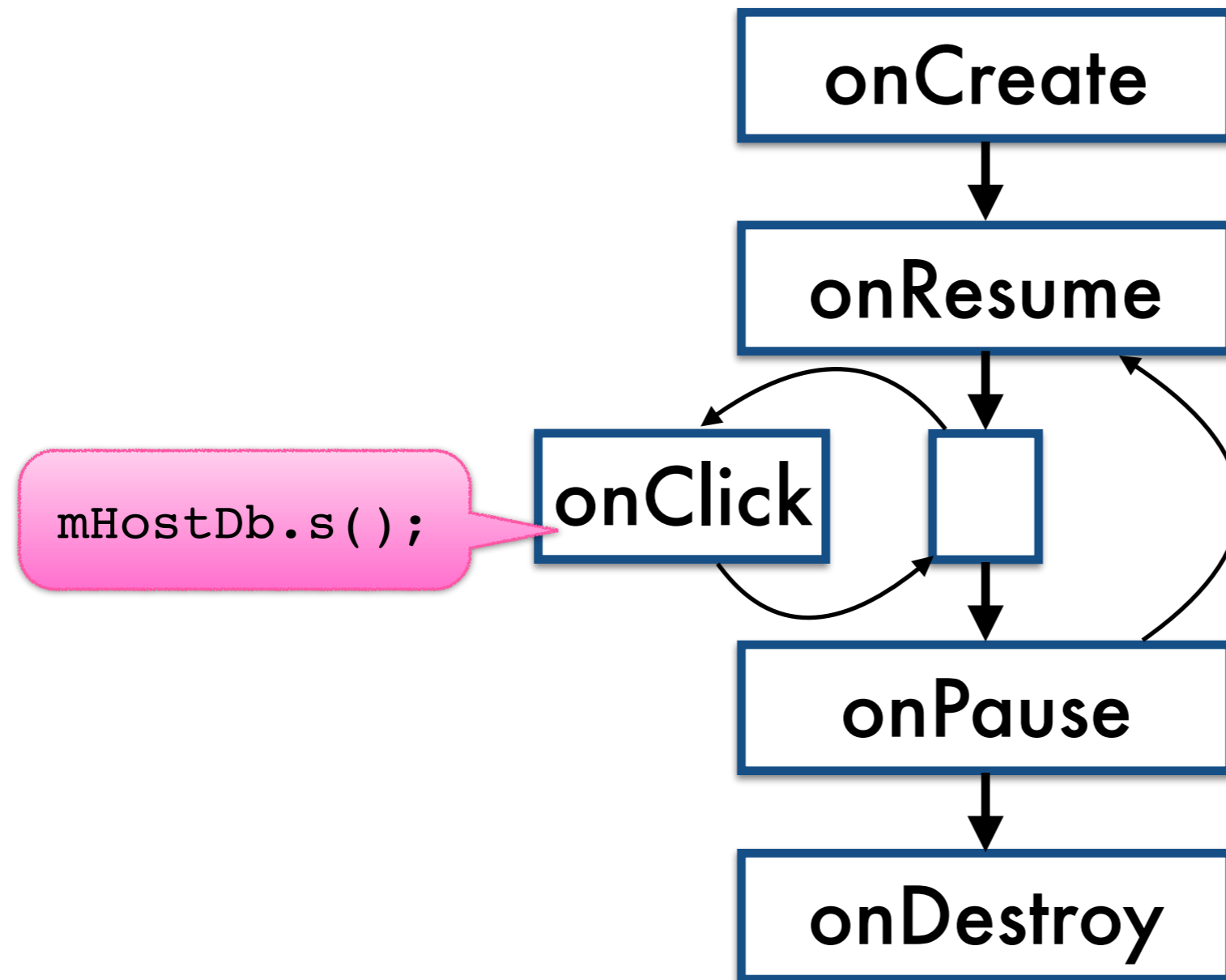
Callback-oriented programming



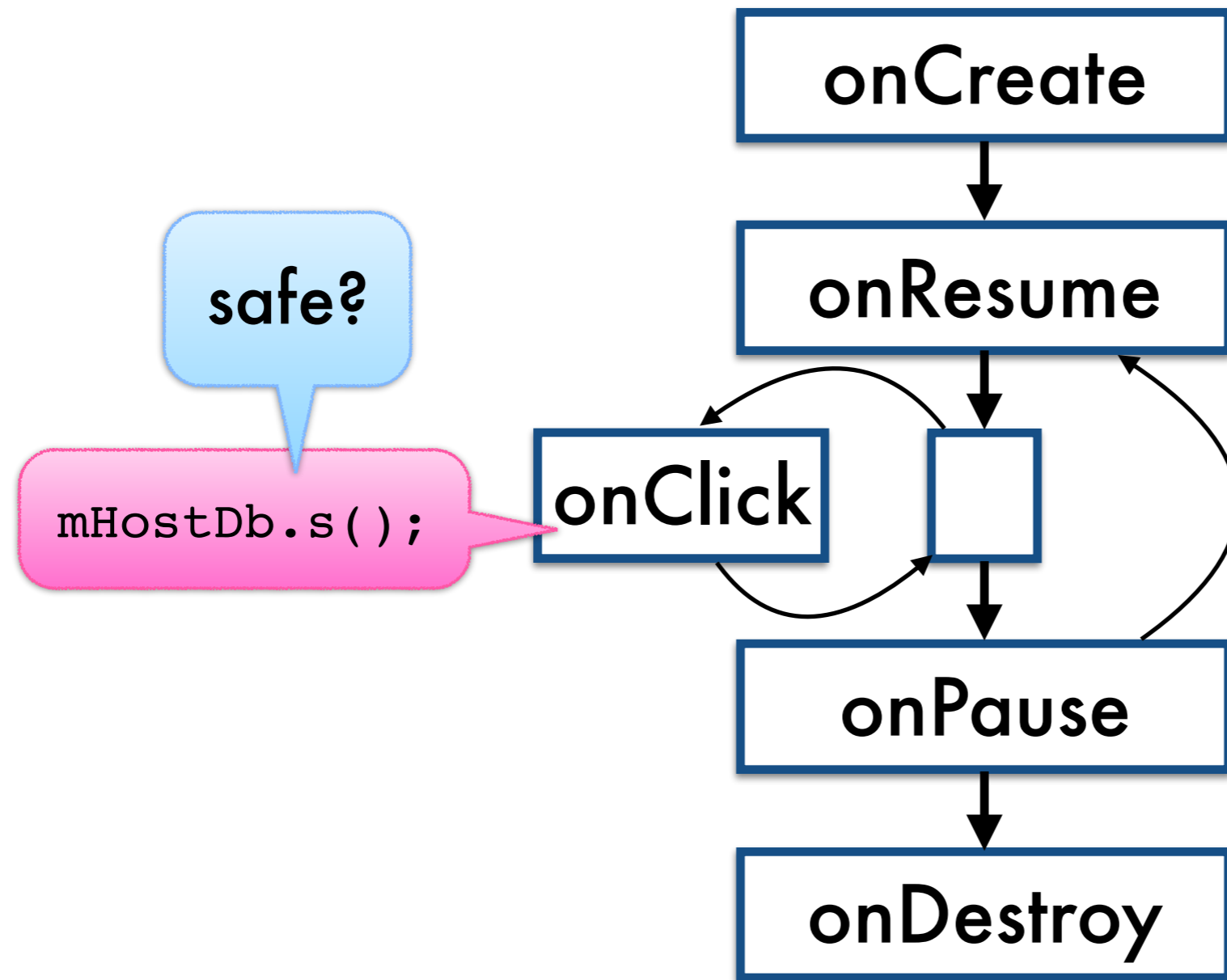
But it shouldn't be so hard ...



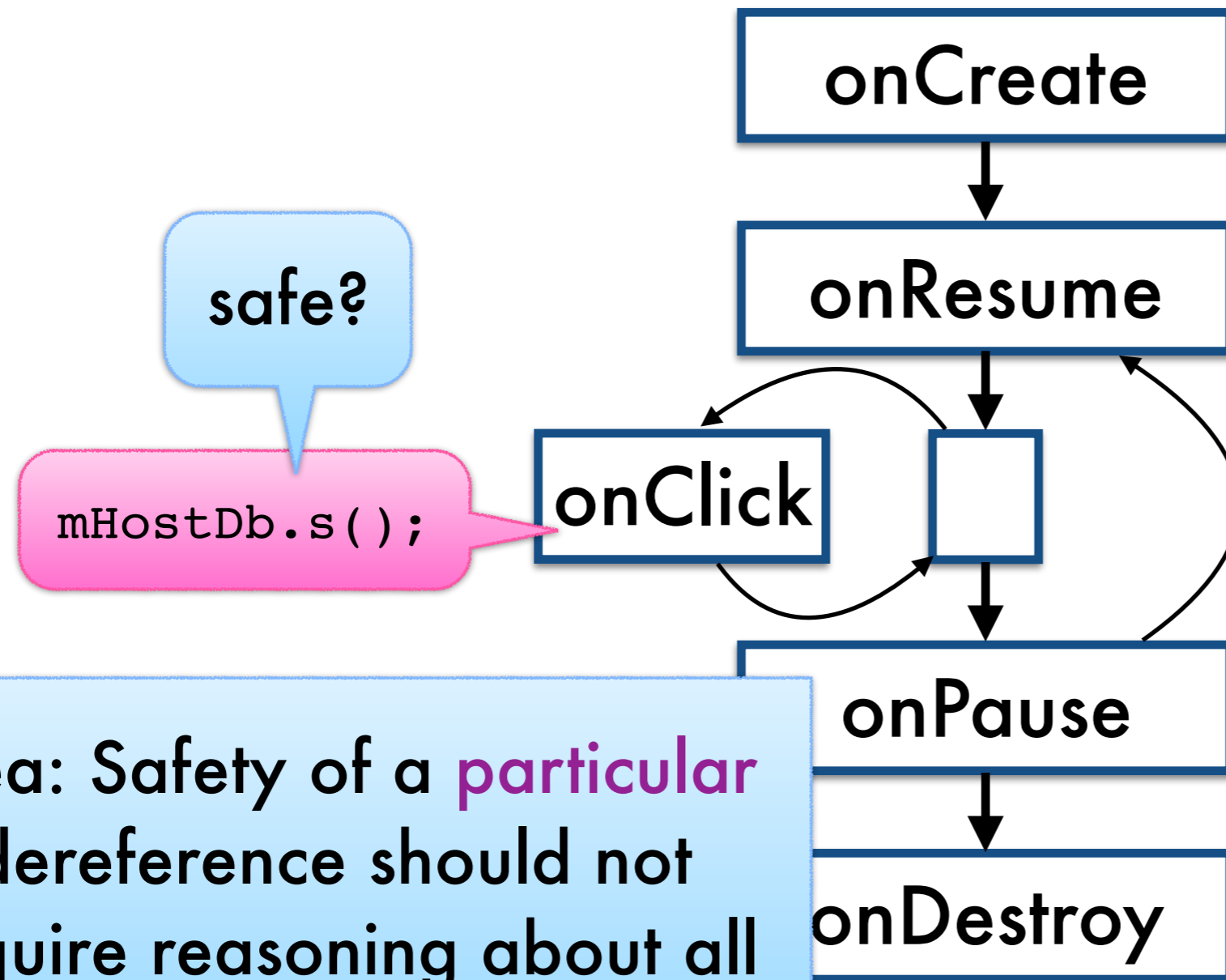
But it shouldn't be so hard ...



But it shouldn't be so hard ...

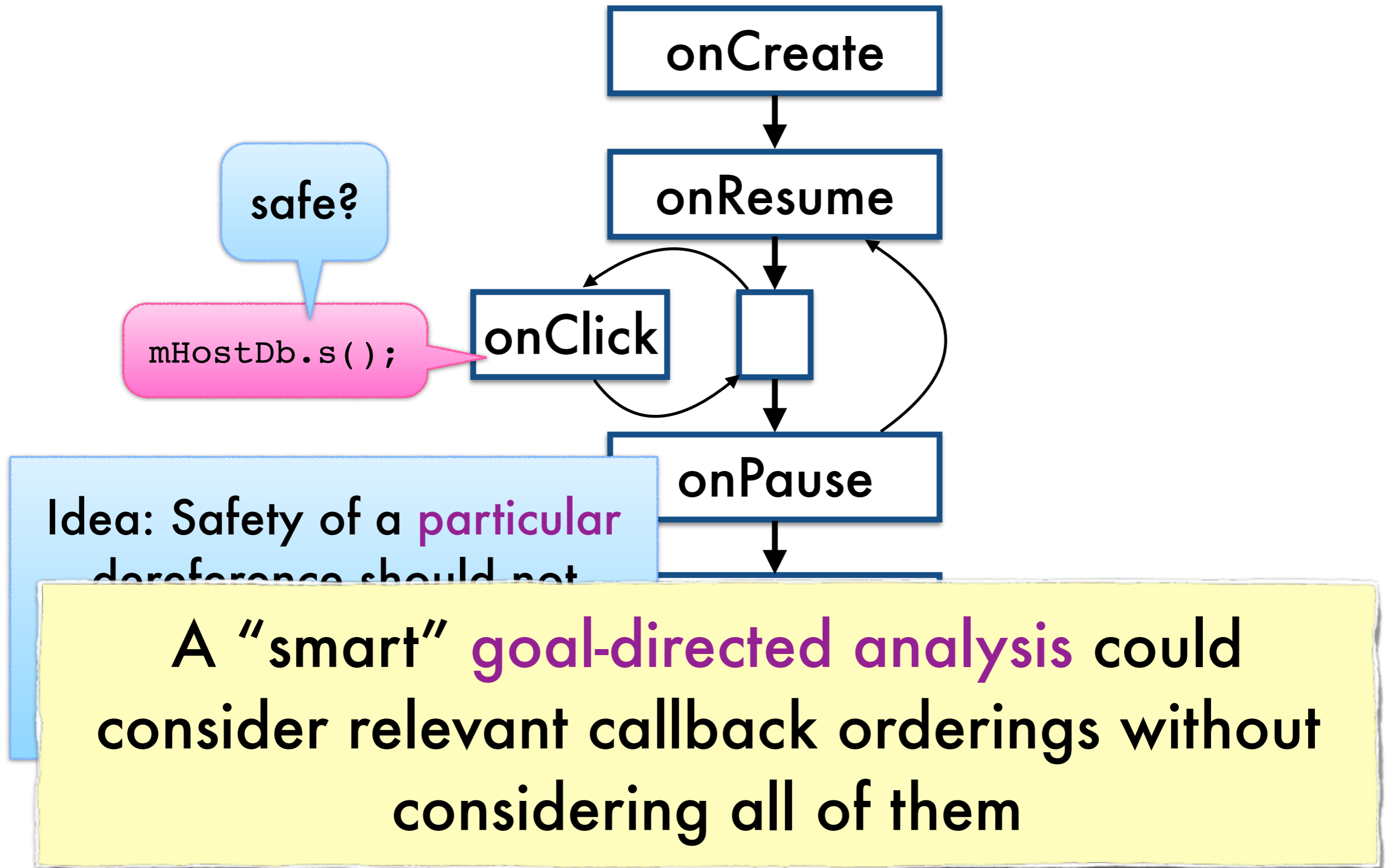


But it shouldn't be so hard ...



Idea: Safety of a **particular** dereference should not require reasoning about all callback interleavings

But it shouldn't be so hard ...



Goal-directed program analysis

safe?

`mHostDb.s();`

Goal-directed program analysis

Given a program configuration **goal**,
derive a **contradiction**
w.r.t. its reachability

safe?

`mHostDb.s();`

Goal-directed program analysis

Given a program configuration **goal**,
derive a **contradiction**
w.r.t. its reachability

safe?

`mHostDb.s();`

`mHostDb == null`

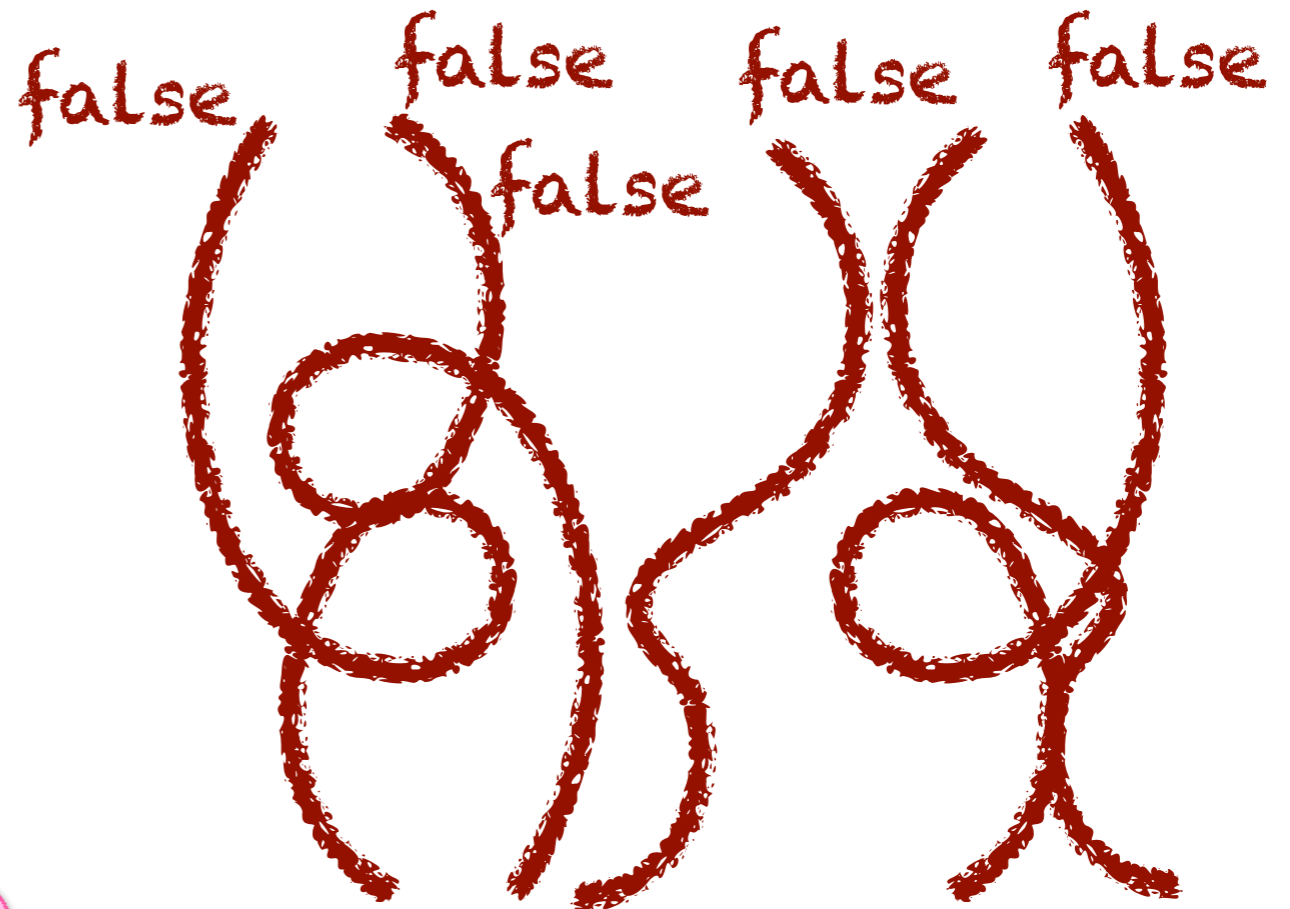
Goal-directed program analysis

Given a program configuration **goal**, derive a **contradiction** w.r.t. its reachability

safe?

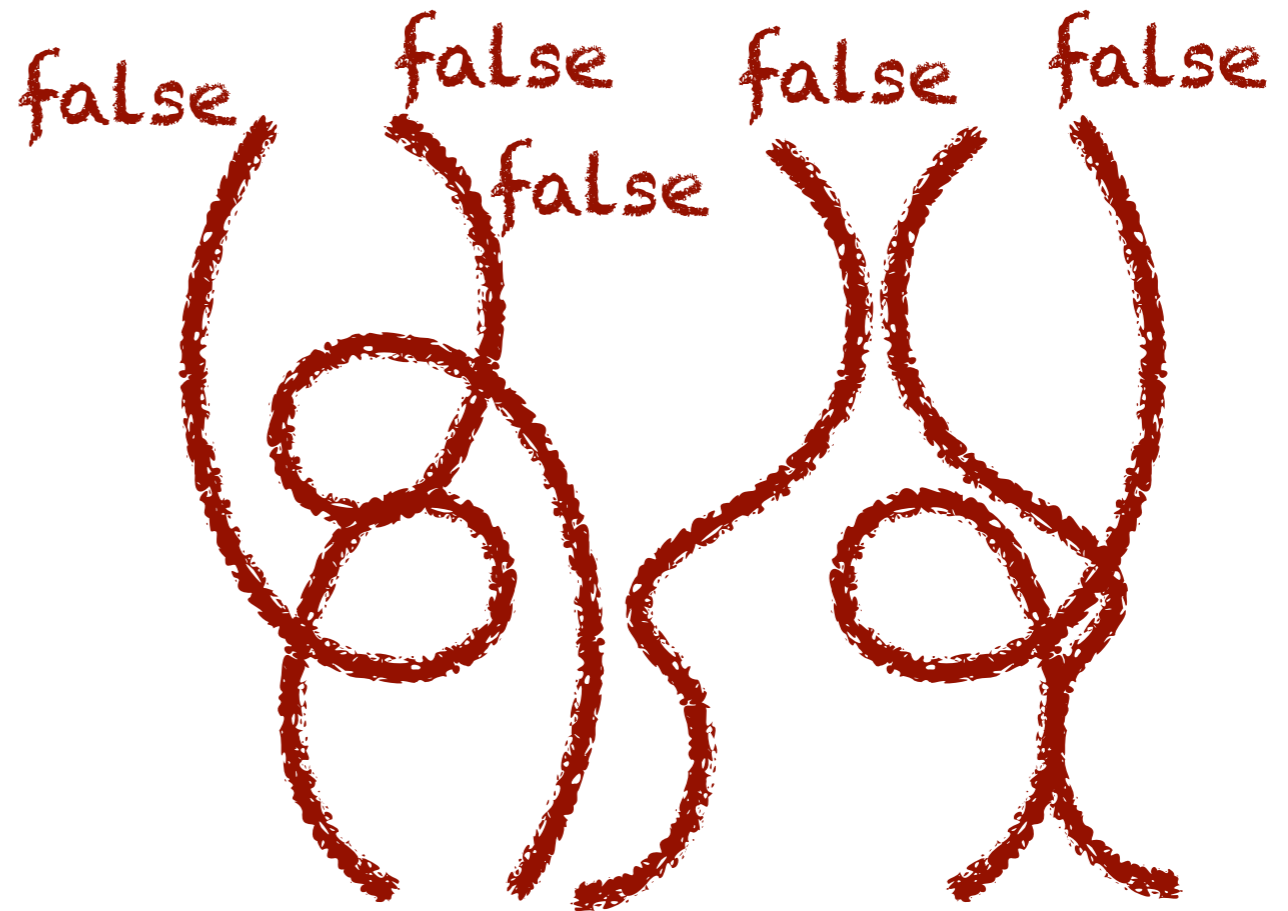
`mHostDb.s();`

`mHostDb == null`



Goal-directed program analysis

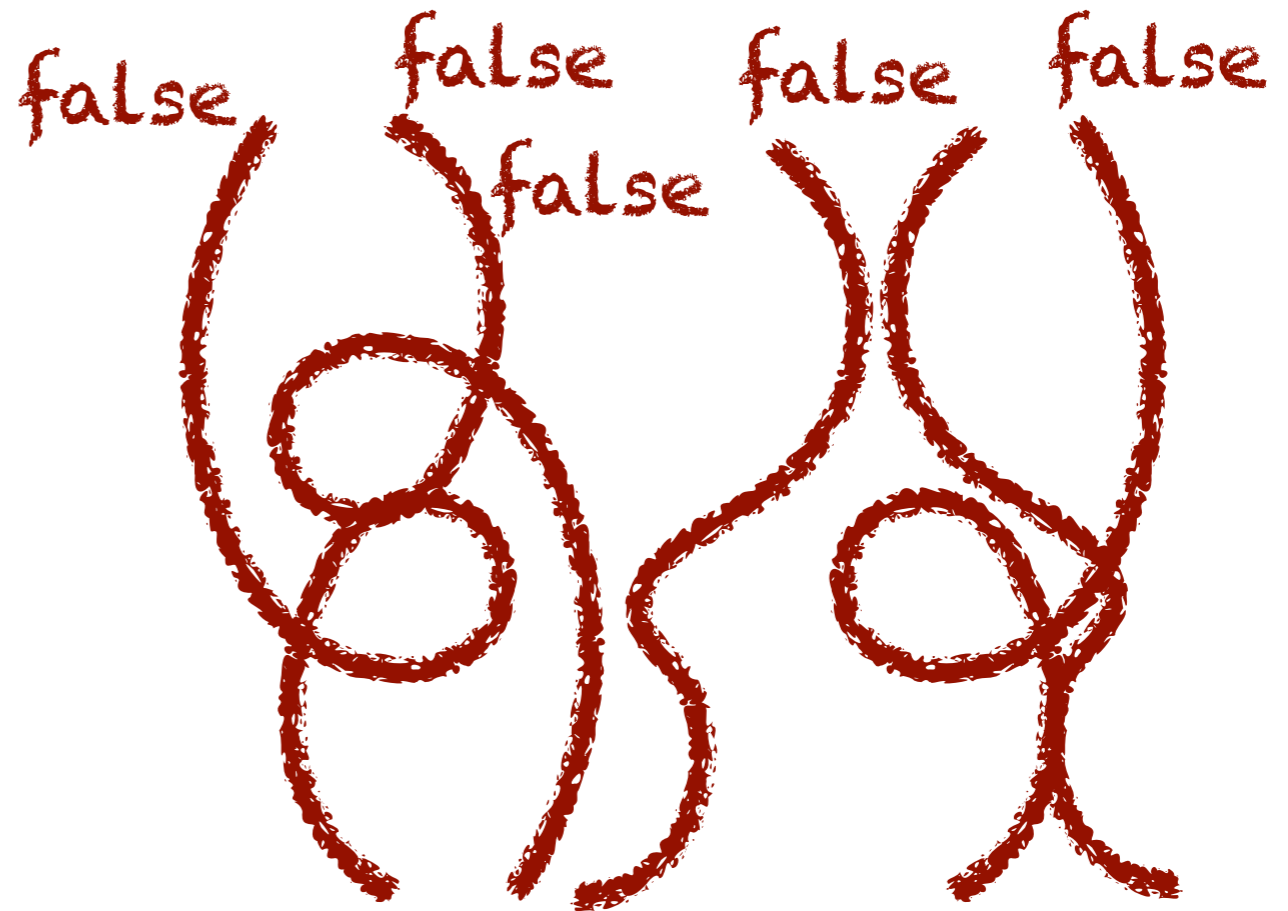
Given a program configuration **goal**,
derive a **contradiction**
w.r.t. its reachability



$$(\mathbf{this} \mapsto \hat{t} * \hat{t} \cdot \mathbf{mHostDb} \mapsto \hat{a} * \mathbf{true}) \wedge \hat{a} = \mathbf{null}$$

Goal-directed program analysis

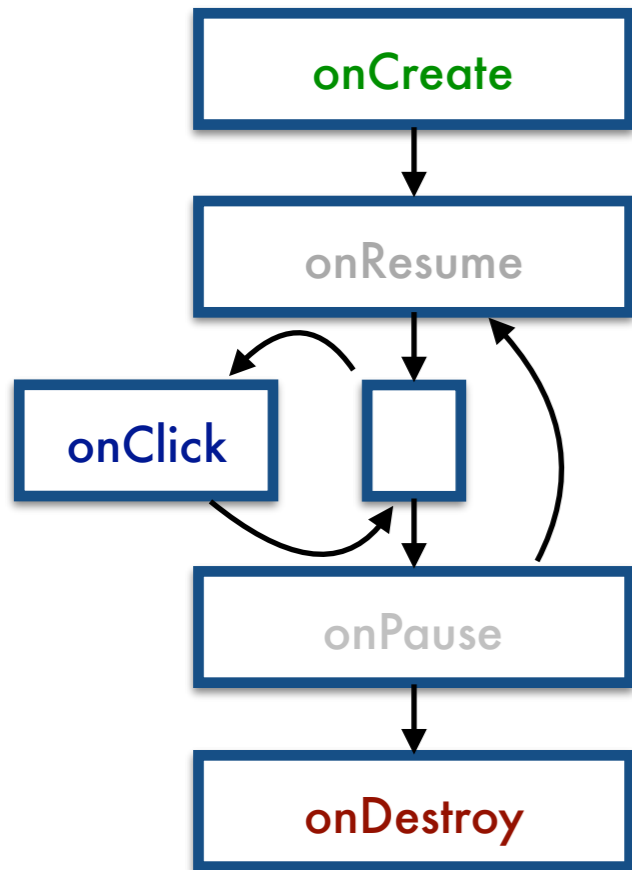
Given a program configuration **goal**, derive a **contradiction** w.r.t. its reachability



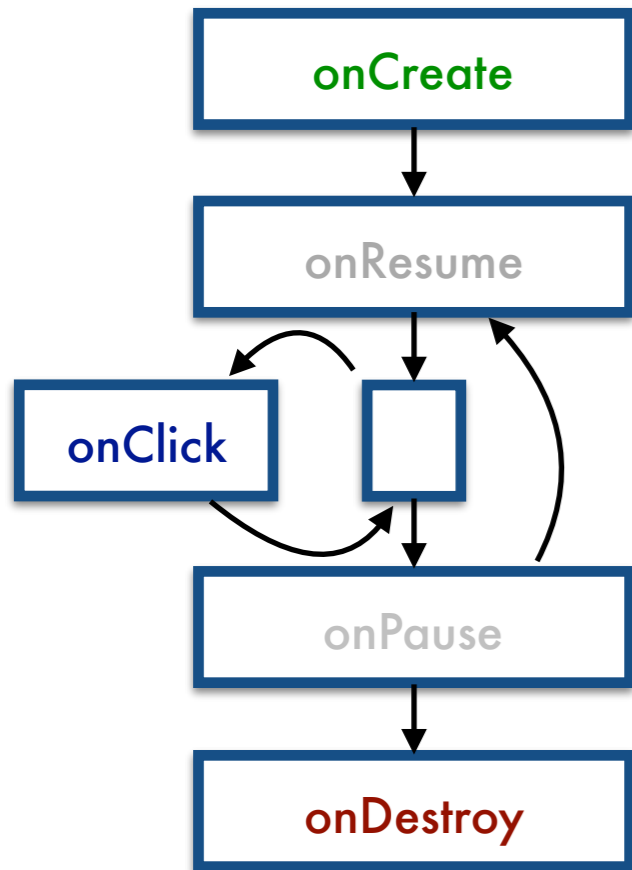
$$(\text{this} \mapsto \hat{t} * \hat{t} \cdot \text{mHostDb} \mapsto \hat{a} * \text{true}) \wedge \hat{a} = \text{null}$$

A precise backwards abstract interpretation (with separation logic constraints) to **refute** error conditions [PLDI'13]

Two dereferences: one safe and one buggy

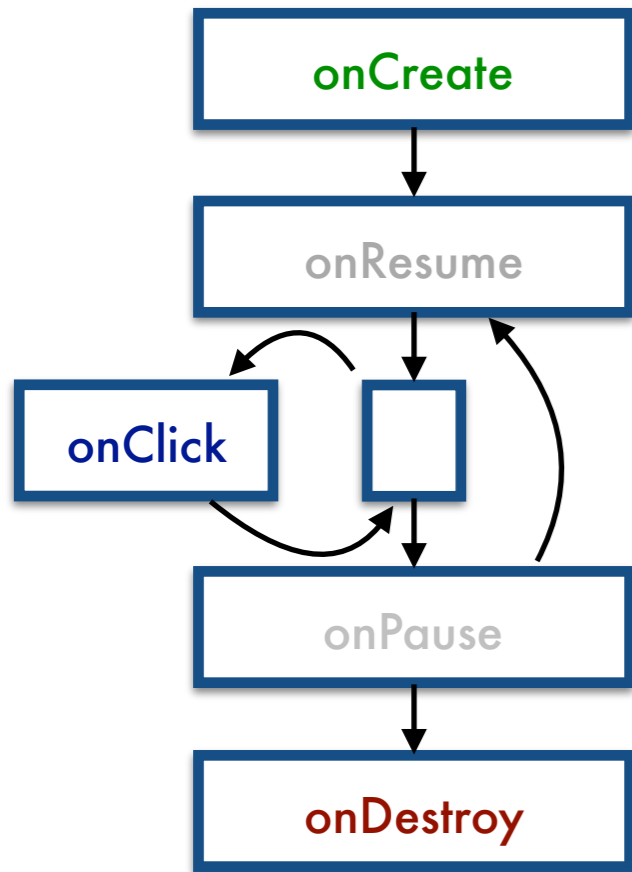


Two dereferences: one safe and one buggy



```
void onClick(...) {  
    mHostDb.s(mService.g());  
}
```

Two dereferences: one safe and one buggy

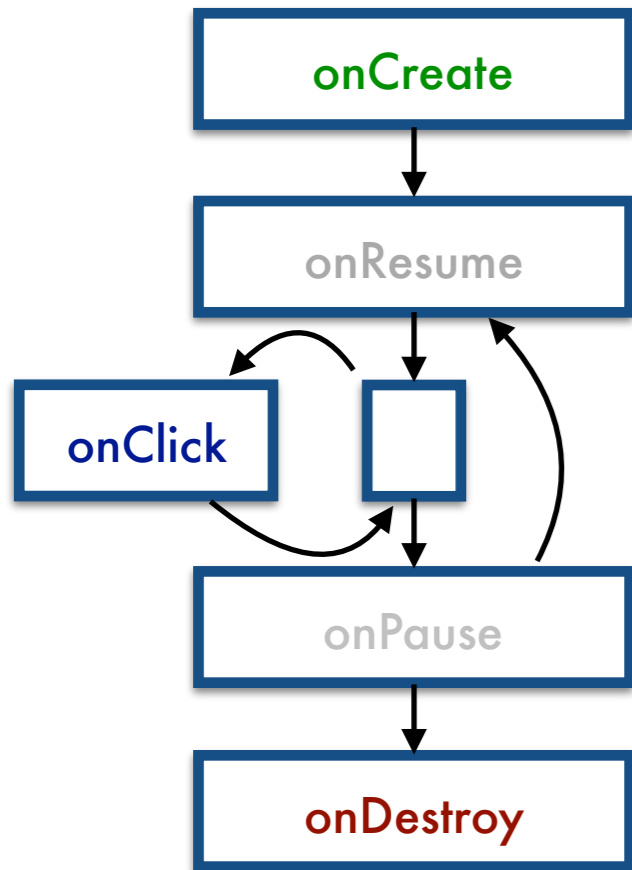


safe?

```
void onClick(...) {  
    mHostDb.s(mService.g());  
}
```

1 2

Two dereferences: one safe and one buggy



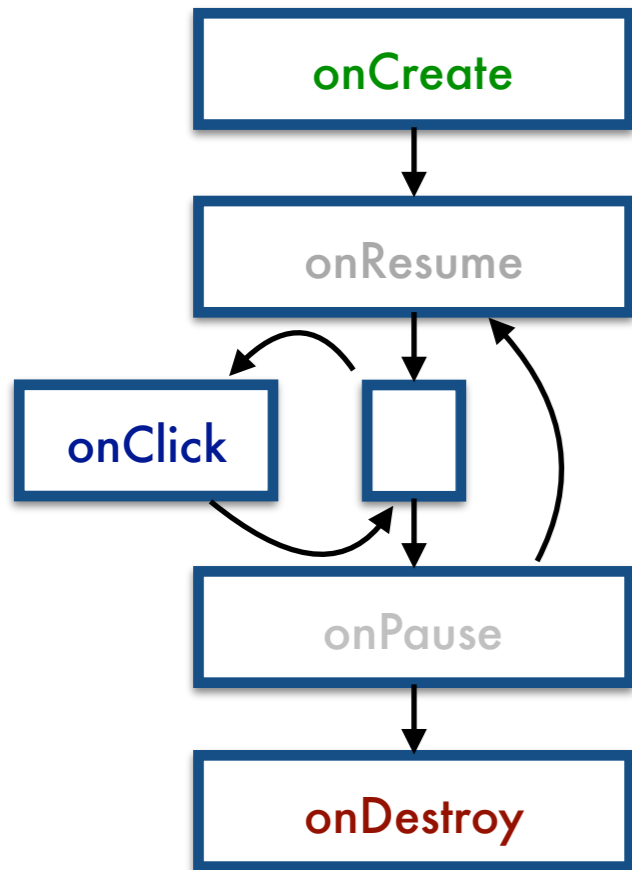
safe?

```
void onClick(...) {  
    mHostDb.s(mService.g());  
}
```

1 2

```
void onDestroy(...) {  
    mHostDb = null;  
    mService = null;  
}
```

Two dereferences: one safe and one buggy



safe?

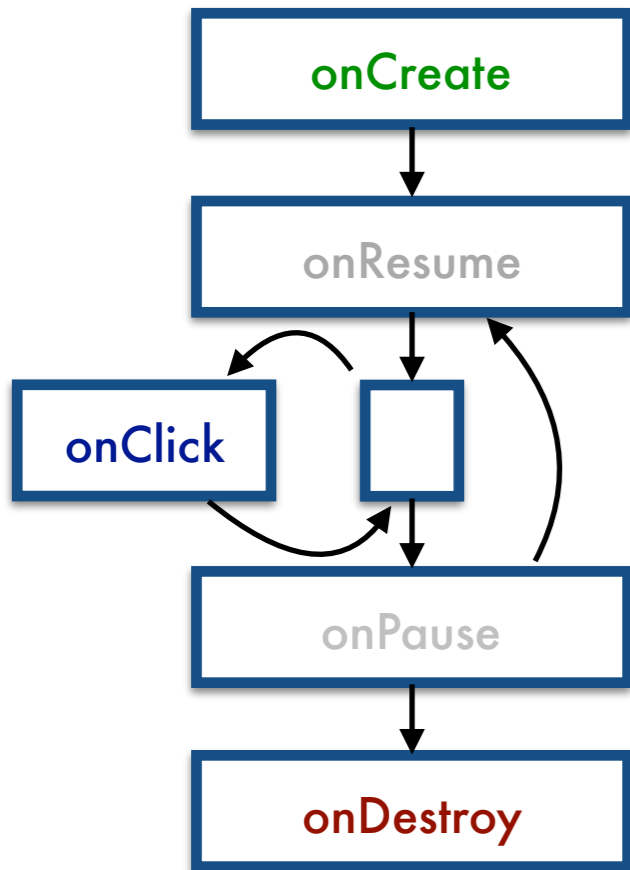
```
void onClick(...) {  
    mHostDb.s(mService.g());  
}
```

1 2

lifecycle
constraints
relevant

```
void onDestroy(...) {  
    mHostDb = null;  
    mService = null;  
}
```

Two dereferences: one safe and one buggy



safe?

lifecycle constraints relevant

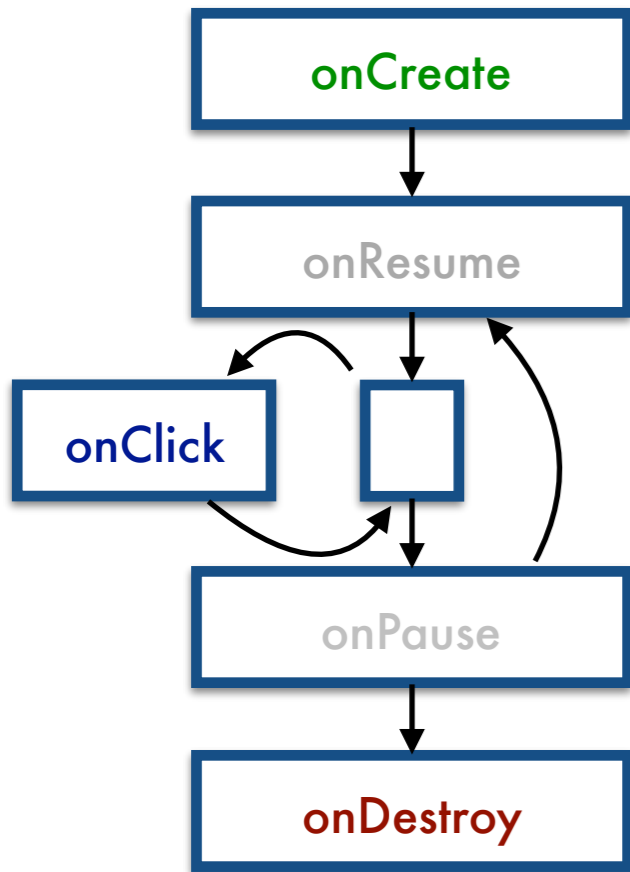
```
void onCreate() {
    bindService(..., new ServiceConn {
        void onConnected(@NonNull Service s) {
            mService = s;
        }
    });
    mHostDb = new Db();
}
```

```
void onClick(...) {
    mHostDb.s(mService.g());
}
```

1 2

```
void onDestroy(...) {
    mHostDb = null;
    mService = null;
}
```

Two dereferences: one safe and one buggy



safe?

lifecycle constraints relevant

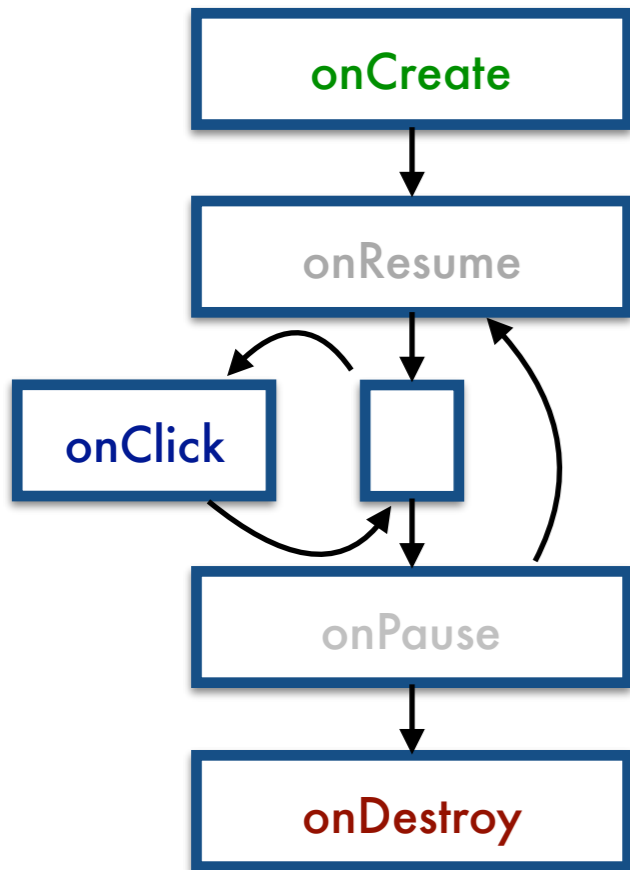
```
void onCreate() {
    bindService(..., new ServiceConn {
        void onConnected(@NonNull Service s) {
            mService = s;
        }
    });
    mHostDb = new Db();
}
```

```
void onClick(...) {
    mHostDb.s(mService.g());
}
```

1 ✓ 2

```
void onDestroy(...) {
    mHostDb = null;
    mService = null;
}
```

Two dereferences: one safe and one buggy




safe?

lifecycle constraints relevant

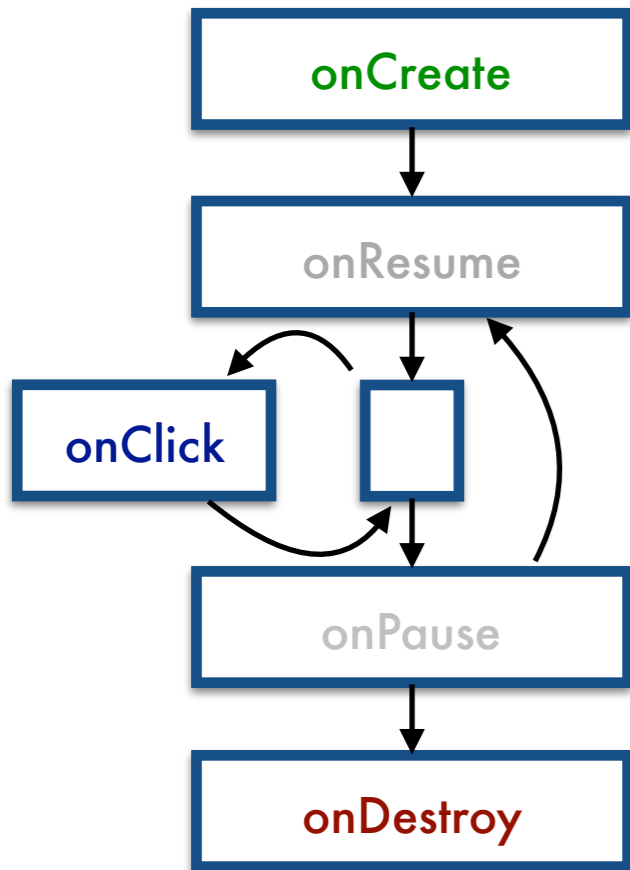
```
void onCreate() {
    bindService(..., new ServiceConn {
        void onConnected(@NonNull Service s) {
            mService = s;
        }
    });
    mHostDb = new Db();
}
```

```
void onClick(...) {
    mHostDb.s(mService.g());
}
```

1 ✓ 2 BUG 

```
void onDestroy(...) {
    mHostDb = null;
    mService = null;
}
```


Two dereferences: one safe and one buggy



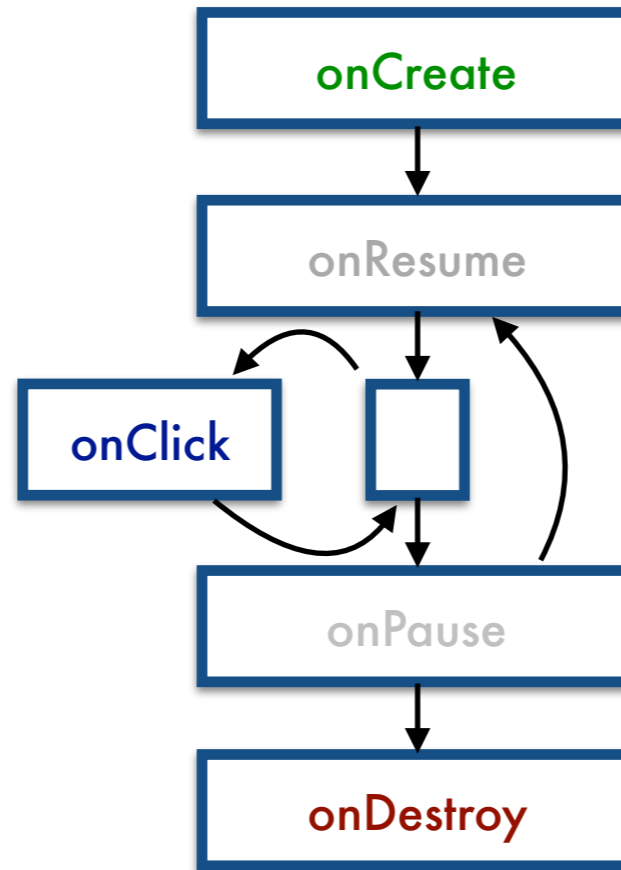
```
void onCreate() {  
    bindService(..., new ServiceConn {  
        void onConnected(@NonNull Service s) {  
            mService = s;  
        }  
    });  
    mHostDb = new Db();  
}
```

```
void onClick(...) {  
    mHostDb.s(mService.g());  
}
```

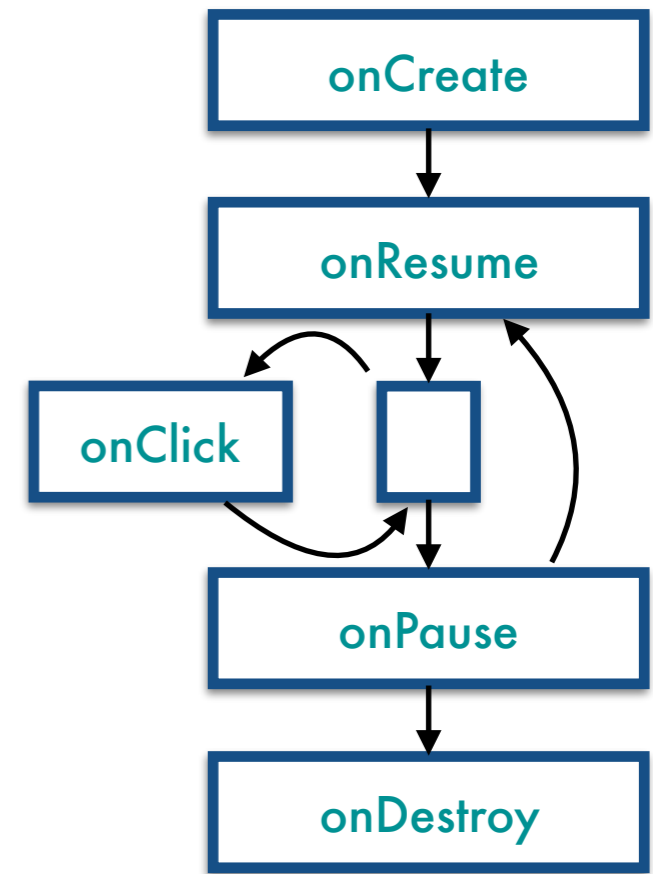
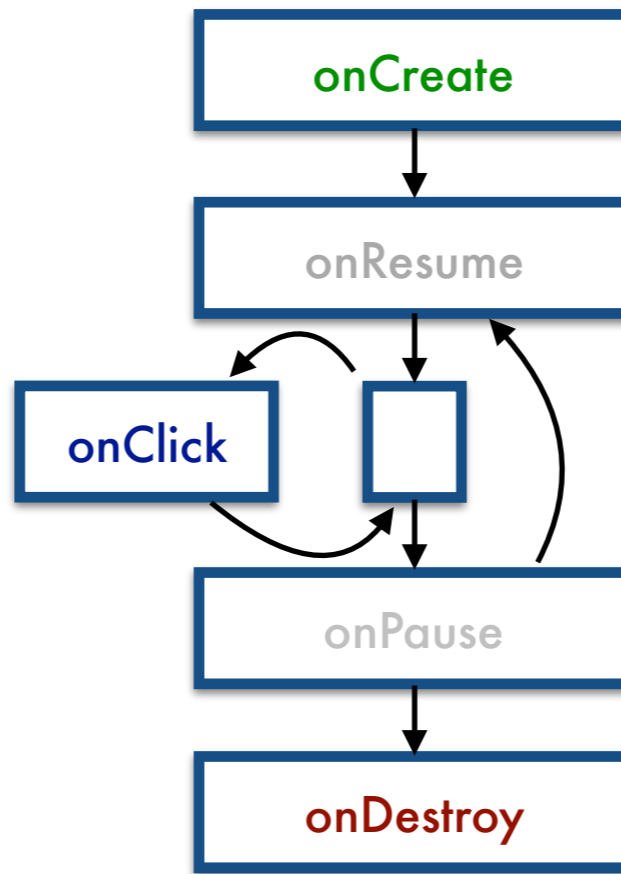
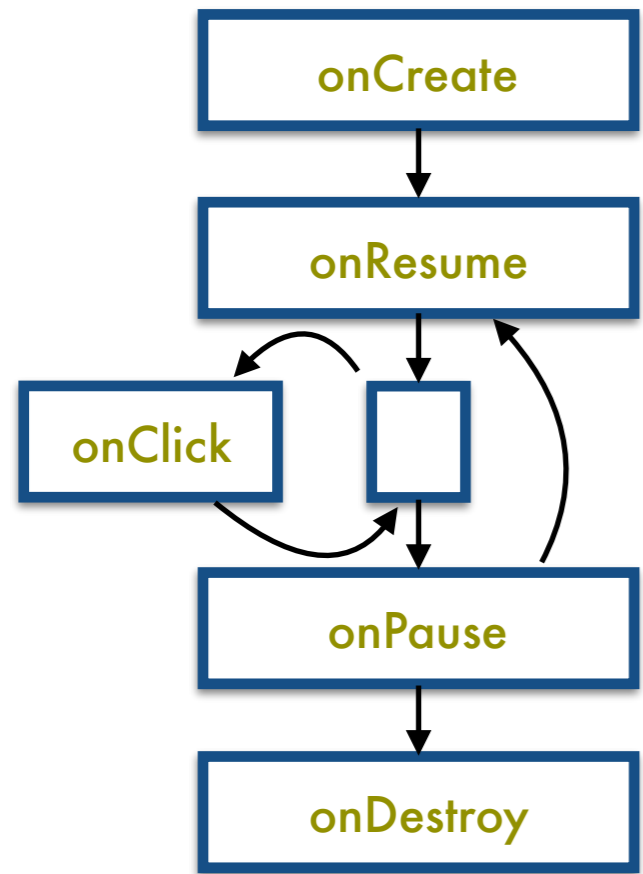
safe?

Need to consider **some but not all** callback ordering constraints

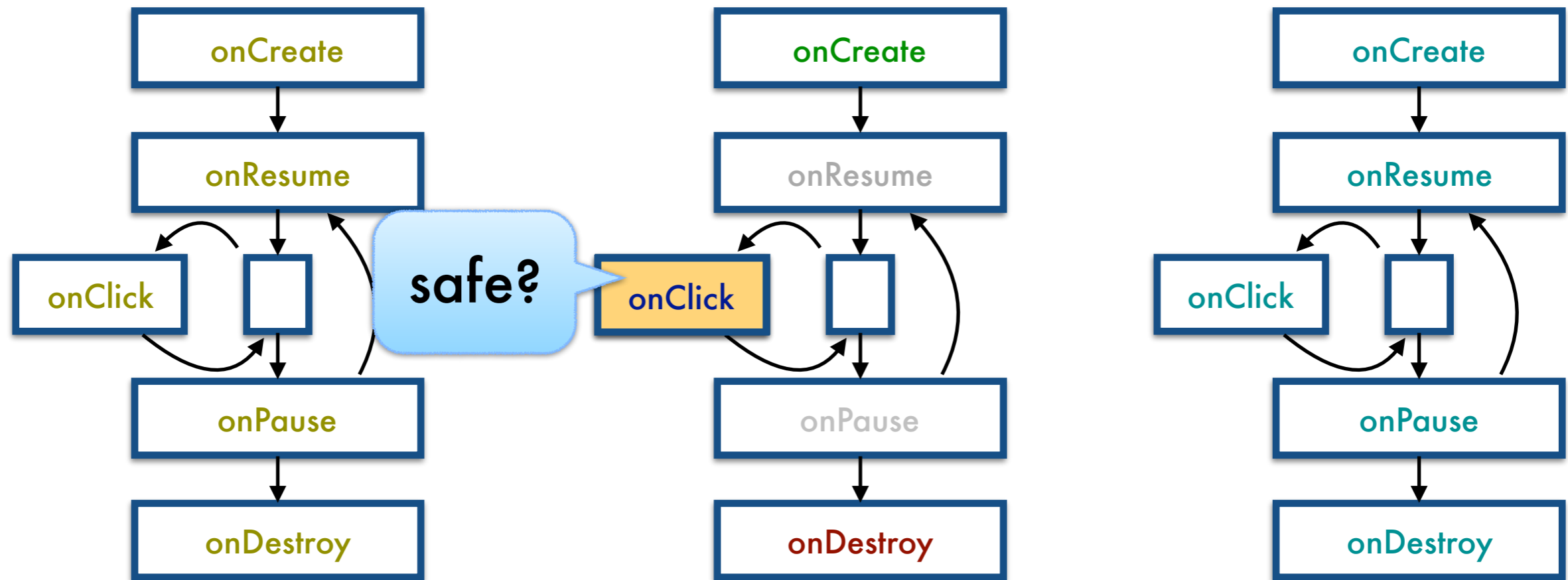
Idea: **Jumping** to relevant callbacks



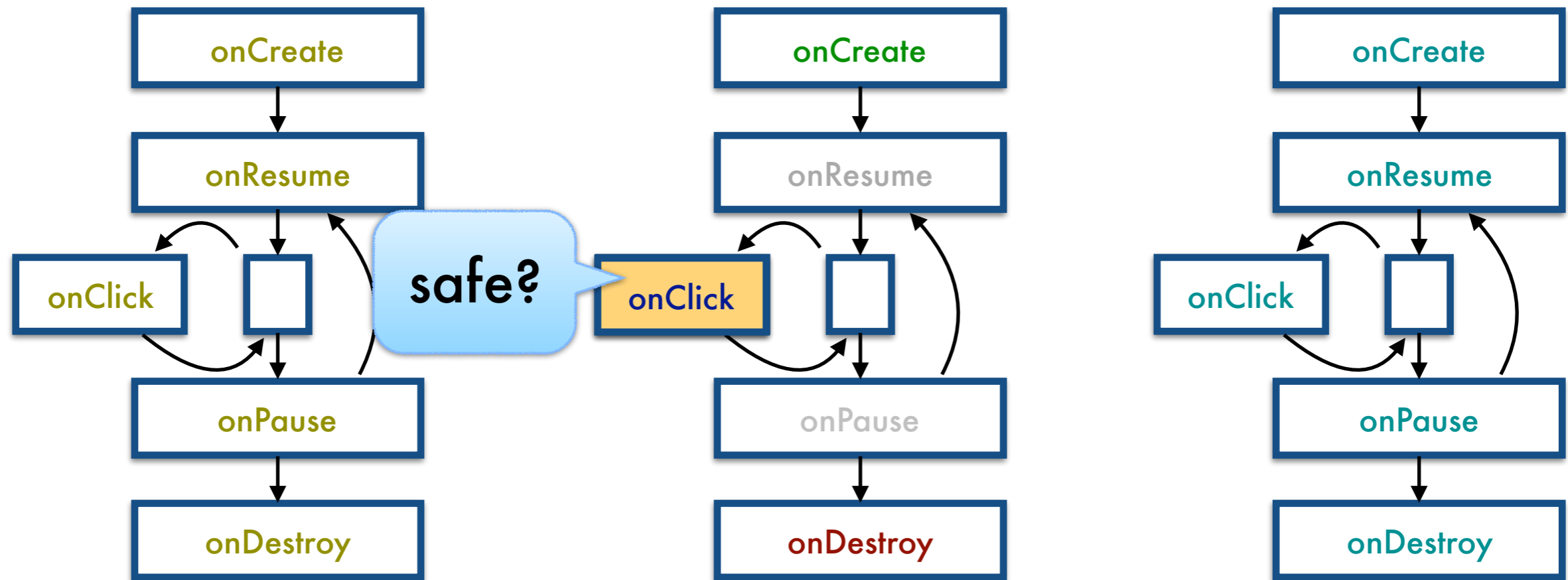
Idea: **Jumping** to relevant callbacks



Idea: **Jumping** to relevant callbacks

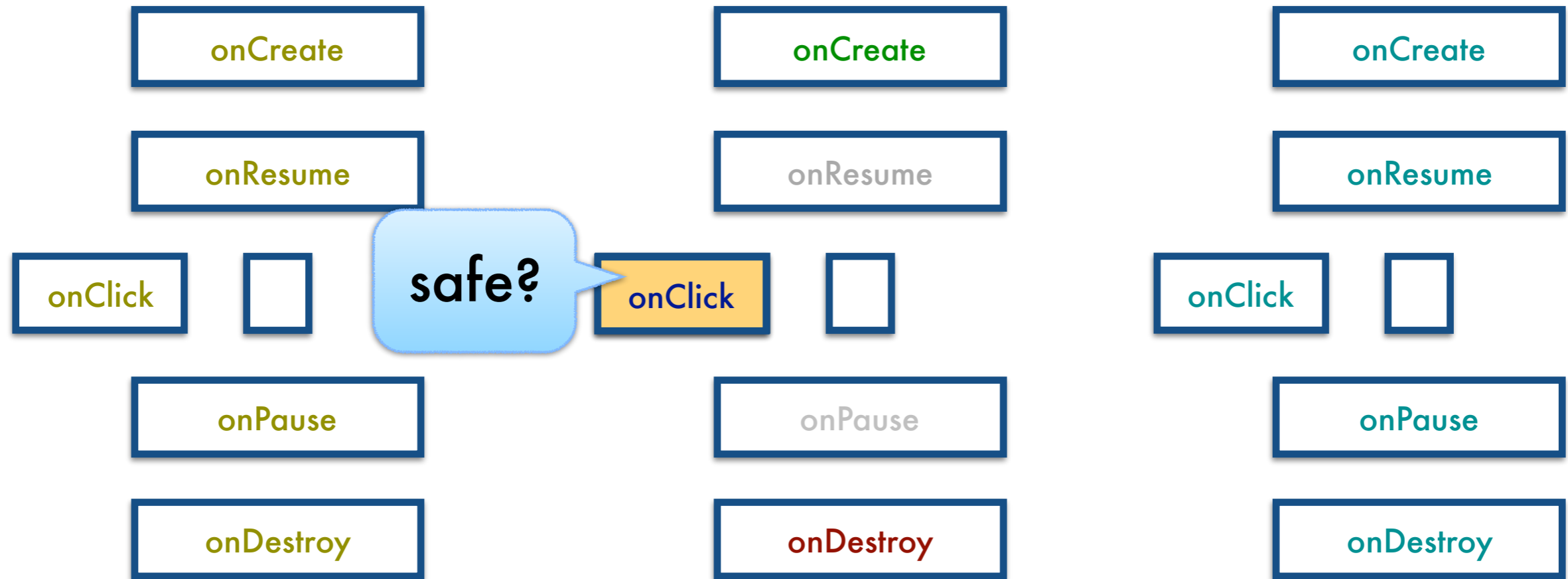


Idea: **Jumping** to relevant callbacks



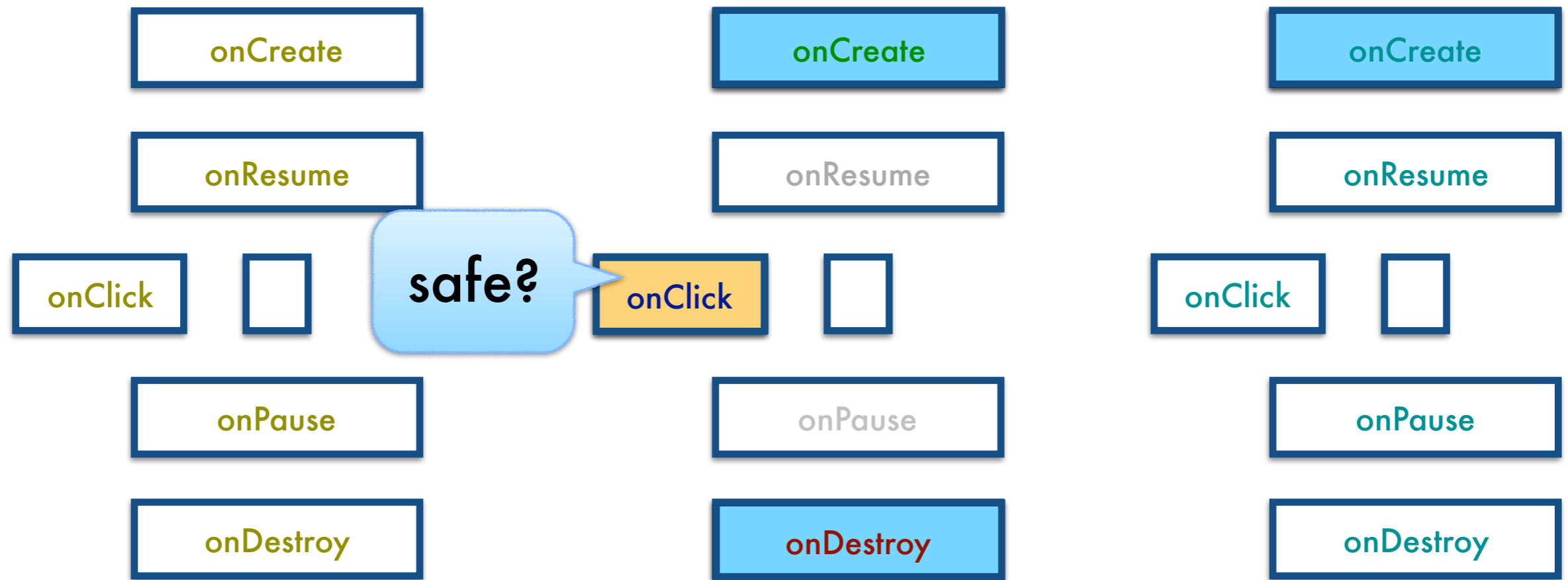
Find **data relevant** callbacks

Idea: **Jumping** to relevant callbacks



Find **data relevant** callbacks

Idea: **Jumping** to relevant callbacks



Find **data relevant** callbacks

Idea: **Jumping** to relevant callbacks



Find **data relevant** callbacks

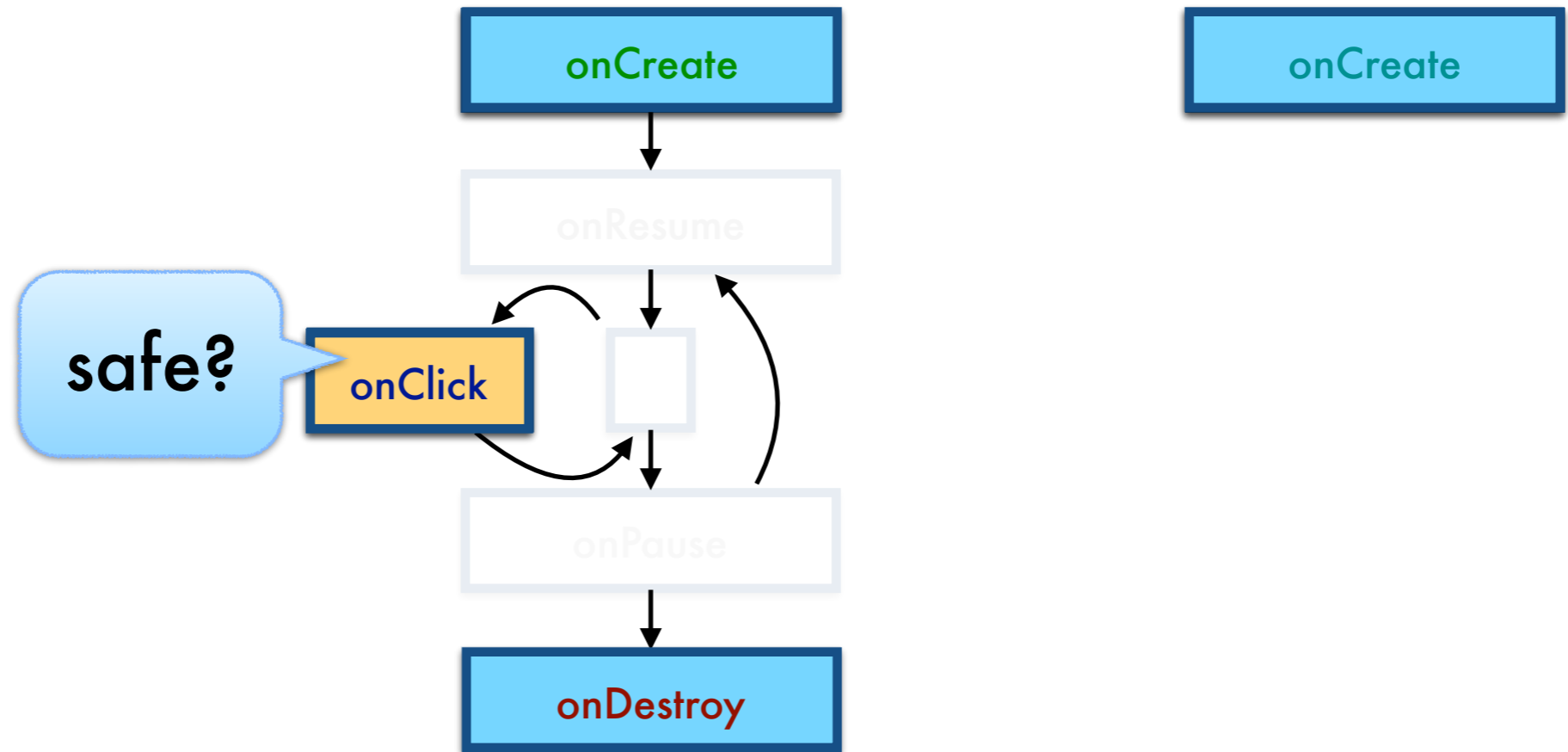
Idea: **Jumping** to relevant callbacks



Find **data relevant** callbacks

Filter using **control-feasibility**

Idea: **Jumping** to relevant callbacks



Find **data relevant** callbacks

Filter using **control-feasibility**

Idea: **Jumping** to relevant callbacks

onCreate

onCreate

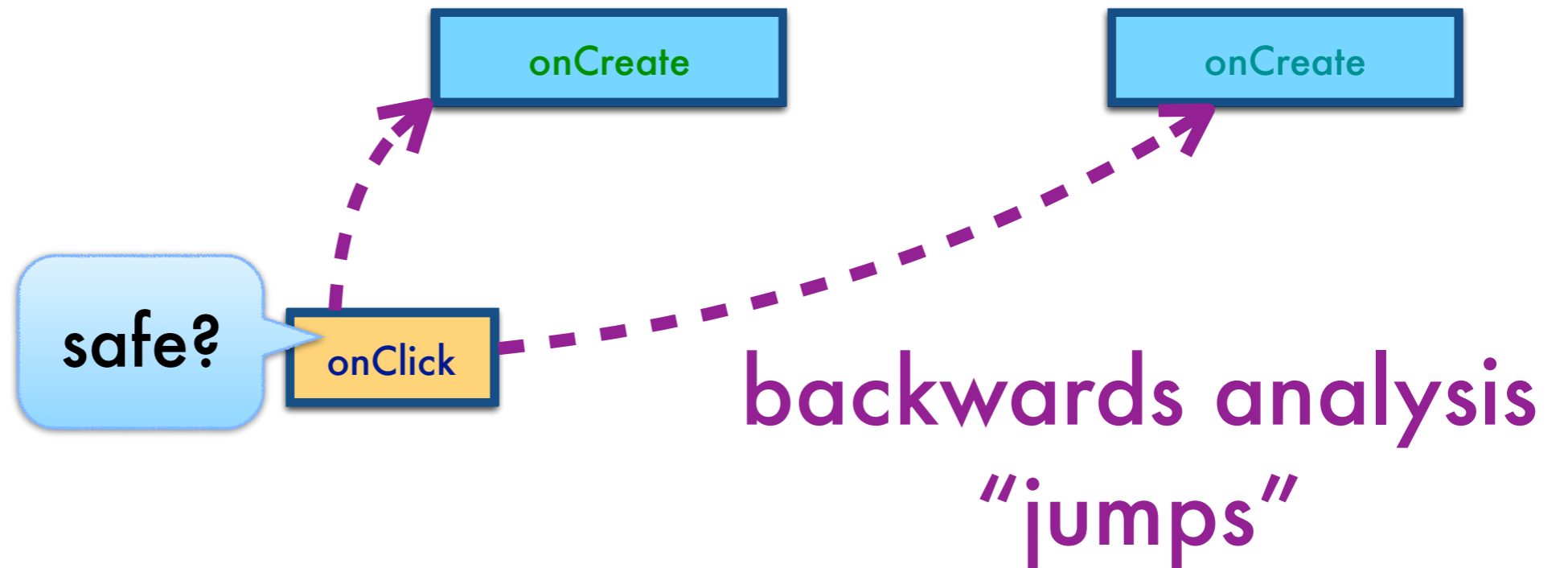
safe?

onClick

Find **data relevant** callbacks

Filter using **control-feasibility**

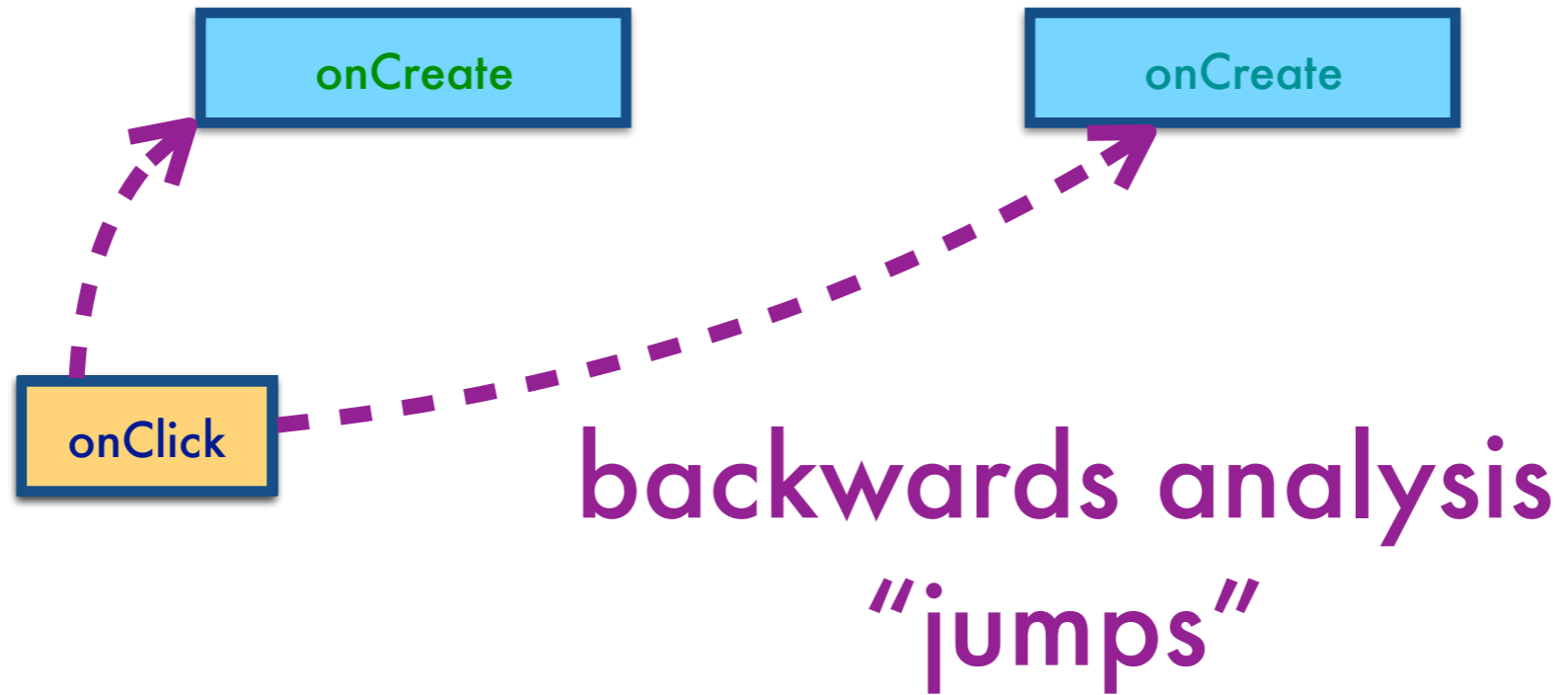
Idea: **Jumping** to relevant callbacks



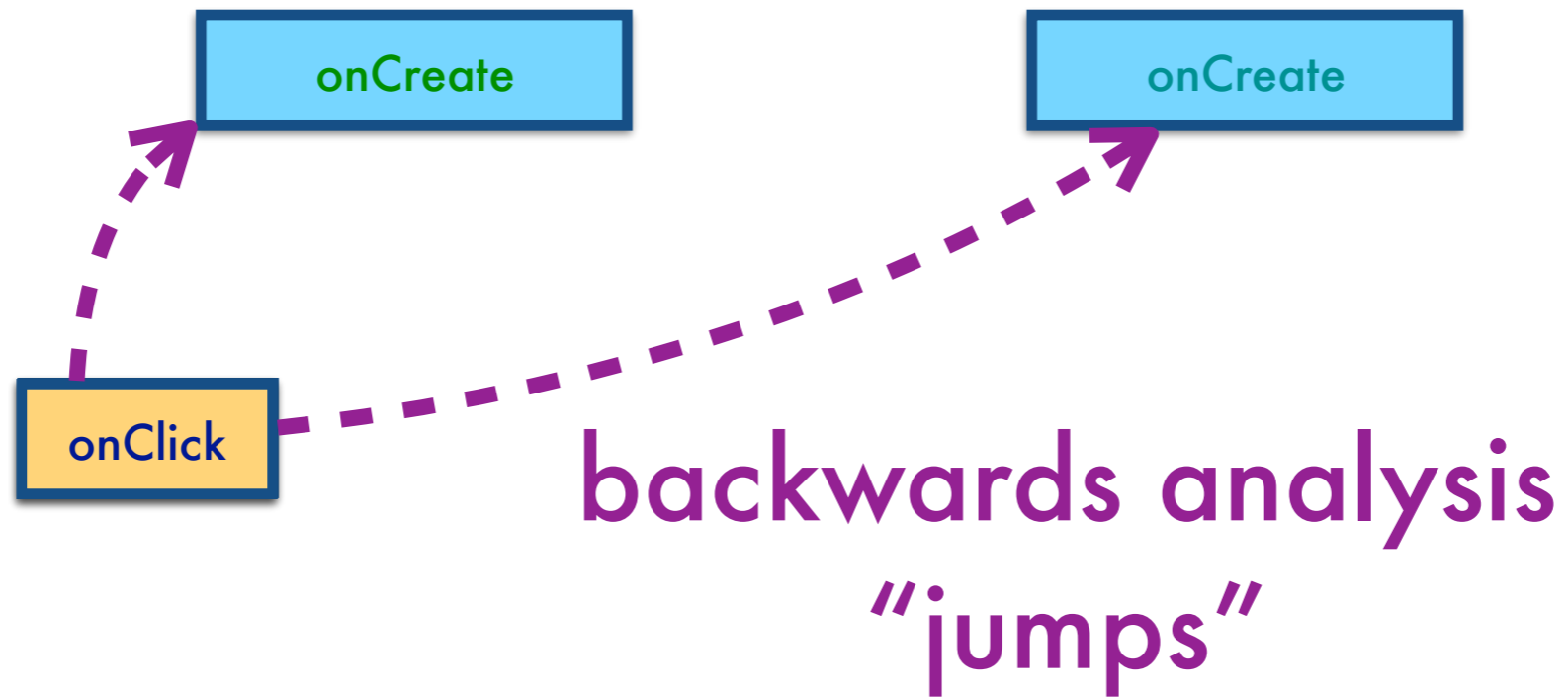
Find **data relevant** callbacks

Filter using **control-feasibility**

Contributions

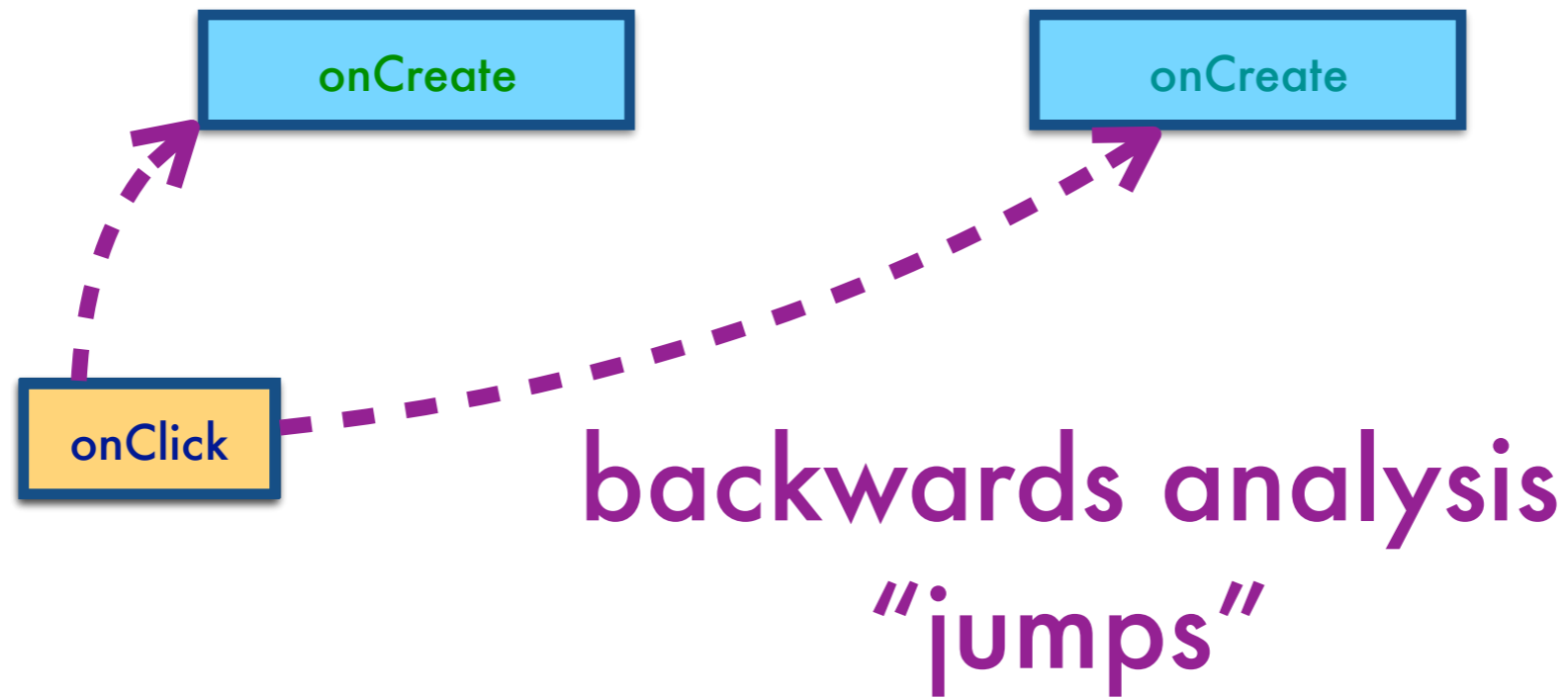


Contributions



① $l_1 \xrightarrow{[c]} l_2$ Framework for jumping analyses

Contributions



1 $l_1 \xrightarrow{[c]} l_2$ Framework for jumping analyses



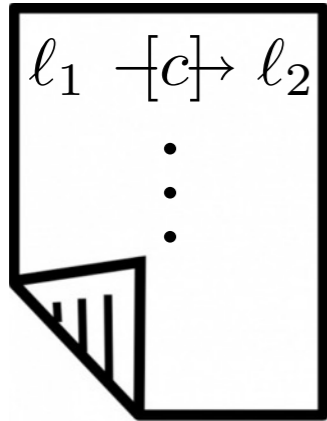
Applied to Android events

A program model for a jumping analysis

$$\overset{1}{\circ} l_1 \xrightarrow{[c]} l_2$$

A program model for a jumping analysis

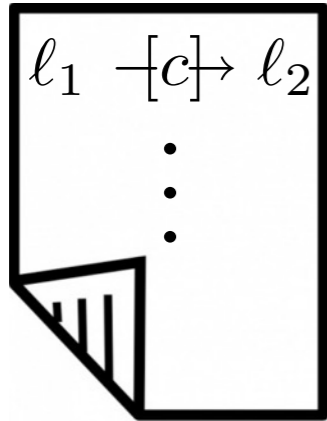
$$\overset{1}{\circ} l_1 \xrightarrow{[c]} l_2$$



Program

A program model for a jumping analysis

$$\overset{1}{\circ} l_1 \xrightarrow{[c]} l_2$$

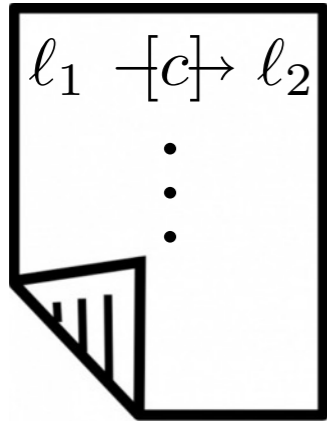


Program

**Unstructured control-flow: atomic commands
connected by pre/post labels**

A program model for a jumping analysis

$$\overset{1}{\circ} l_1 \xrightarrow{[c]} l_2$$



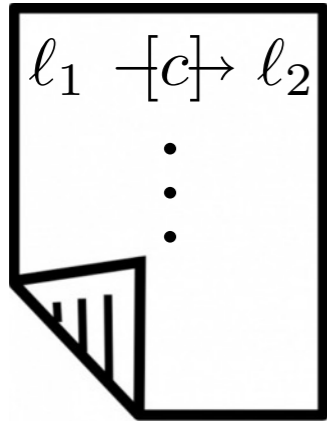
Program

Unstructured control-flow: atomic commands connected by pre/post labels

Represent conditionals with assume, loops with assume and a back edge.

A program model for a jumping analysis

$$\overset{1}{\circ} l_1 \xrightarrow{[c]} l_2$$



Program

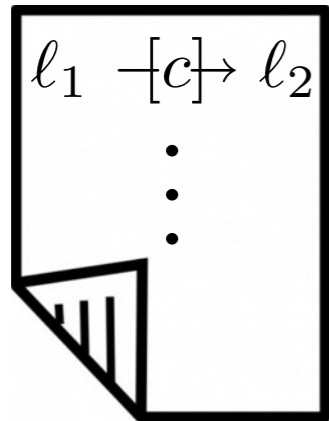
Unstructured control-flow: atomic commands connected by pre/post labels

Represent conditionals with assume, loops with assume and a back edge.

Analysis parameters

A program model for a jumping analysis

$$\overset{1}{\circ} l_1 \xrightarrow{[c]} l_2$$



Program

Unstructured control-flow: atomic commands connected by pre/post labels

Represent conditionals with assume, loops with assume and a back edge.

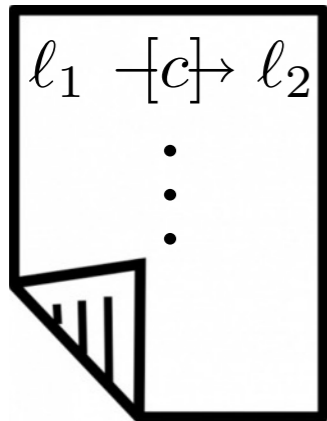
Analysis parameters

$$\{Q'\} c \{Q\}$$

Backward
transfer functions

A program model for a jumping analysis

$$\overset{1}{\circ} l_1 \xrightarrow{[c]} l_2$$



Program

Unstructured control-flow: atomic commands connected by pre/post labels

Represent conditionals with assume, loops with assume and a back edge.

Analysis parameters

$$\{Q'\} \ c \ \{Q\}$$

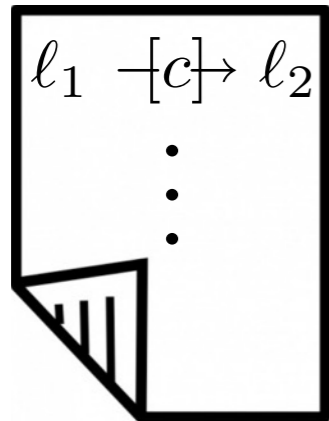
Backward
transfer functions

$$l_{\text{cur}}$$

Current
location

A program model for a jumping analysis

$$\overset{1}{\circ} l_1 \xrightarrow{[c]} l_2$$



Program

Unstructured control-flow: atomic commands connected by pre/post labels

Represent conditionals with assume, loops with assume and a back edge.

Analysis parameters

$$\{Q'\} \ c \ \{Q\}$$

Backward
transfer functions

$$l_{\text{cur}}$$

Current
location

$$Q_{\text{cur}}$$

Current
query

Data-relevance identifies writes

$$\overset{1}{\circ} l_1 \xrightarrow{[c]} l_2$$

(Q, l_{cur})

Data-relevance identifies writes

$$\overset{1}{\circ} l_1 \xrightarrow{[c]} l_2$$

(Q, l_{cur})

Data-relevance

$l_i \quad l_j \quad l_k$
 \dots

Identify locations that can write to the query Q .

Data-relevance identifies writes

$$\textcircled{1} l_1 \xrightarrow{[c]} l_2$$

(Q, l_{cur})

Data-relevance

l_i l_j l_k
...

Identify locations that can write to the query Q .

Computed using pre-pass points-to analysis, types, field-based, ...

Data-relevance identifies writes

$$\textcircled{1} \ell_1 \xrightarrow{[c]} \ell_2$$

(Q, ℓ_{cur})

Data-relevance

$\ell_i \quad \ell_j \quad \ell_k$
...

Identify locations that can write to the query Q .

Computed using pre-pass points-to analysis, types, field-based, ...

Classic idea: Following data dependencies yields a **sparse** analysis (but, here, flow-insensitive)

Control-feasibility **selectively** recovers
flow-sensitivity

$$\textcircled{1} \quad l_1 \xrightarrow{[c]} l_2$$

(Q, l_{cur})



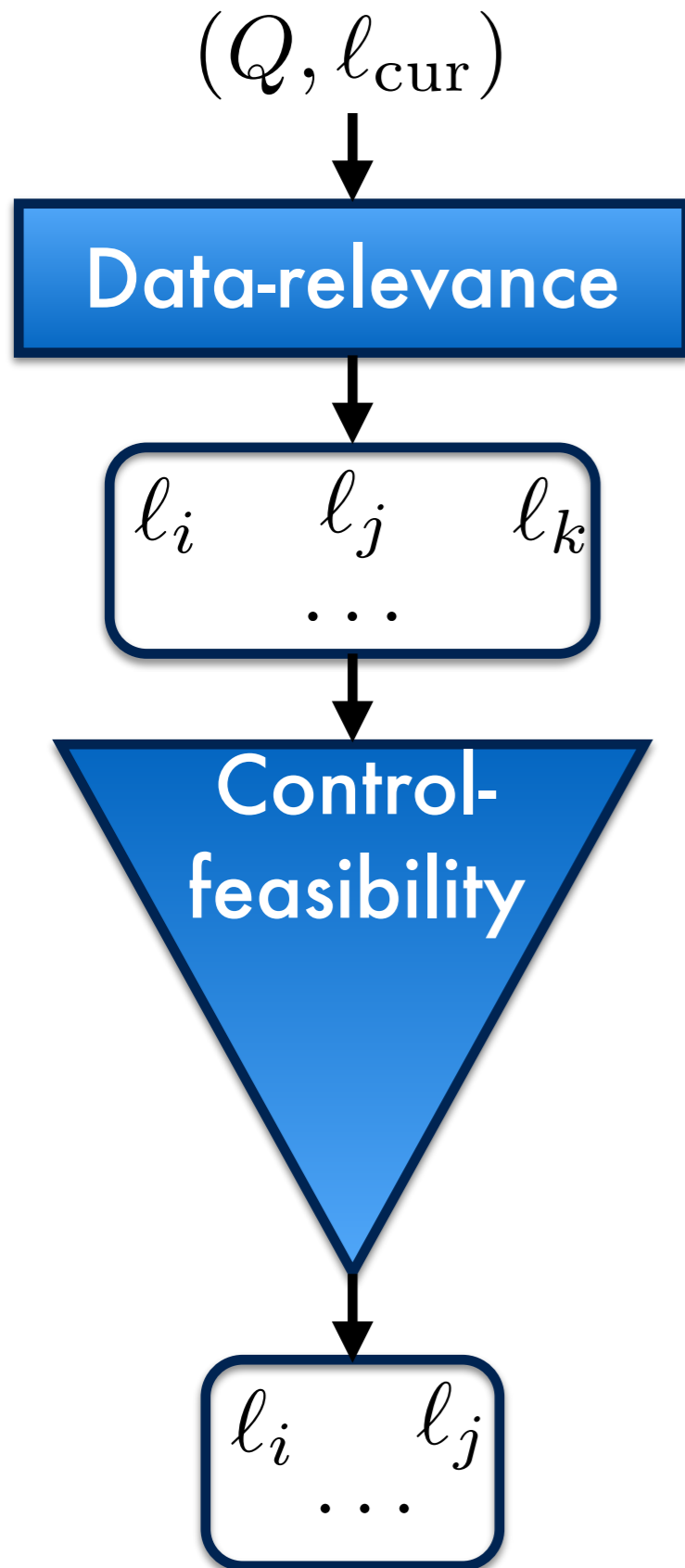
Data-relevance



$l_i \quad l_j \quad l_k$
 \dots

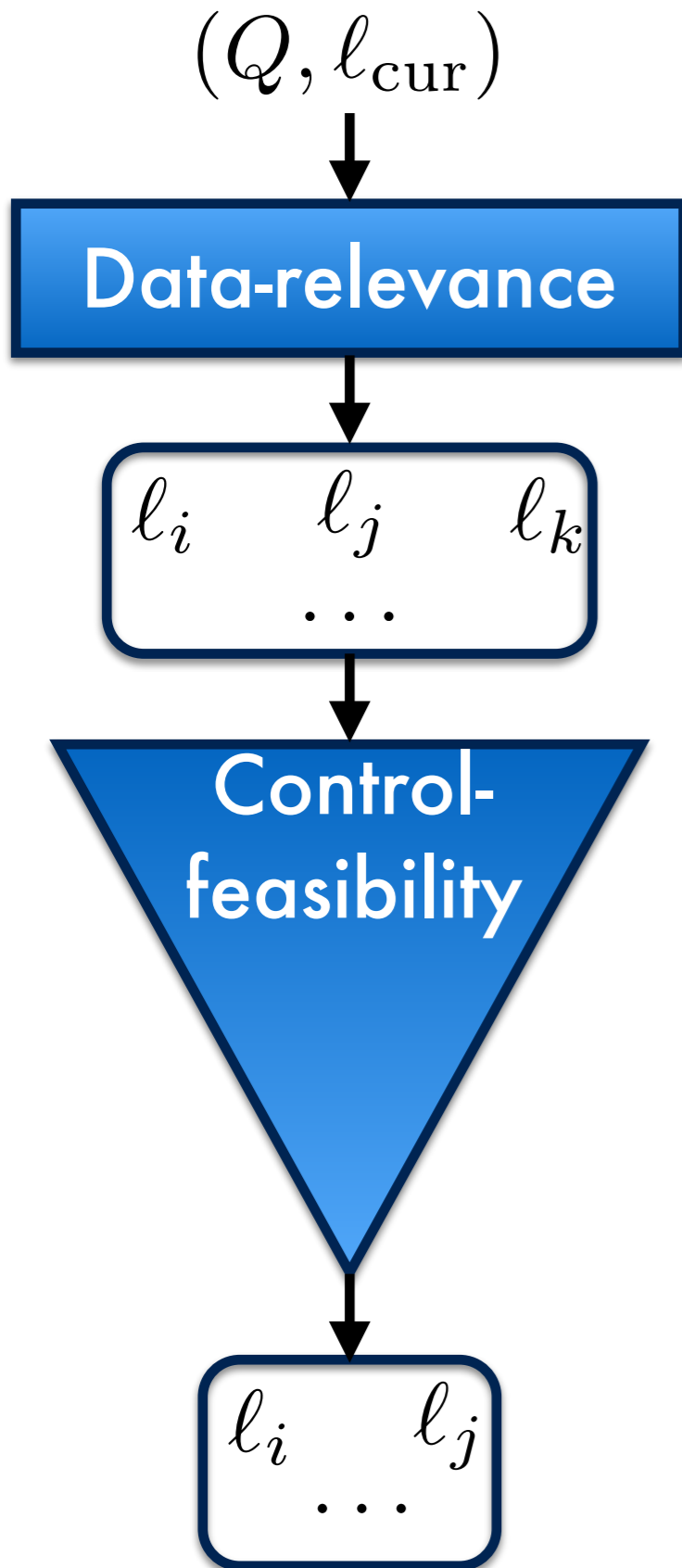
Control-feasibility **selectively** recovers flow-sensitivity

$$\textcircled{1} \ell_1 \xrightarrow{[c]} \ell_2$$



Control-feasibility **selectively** recovers flow-sensitivity

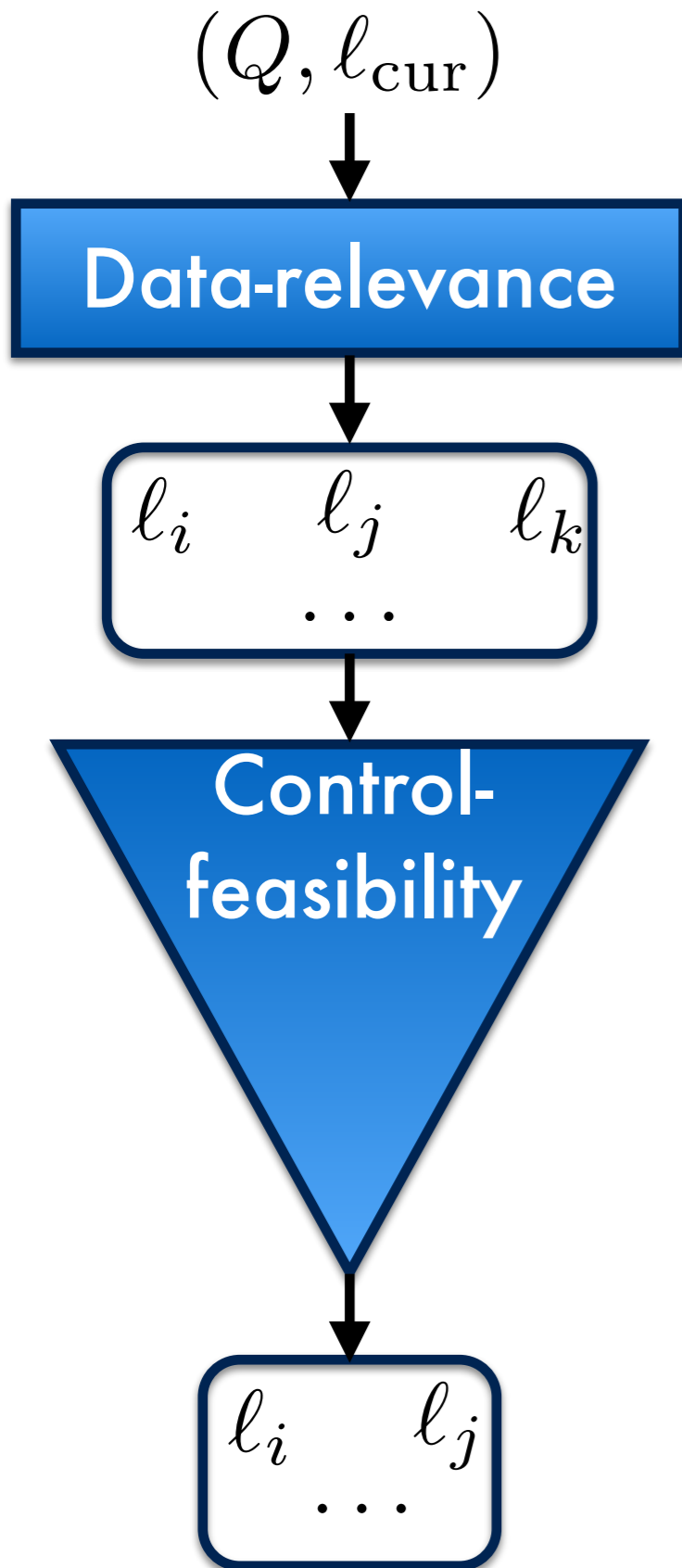
$$\textcircled{1} l_1 \xrightarrow{[c]} l_2$$



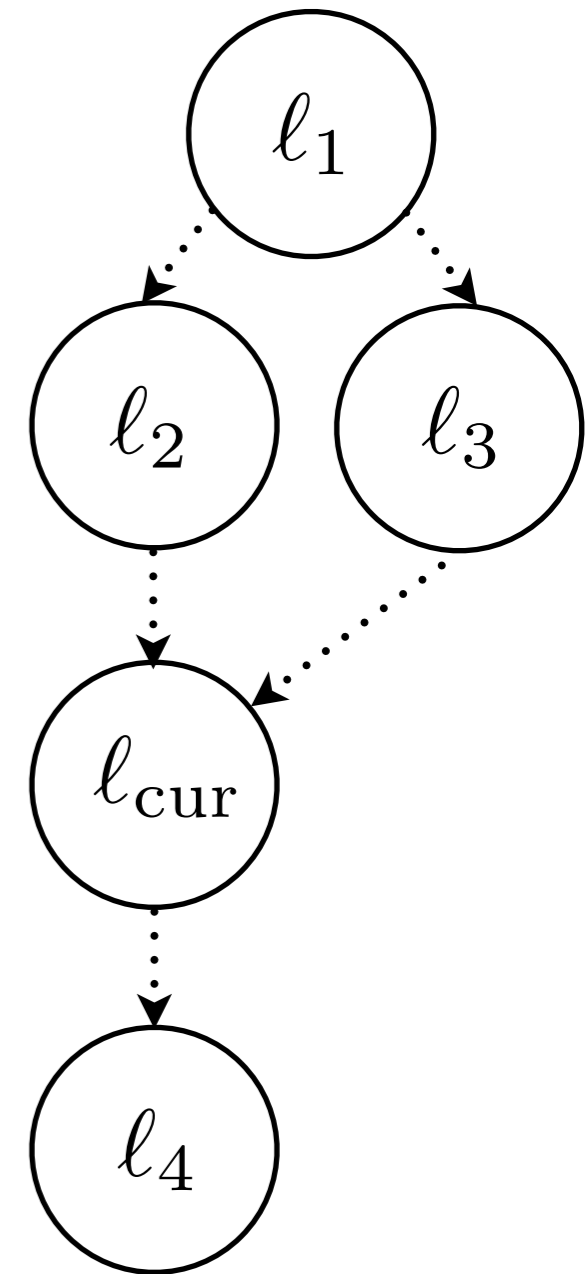
Filter the set of data-relevant locations using control flow and the current program point

Control-feasibility **selectively** recovers flow-sensitivity

① $l_1 \xrightarrow{[c]} l_2$

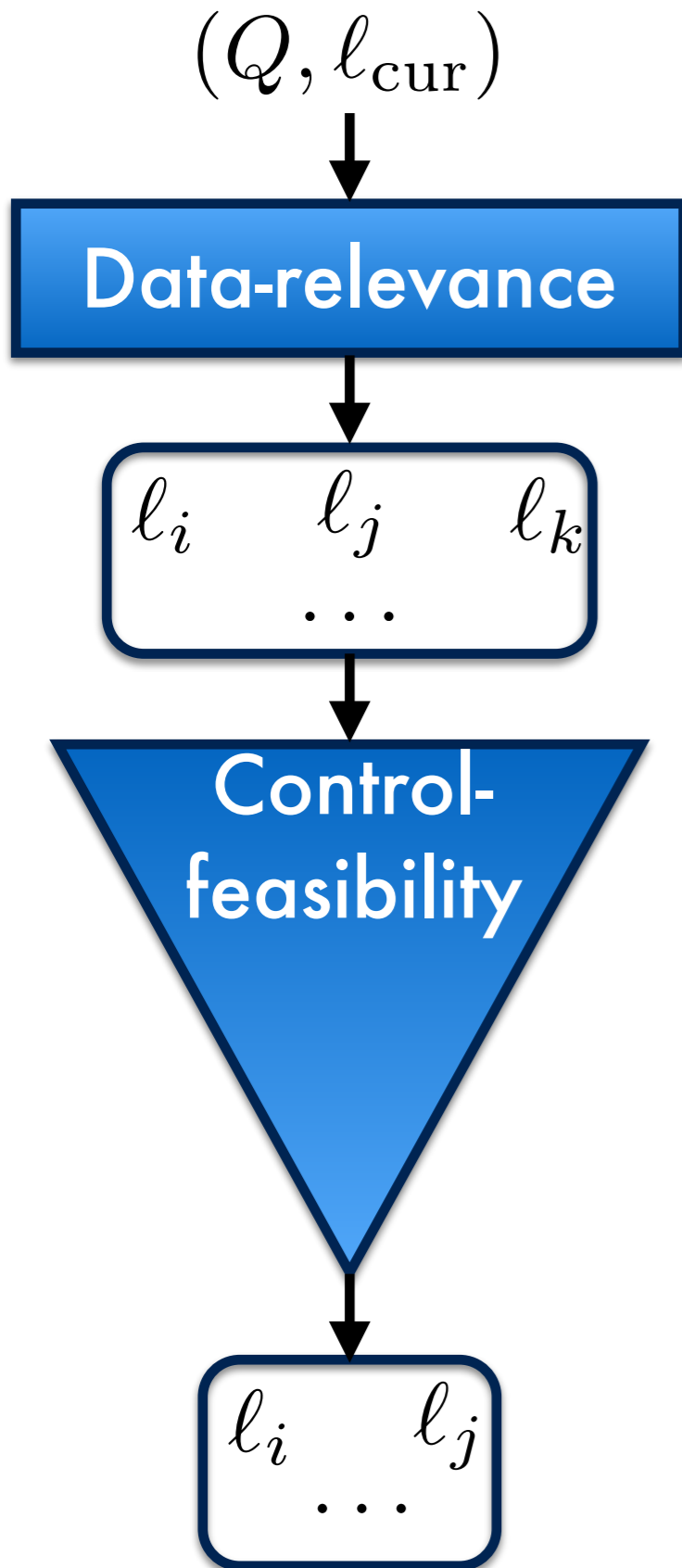


Filter the set of data-relevant locations using control flow and the current program point



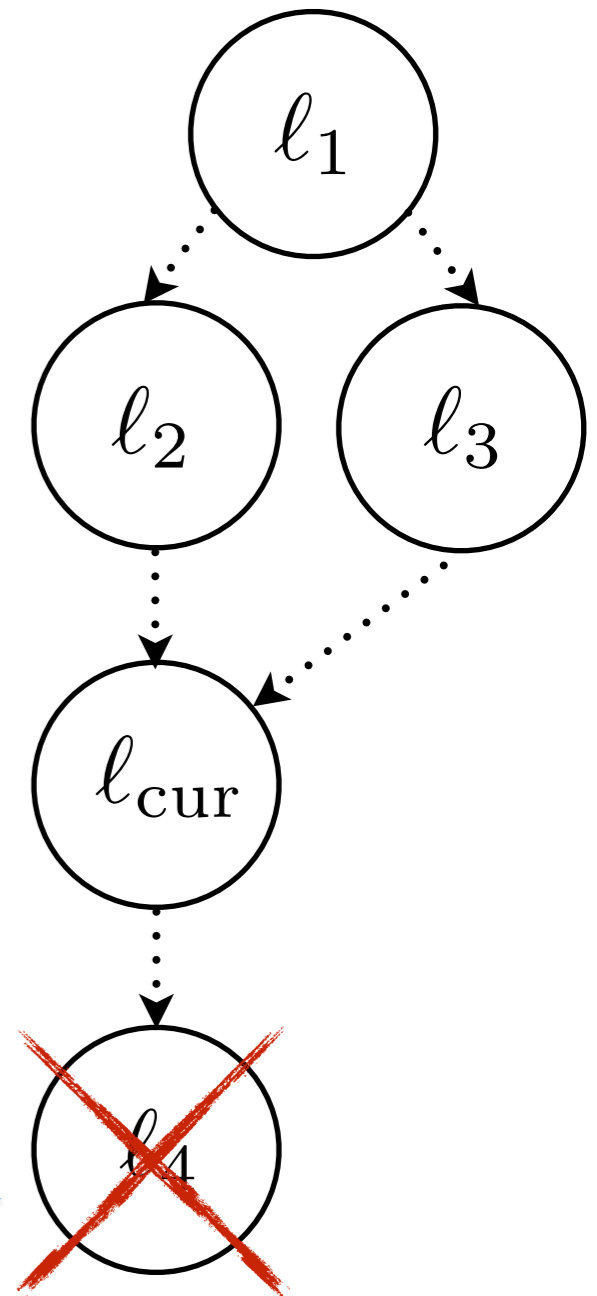
Control-feasibility **selectively** recovers flow-sensitivity

① $l_1 \xrightarrow{[c]} l_2$



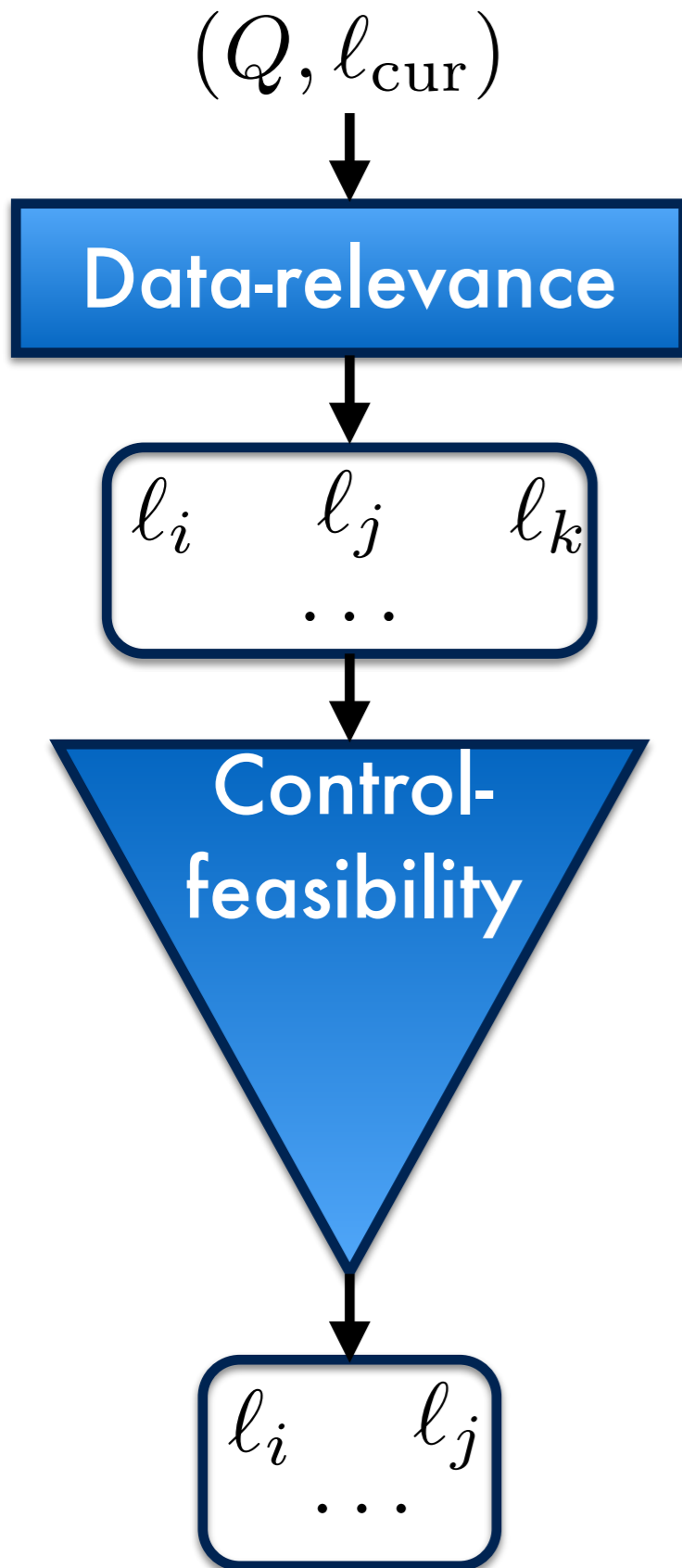
Filter the set of data-relevant locations using control flow and the current program point

Not backward-reachable from current location



Control-feasibility **selectively** recovers flow-sensitivity

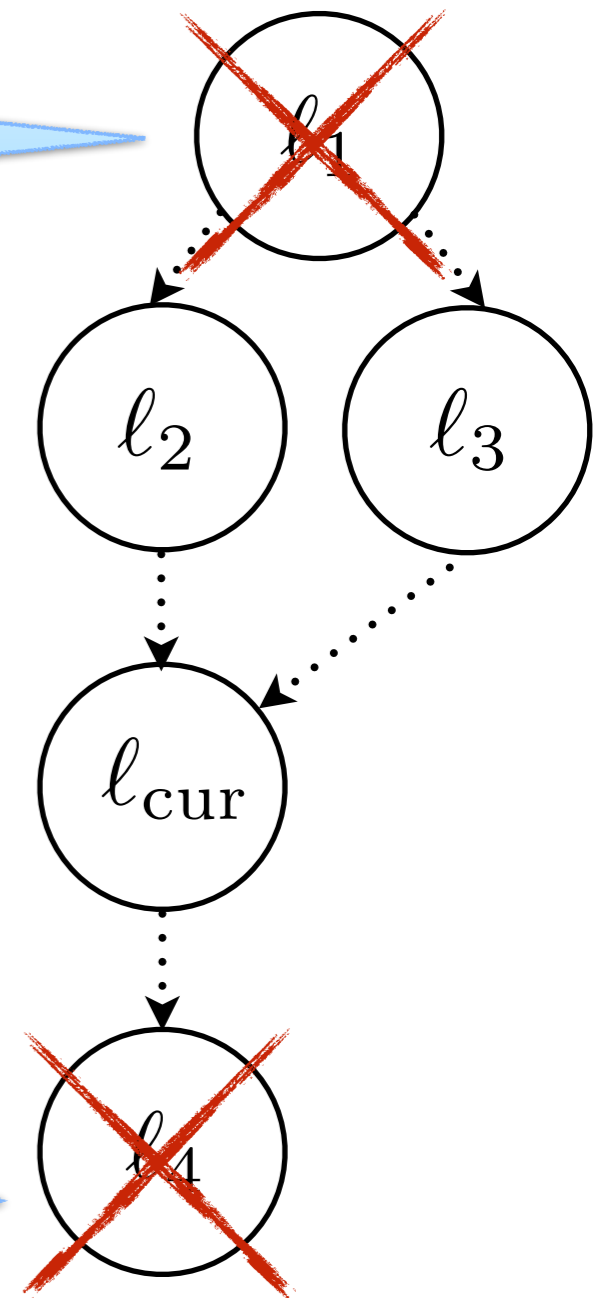
① $l_1 \xrightarrow{[c]} l_2$



Must visit another relevant location first.

Filter the set of data-relevant locations using control flow and the current program point

Not backward-reachable from current location



Jumping enables sparse, selective, on-the-fly
control-flow abstraction

1

$l_1 \dashrightarrow [c] \rightarrow l_2$

Jumping enables sparse, selective, on-the-fly **1**
control-flow abstraction

$l_1 \xrightarrow{[c]} l_2$

l_{cur}

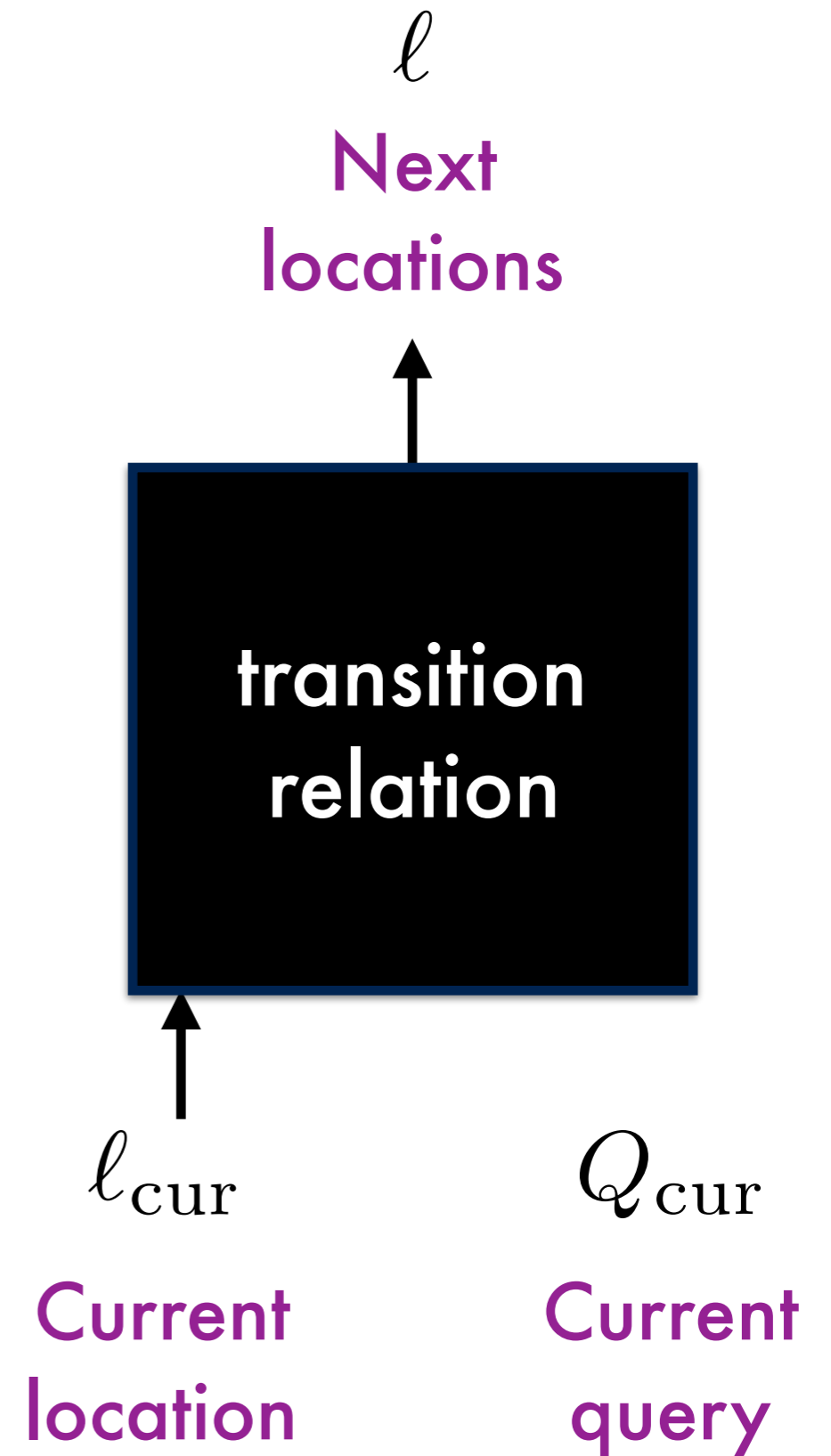
**Current
location**

Q_{cur}

**Current
query**

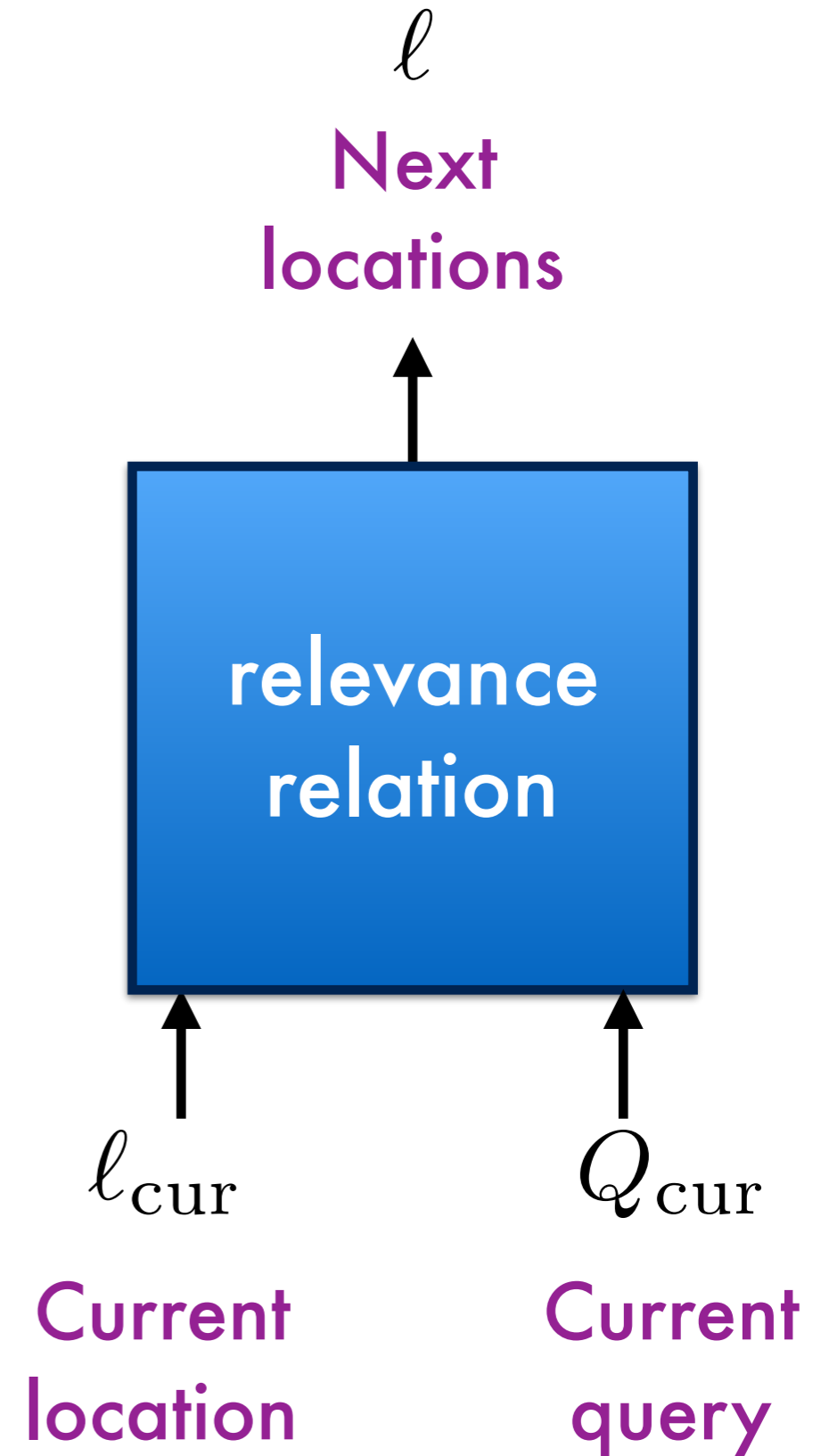
Jumping enables sparse, selective, on-the-fly ¹ control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$



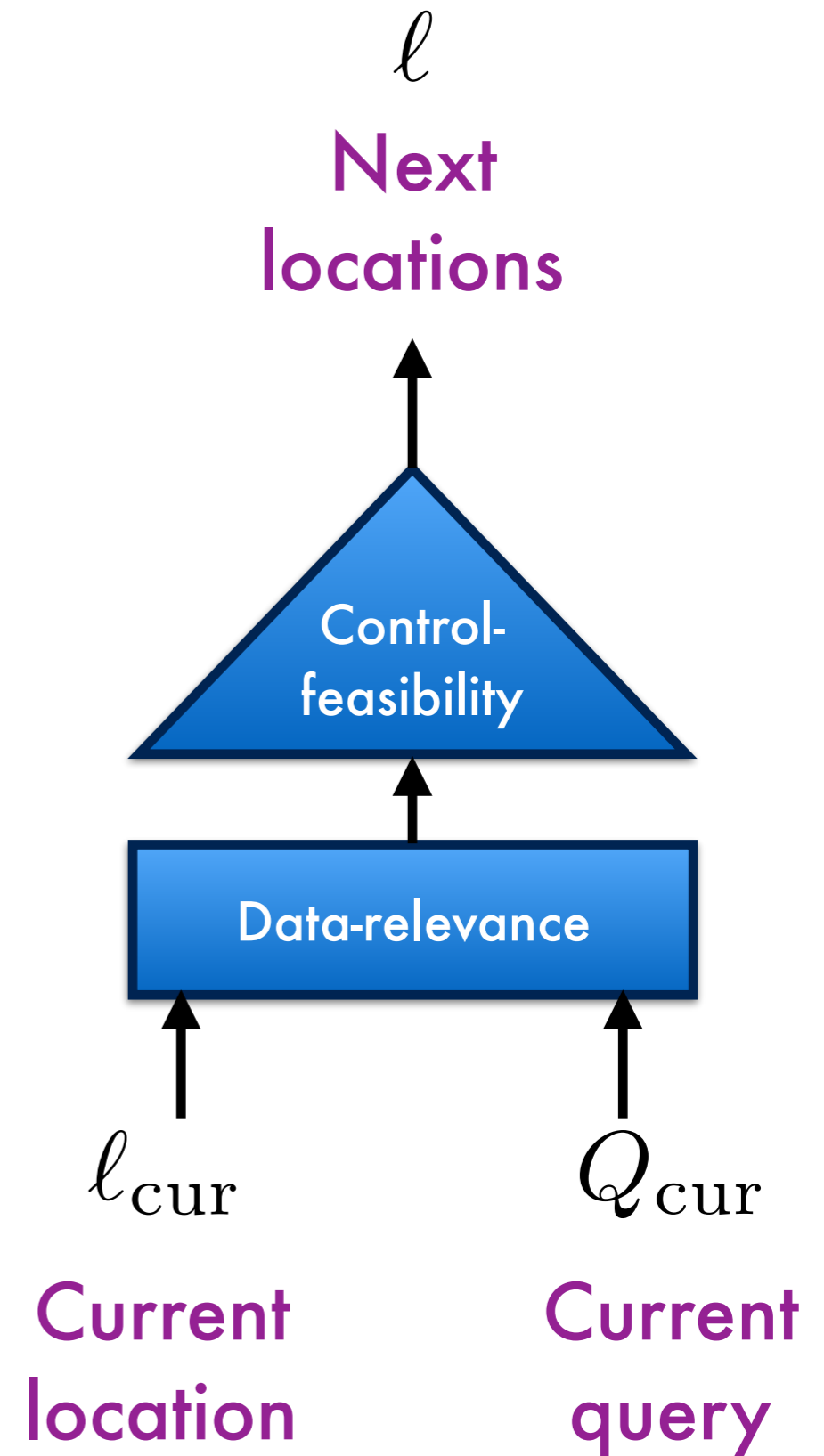
Jumping enables sparse, selective, on-the-fly ¹ control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$



Jumping enables sparse, selective, on-the-fly ¹ control-flow abstraction

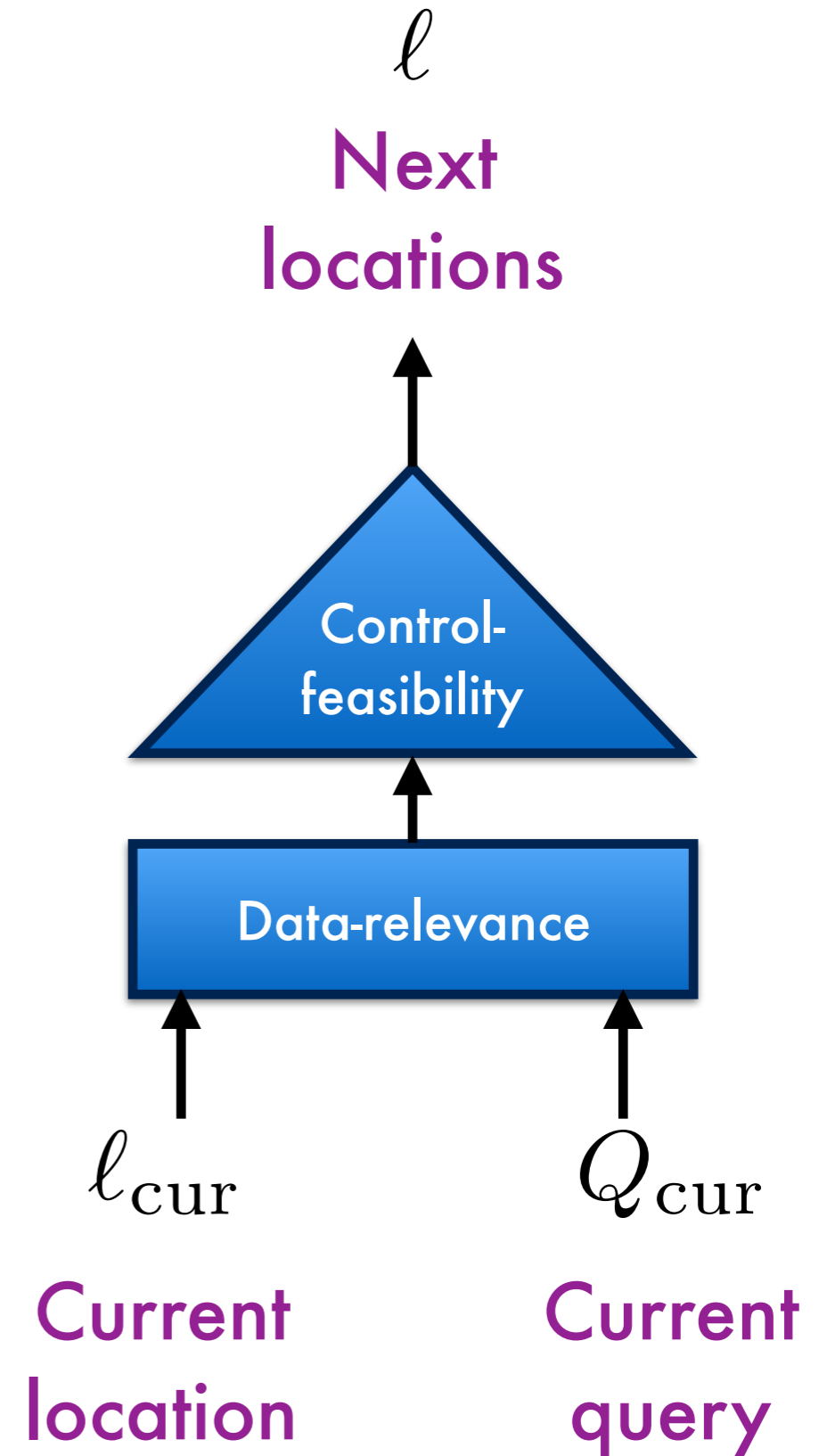
$$l_1 \xrightarrow{[c]} l_2$$



Jumping enables sparse, selective, on-the-fly ¹ control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow		
path		
context		

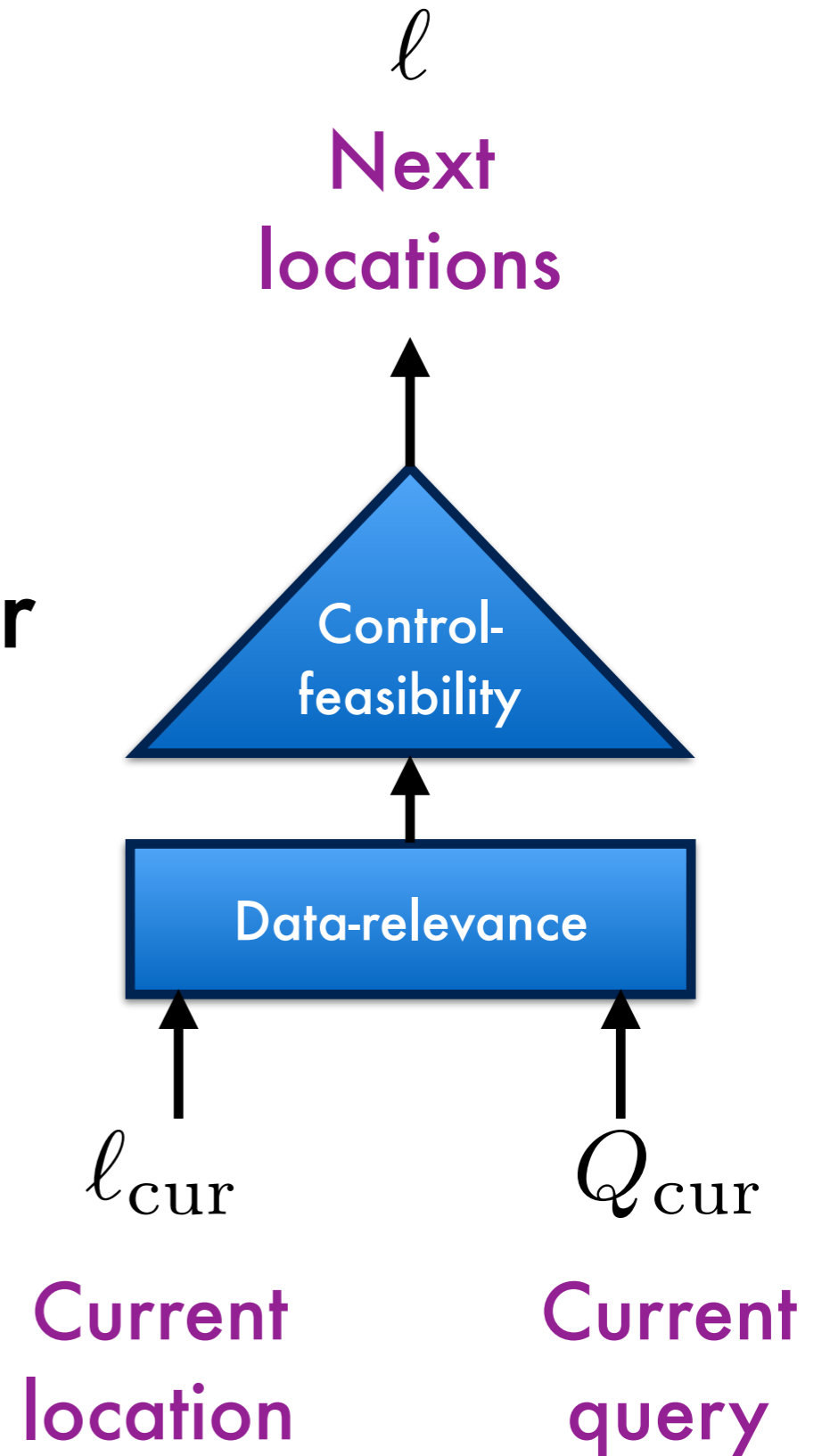


Jumping enables sparse, selective, on-the-fly ¹ control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow		
path		
context		

no filter



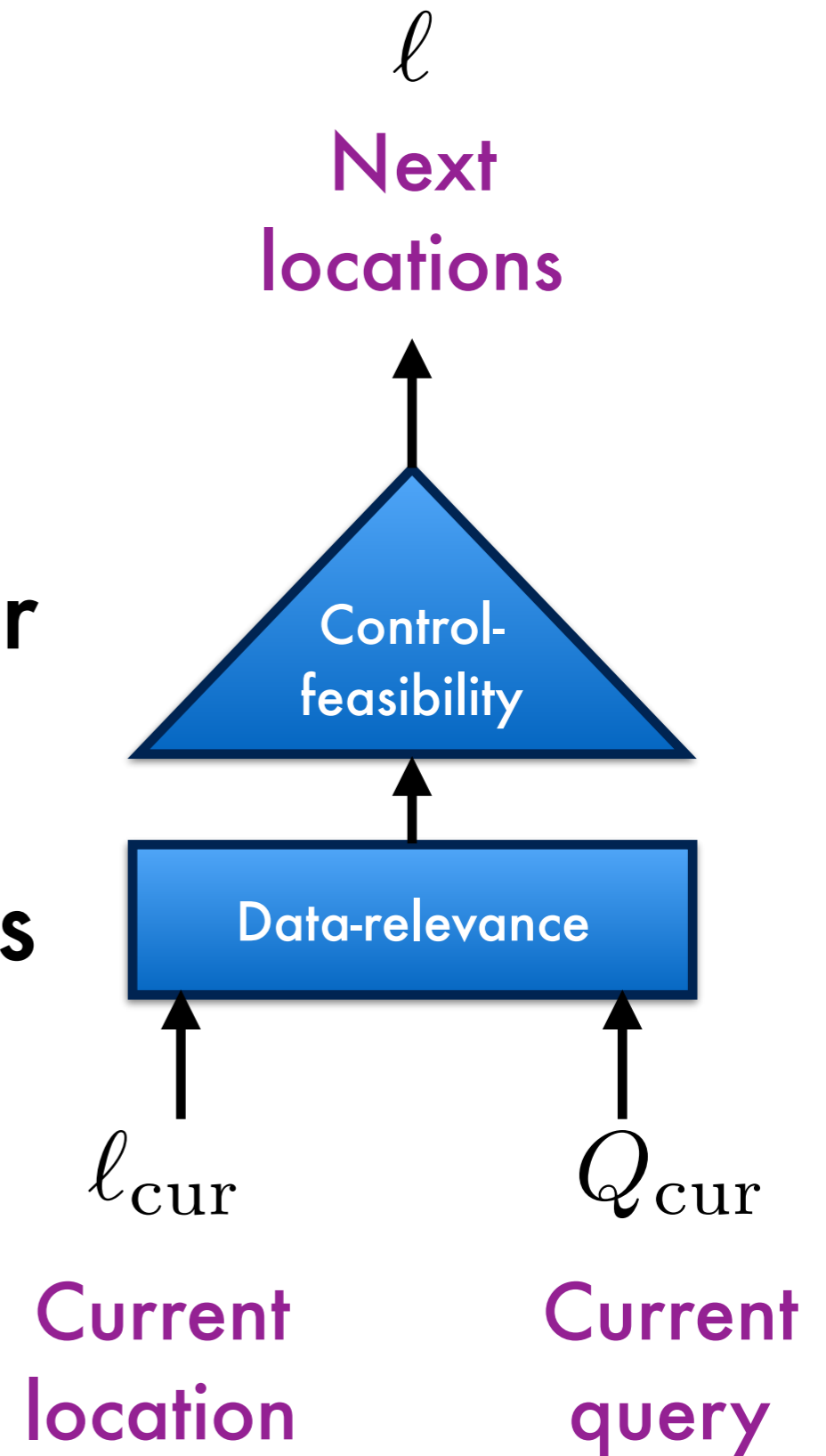
Jumping enables sparse, selective, on-the-fly control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow		
path		
context		

no filter

all locations



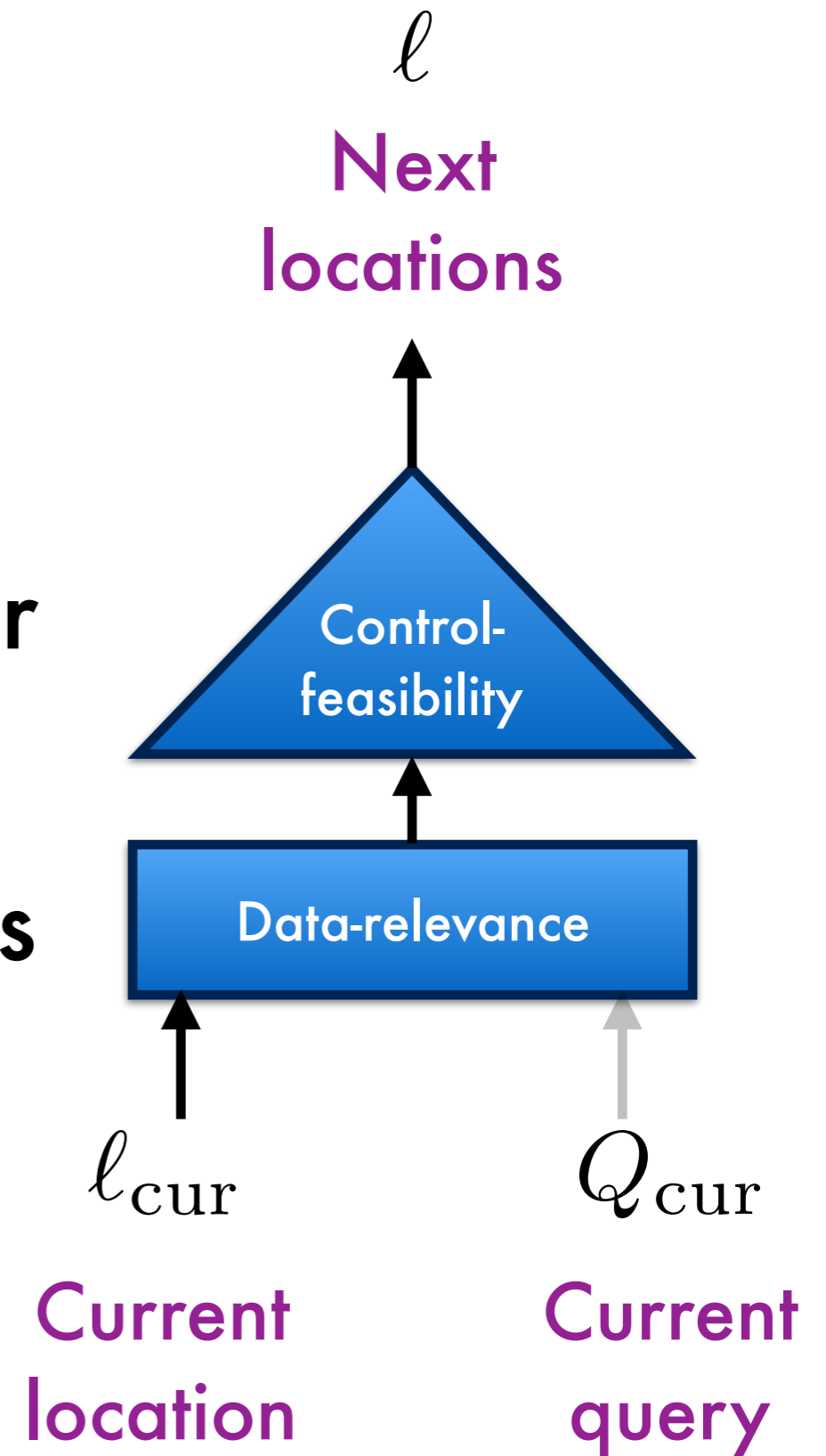
Jumping enables sparse, selective, on-the-fly control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow		
path		
context		

no filter

all locations



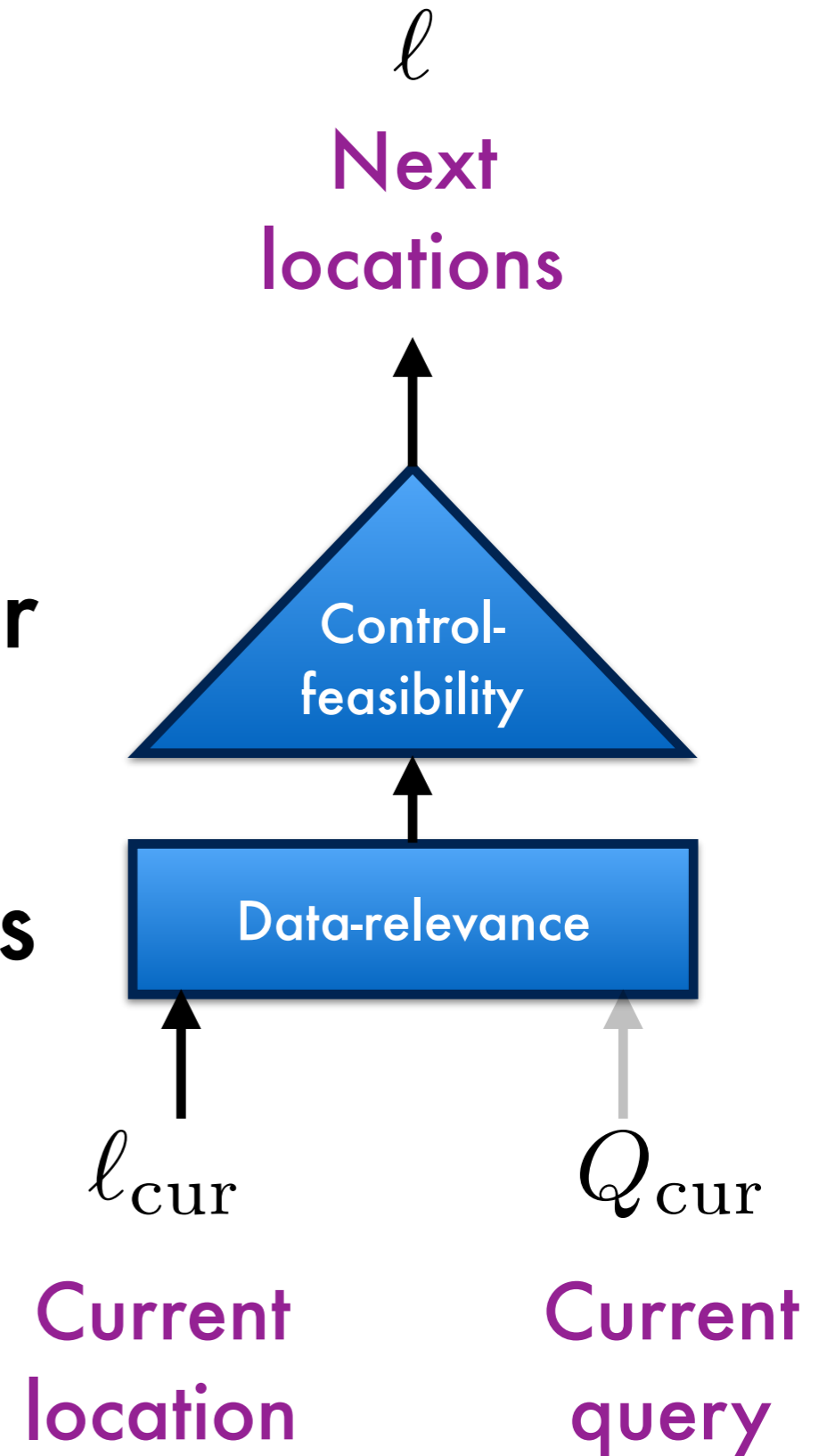
Jumping enables sparse, selective, on-the-fly control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow	✓	
path		
context		

no filter

all locations



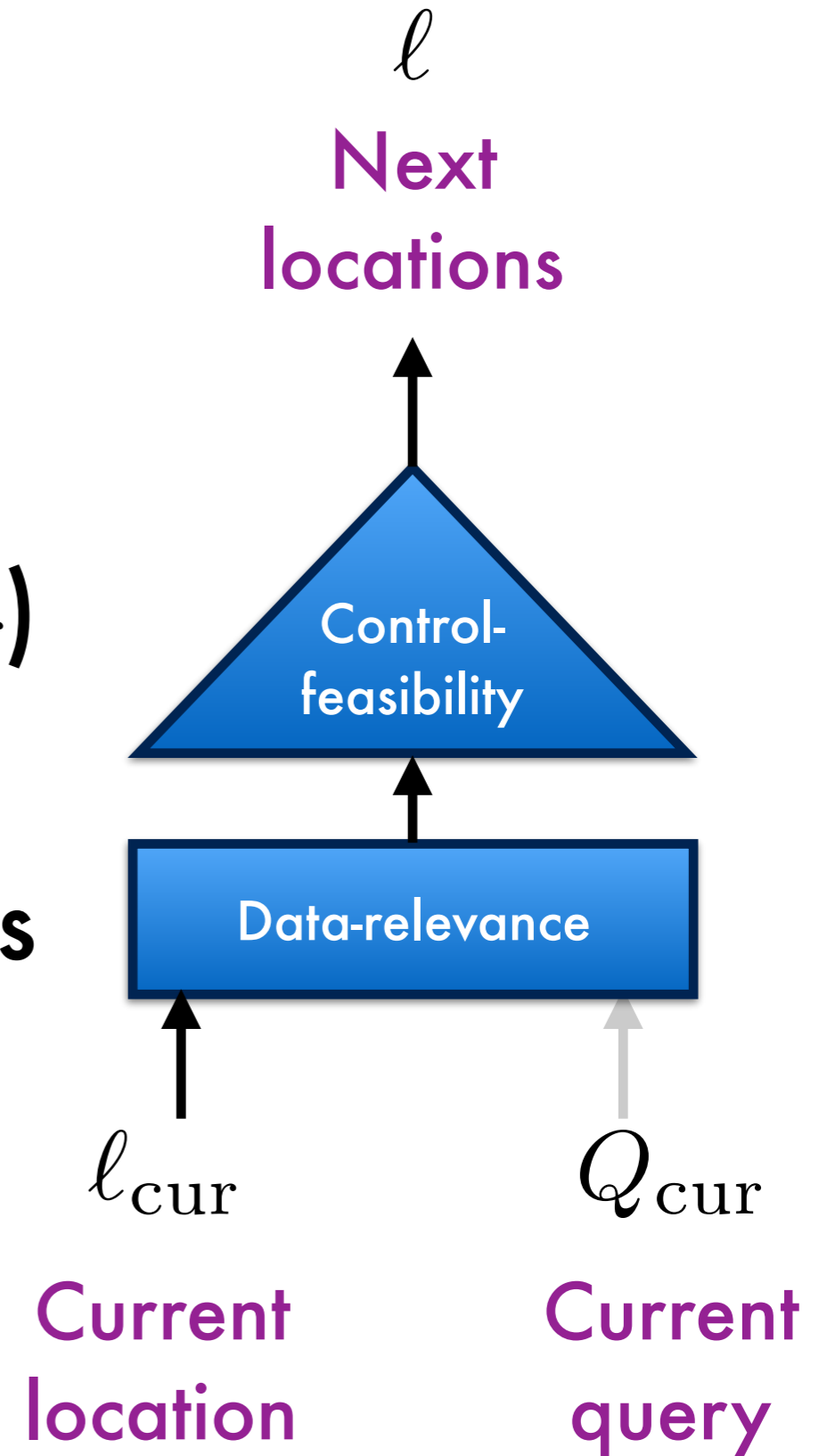
Jumping enables sparse, selective, on-the-fly control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow	✓	
path		
context		

only preds(l_{cur})

all locations



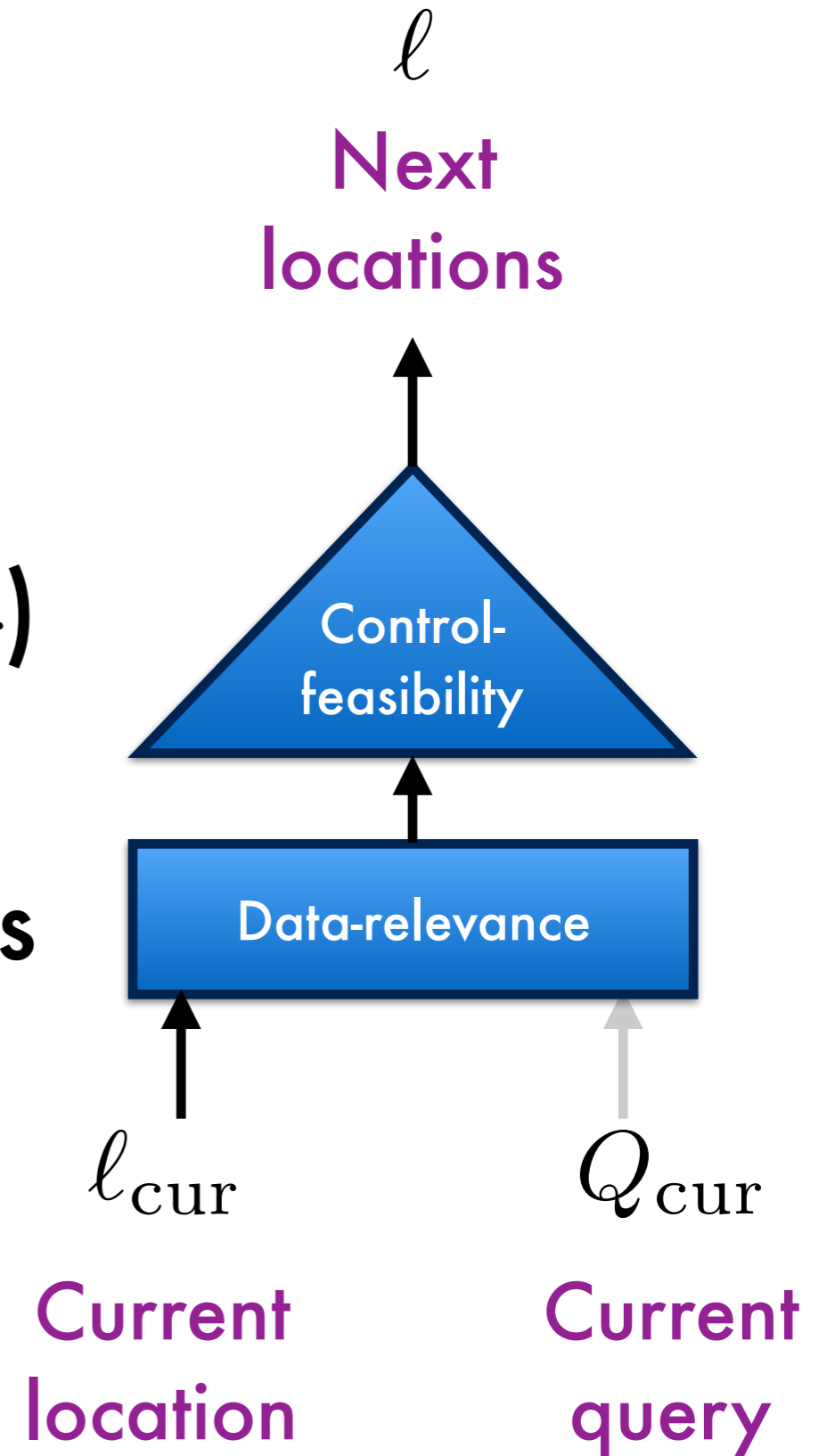
Jumping enables sparse, selective, on-the-fly control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow	✓	
path		
context		

only preds(l_{cur})

all locations



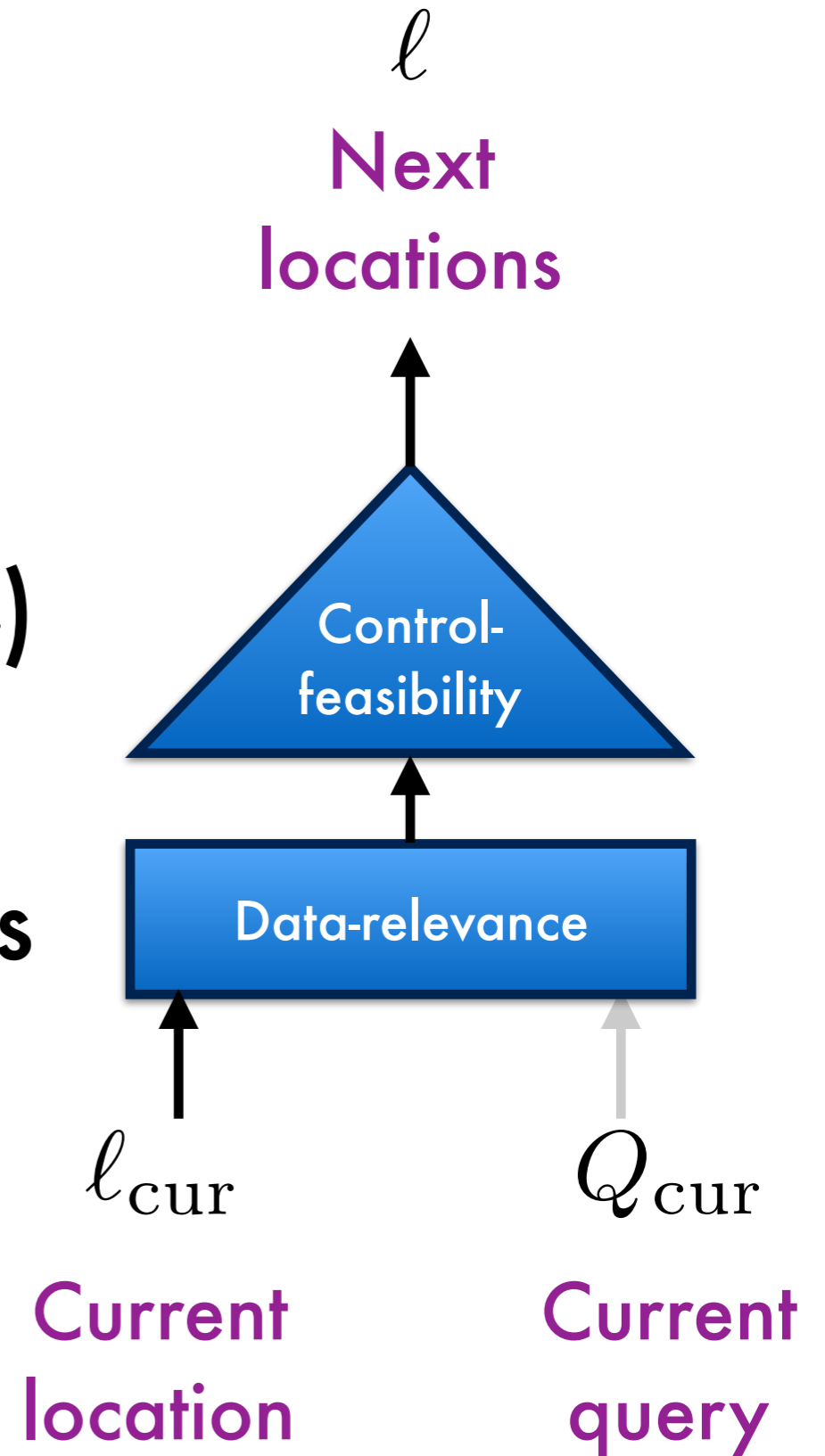
Jumping enables sparse, selective, on-the-fly ¹ control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow	✓	✓
path		✓
context	✓	

only preds(l_{cur})

all locations



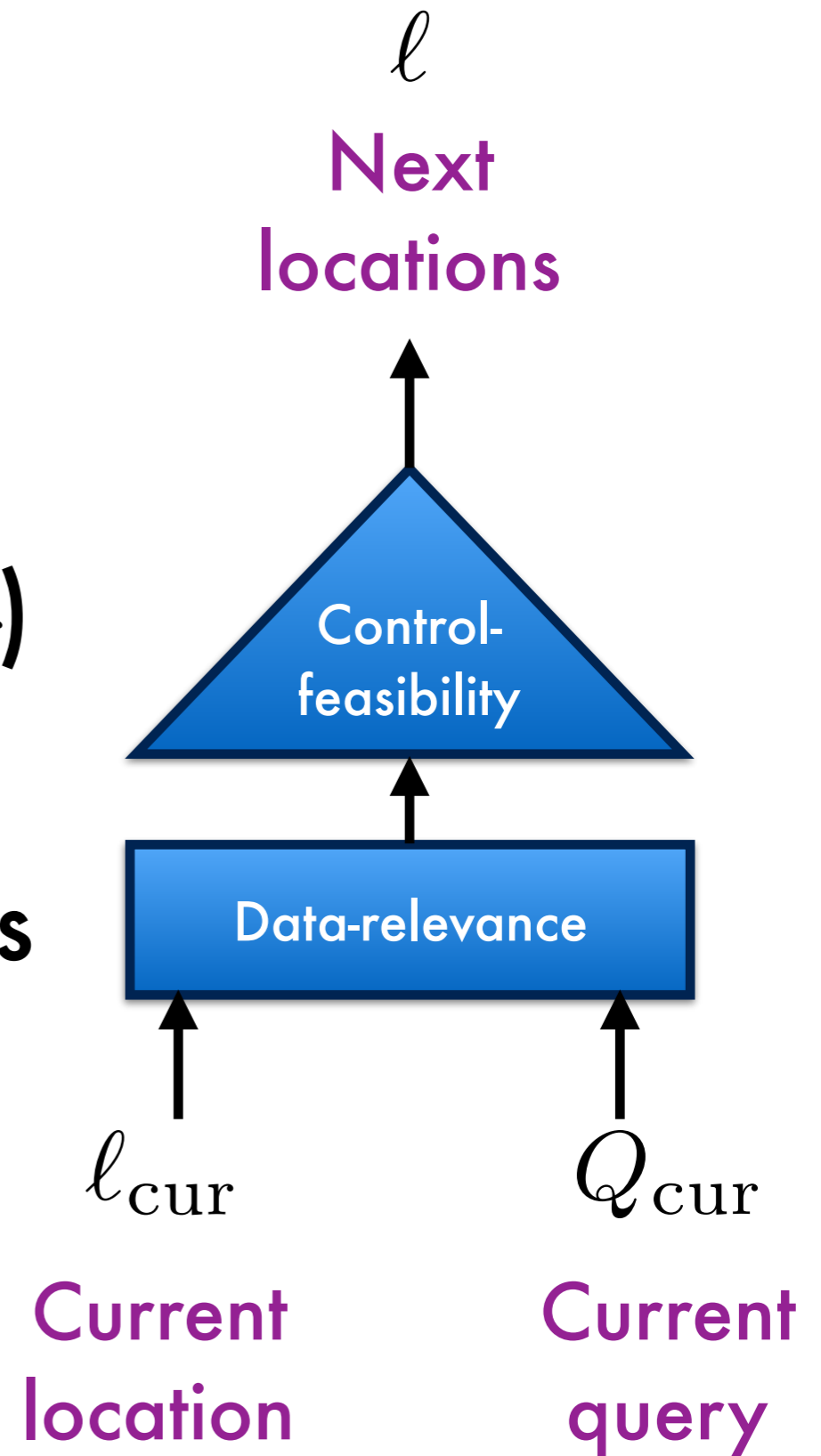
Jumping enables sparse, selective, on-the-fly ¹ control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow	✓	✓
path		✓
context	✓	

only feasible-preds(l_{cur}, Q_{cur})

all locations



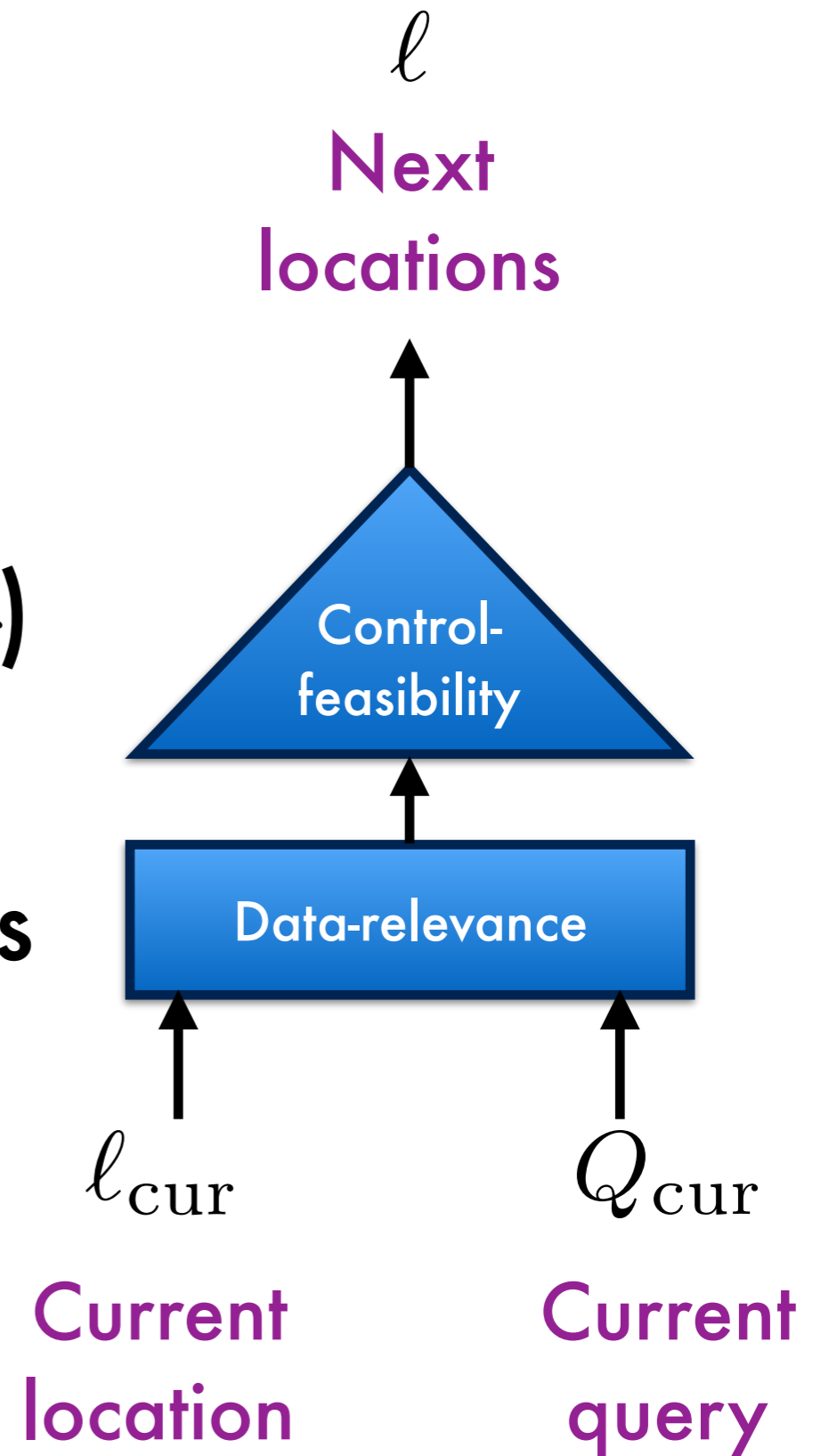
Jumping enables sparse, selective, on-the-fly ¹ control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow	✓	✓
path		✓
context	✓	

only feasible-preds(l_{cur}, Q_{cur})

all locations



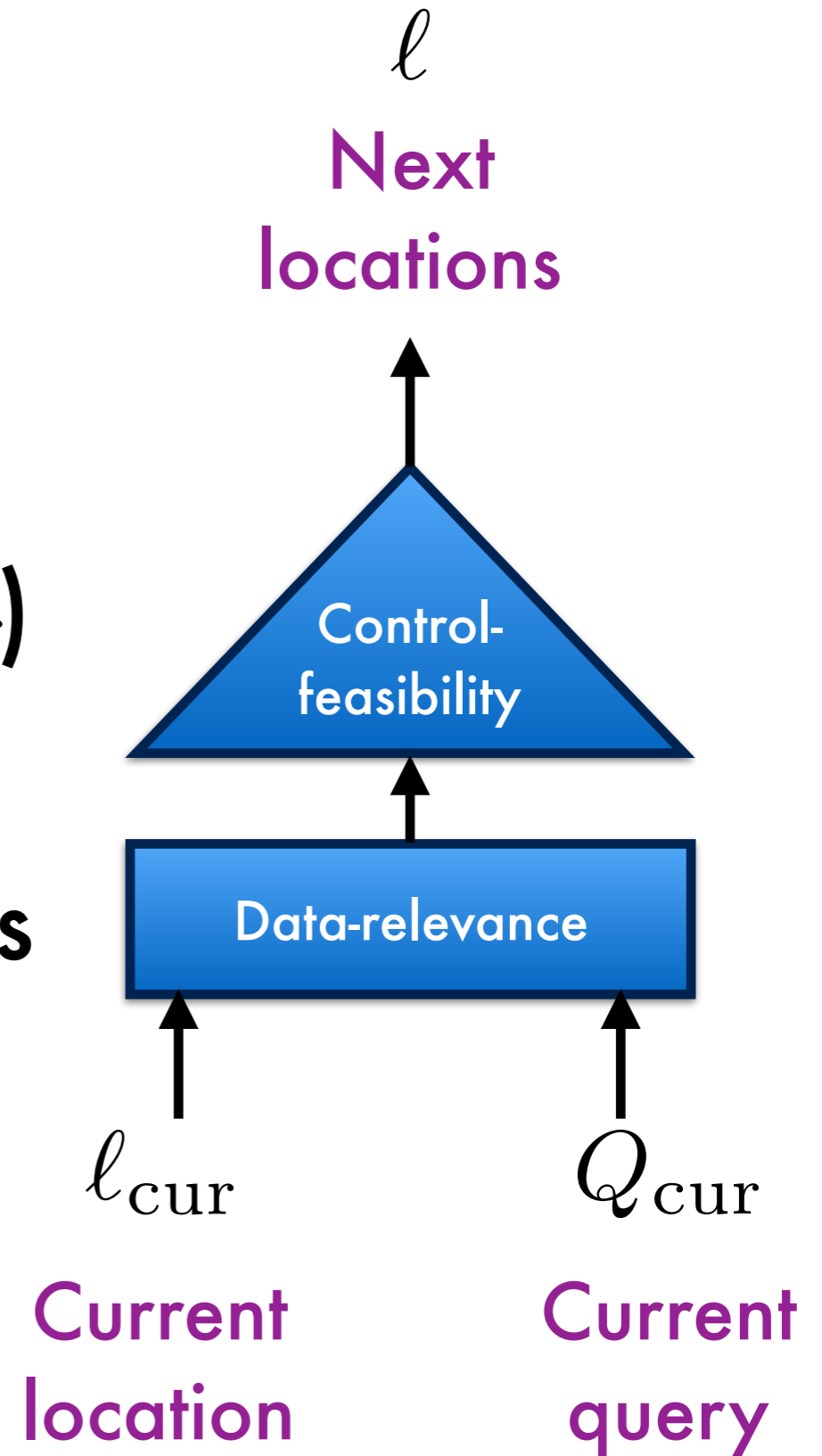
Jumping enables sparse, selective, on-the-fly ¹ control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow	✓	✓
path		✓
context	✓	✓

only feasible-preds(l_{cur}, Q_{cur})

all locations



Jumping enables sparse, selective, on-the-fly

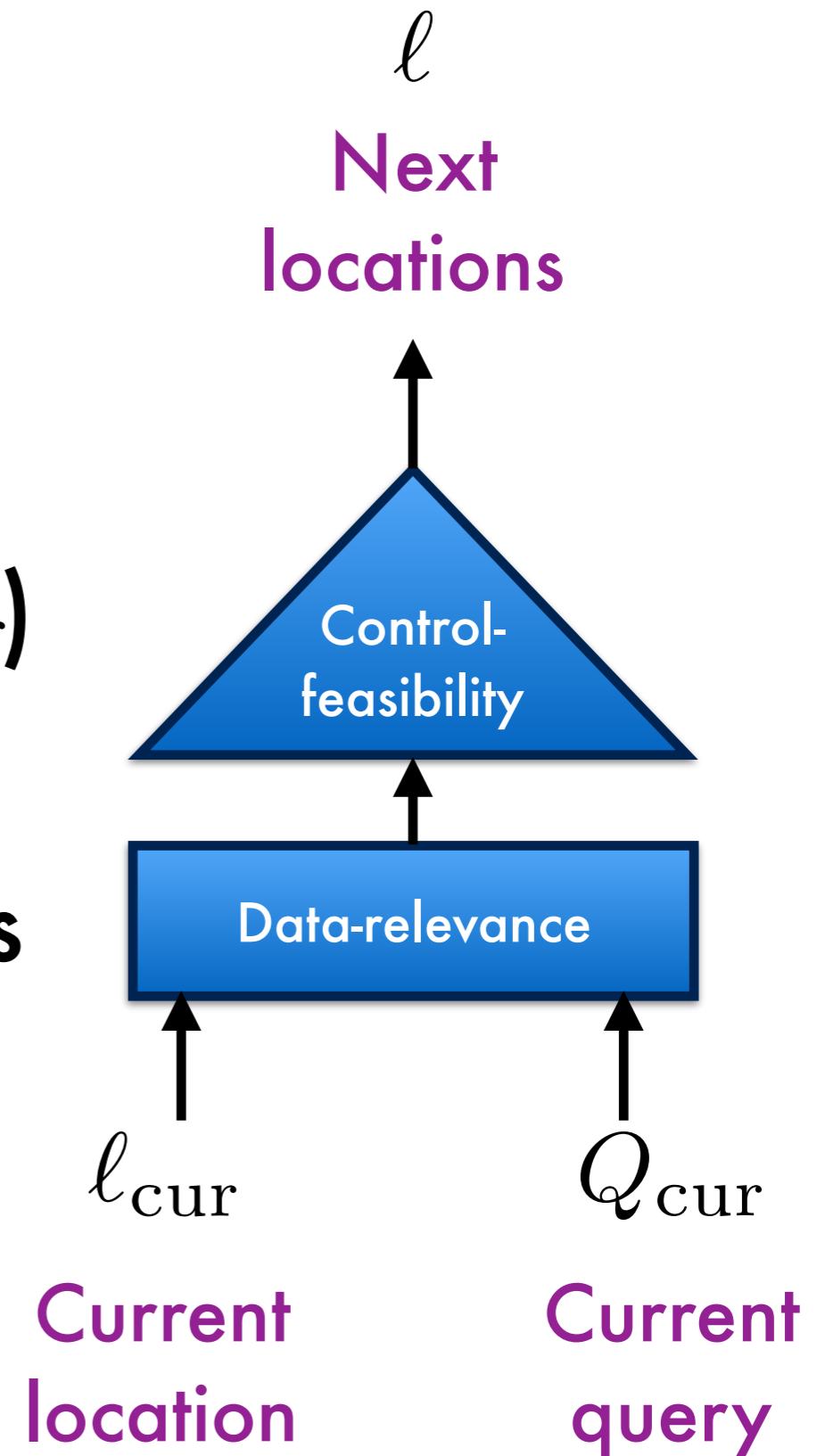
control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow	✓	✓
path		✓
context	✓	✓

only feasible-preds(l_{cur}, Q_{cur})

all locations, except assumes



Jumping enables sparse, selective, on-the-fly

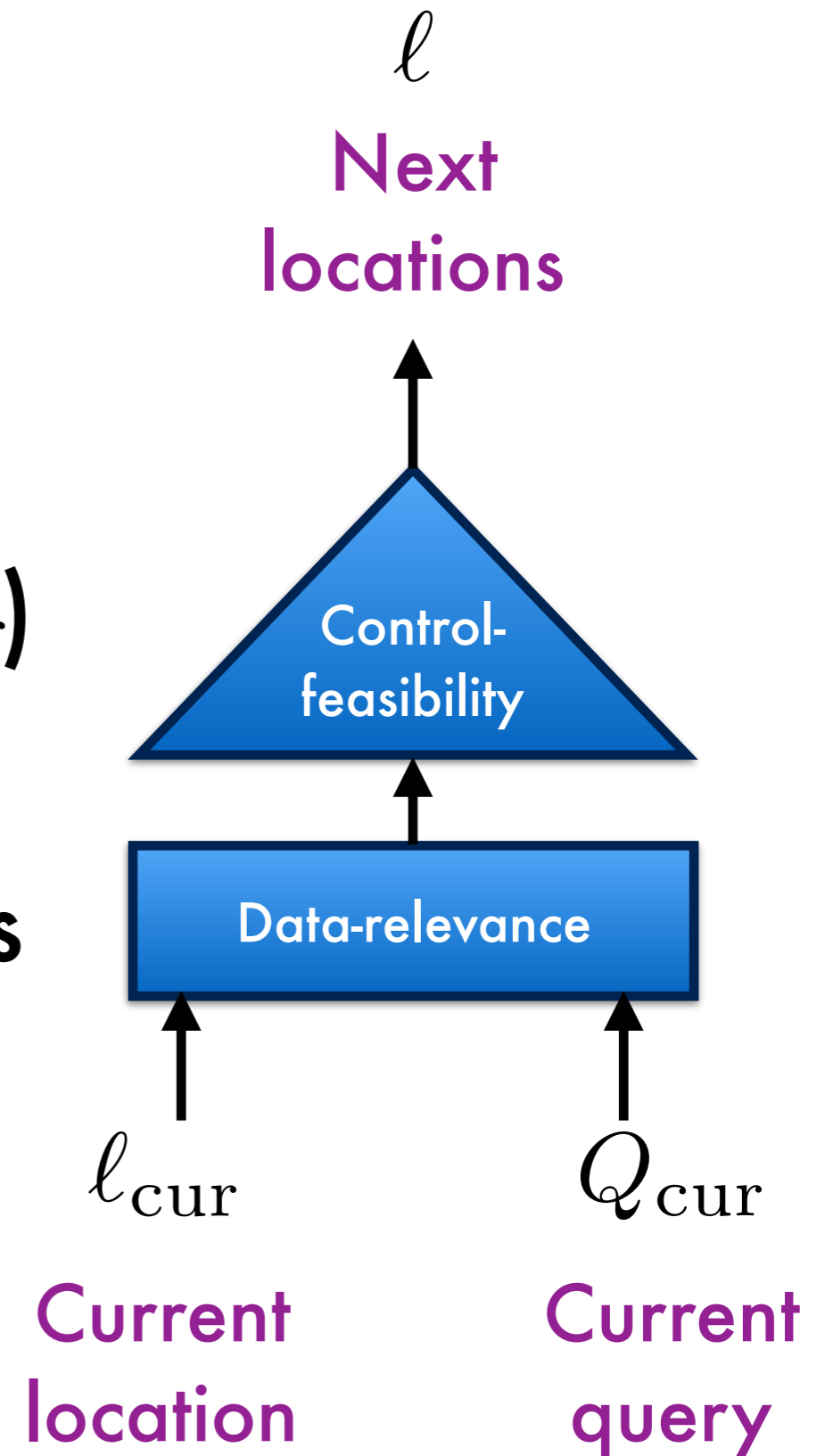
control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow	✓	✓
path		✓
context	✓	✓

only feasible-preds(l_{cur}, Q_{cur})

all locations, except assumes



Jumping enables sparse, selective, on-the-fly

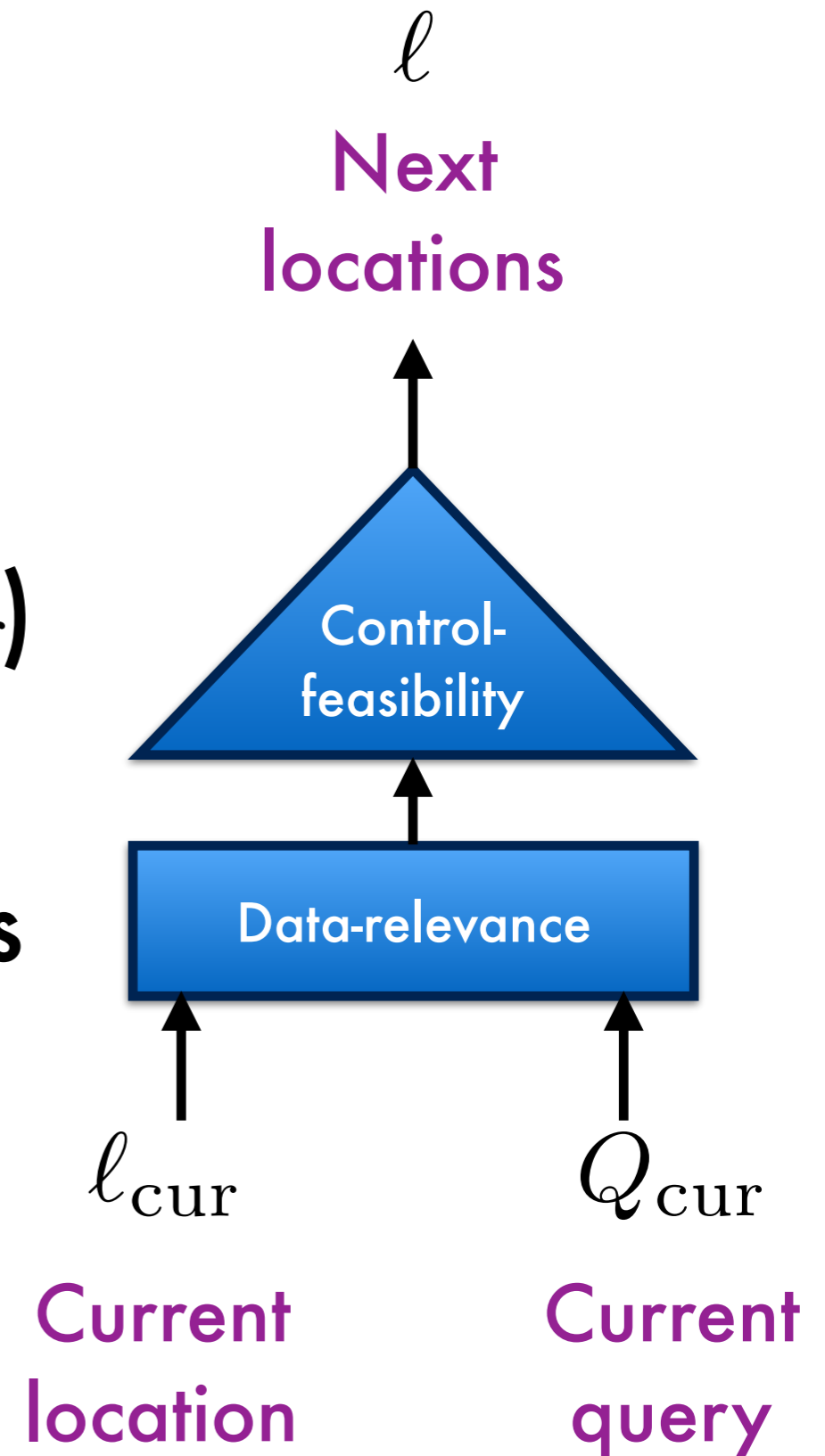
control-flow abstraction

$$l_1 \xrightarrow{[c]} l_2$$

Precision	-insensitive	-sensitive
flow	✓	✓
path	✓	✓
context	✓	✓

only feasible-preds(l_{cur}, Q_{cur})

all locations, except assumes



Jumping enables sparse, selective, on-the-fly

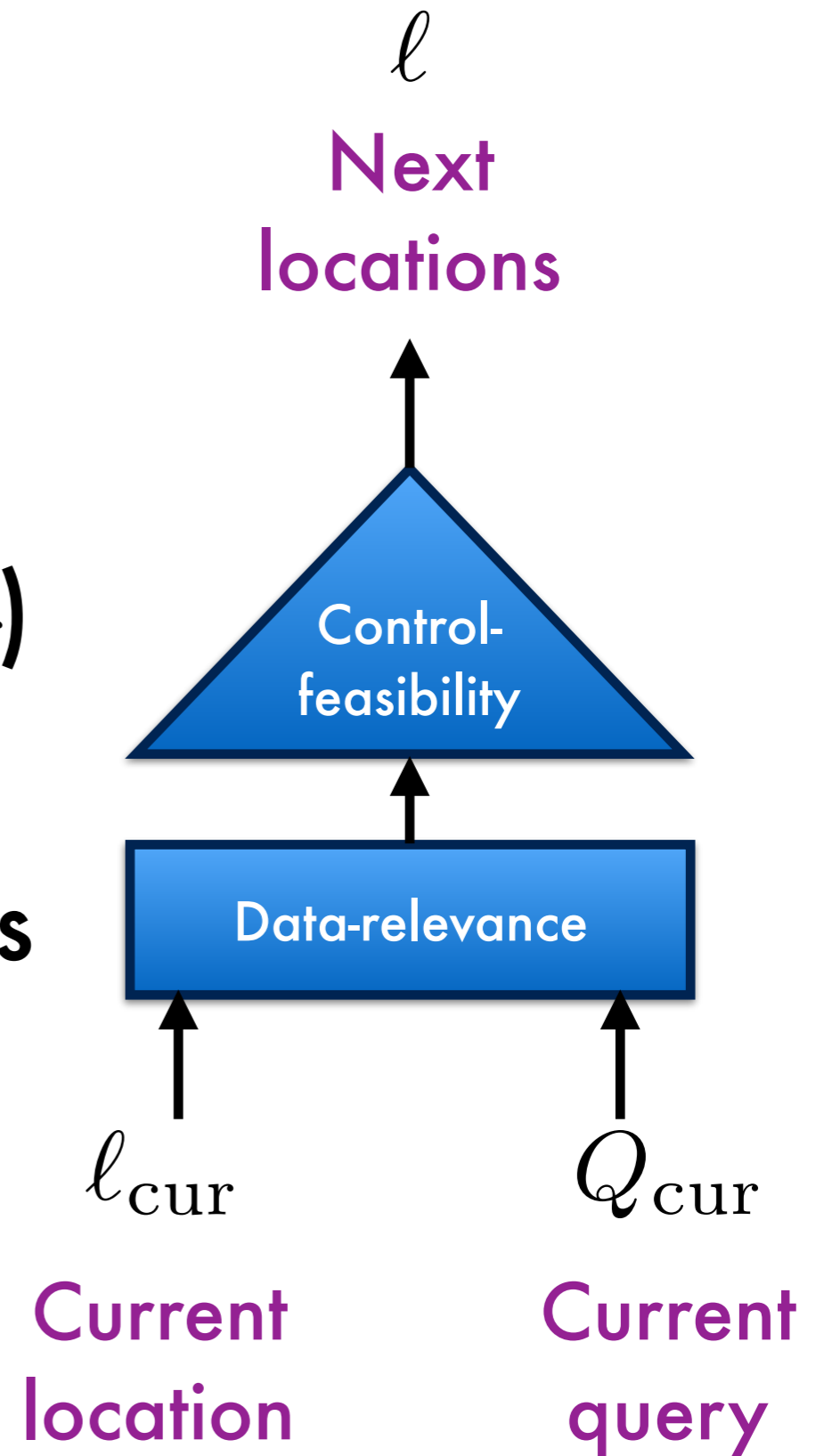
control-flow abstraction

1

$l_1 \xrightarrow{[c]} l_2$

only feasible-preds(l_{cur}, Q_{cur})

all locations

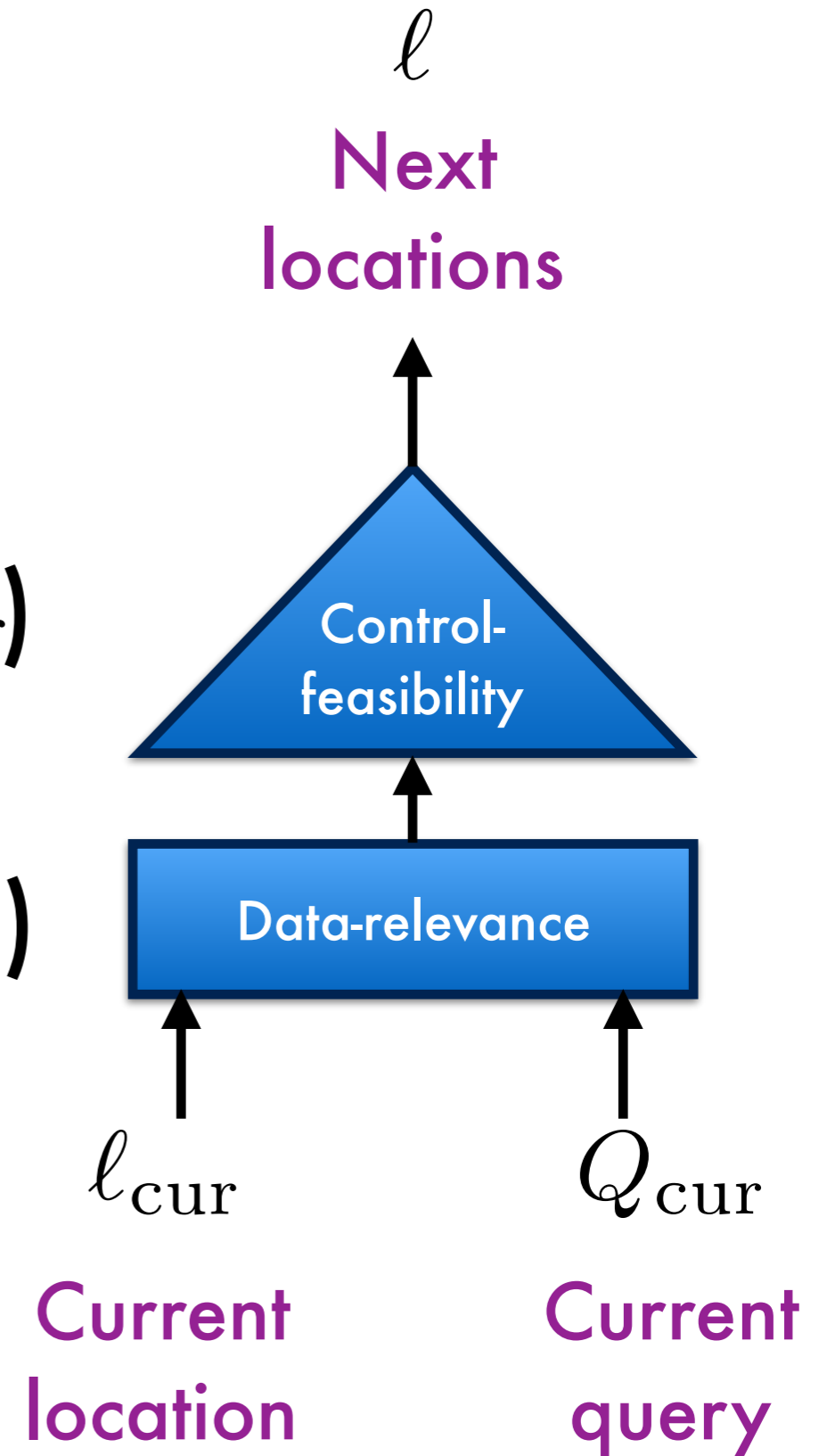


Jumping enables sparse, selective, on-the-fly **control-flow abstraction**

$$l_1 \xrightarrow{[c]} l_2$$

only feasible-preds(l_{cur}, Q_{cur})

all mods(Q_{cur})



Jumping enables sparse, selective, on-the-fly

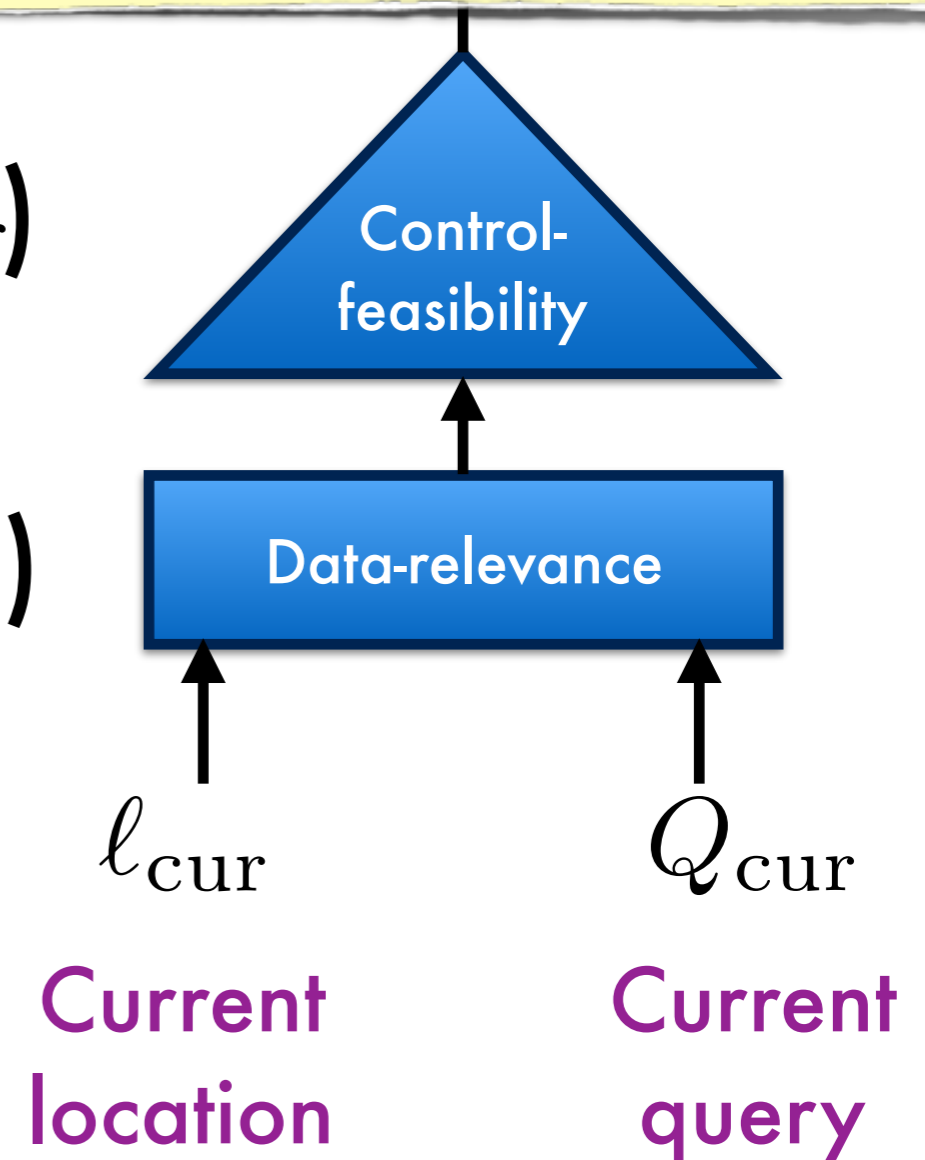
control-flow abstraction

① $l_1 \xrightarrow{[c]} l_2$

Be **sparse** by using data relevance with desired control-flow abstraction

only feasible-preds(l_{cur}, Q_{cur})

all mods(Q_{cur})



Jumping enables sparse, selective, on-the-fly

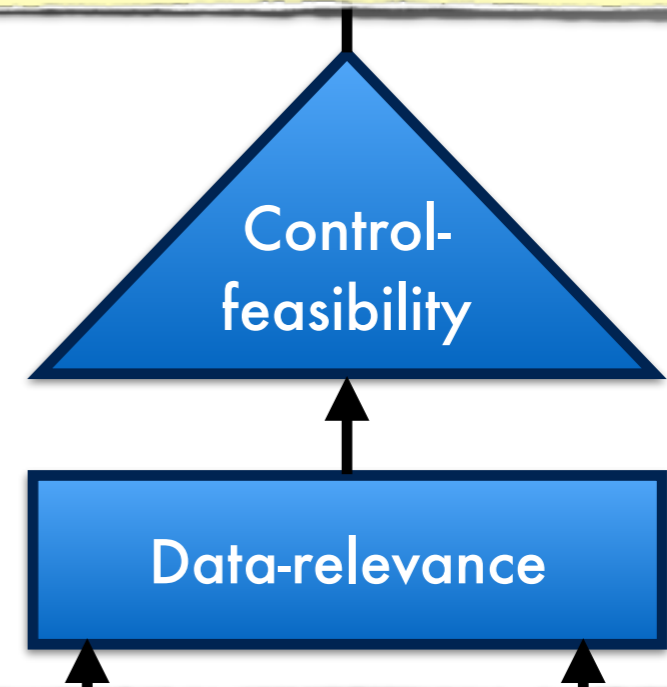
control-flow abstraction

① $l_1 \xrightarrow{[c]} l_2$

Be **sparse** by using data relevance with desired control-flow abstraction

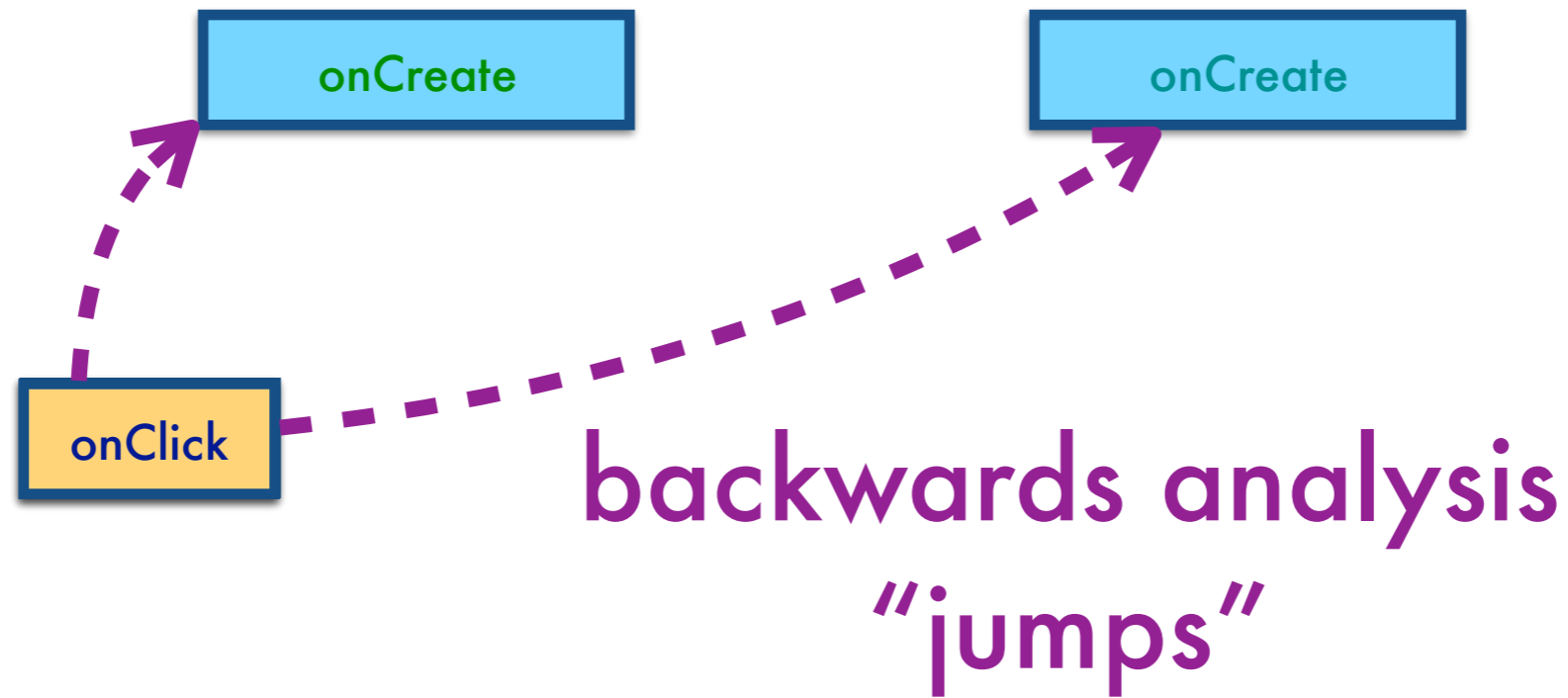
only feasible-preds(l_{cur}, Q_{cur})

all mods(Q_{cur})



Be **selective** by varying the relevance relation at each an analysis step

Contributions

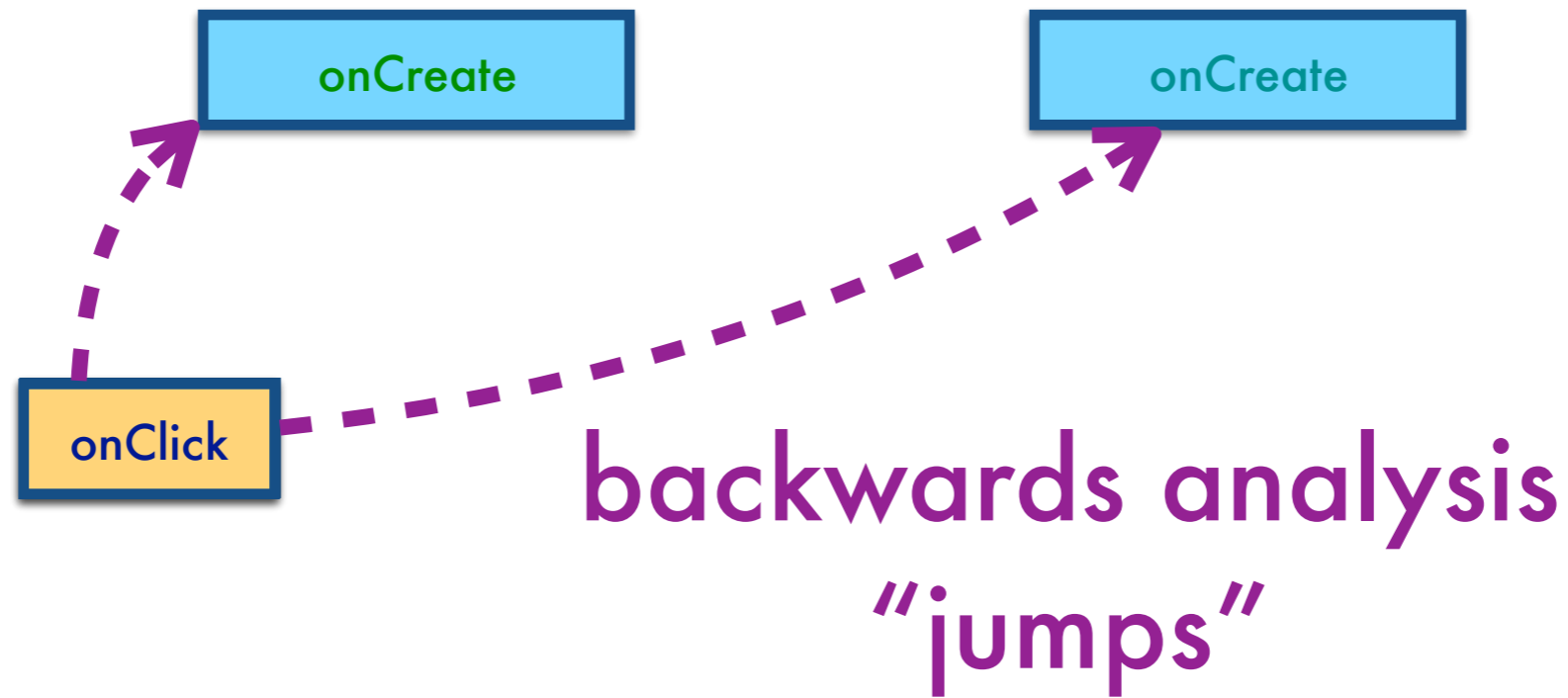


1 $l_1 \xrightarrow{[c]} l_2$ Framework for jumping analyses



Applied to Android events

Contributions

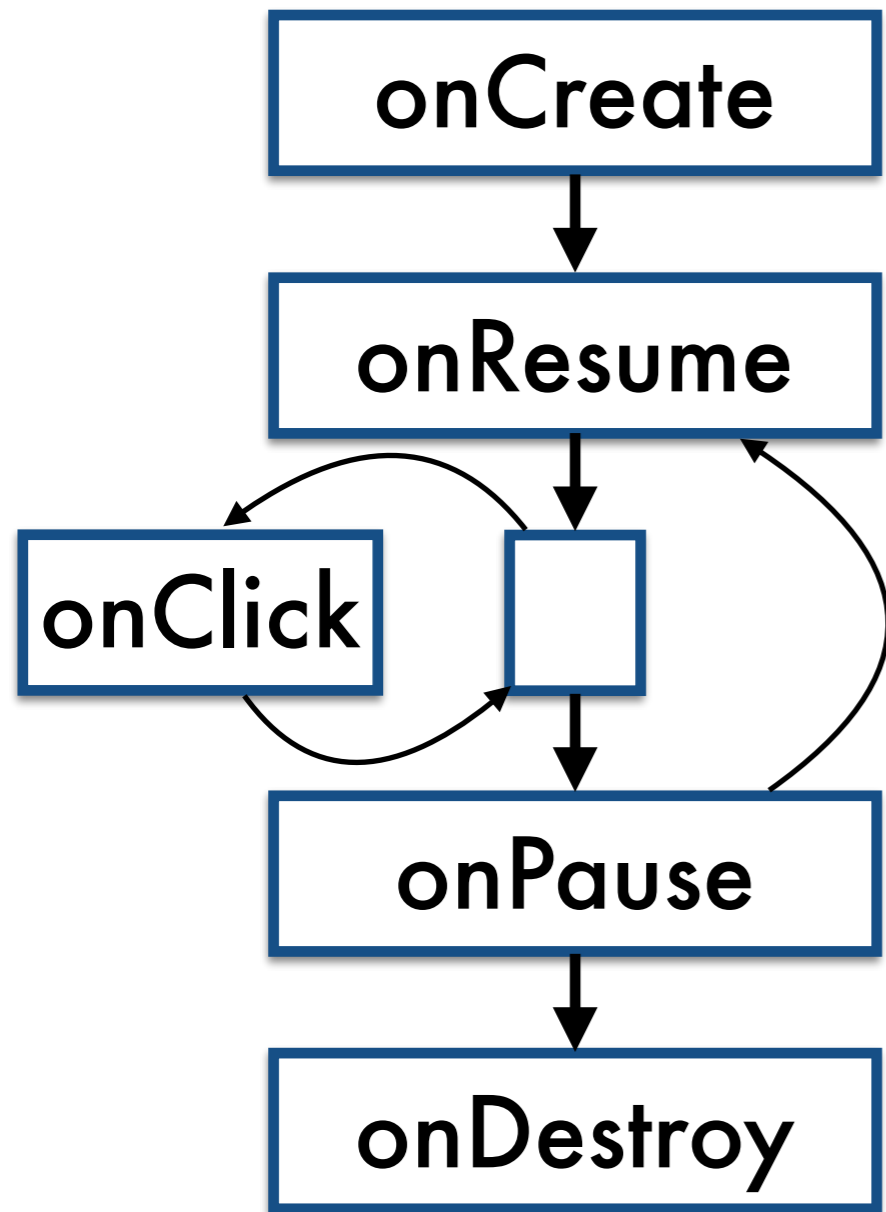


1 $l_1 \xrightarrow{[c]} l_2$ Framework for jumping analyses



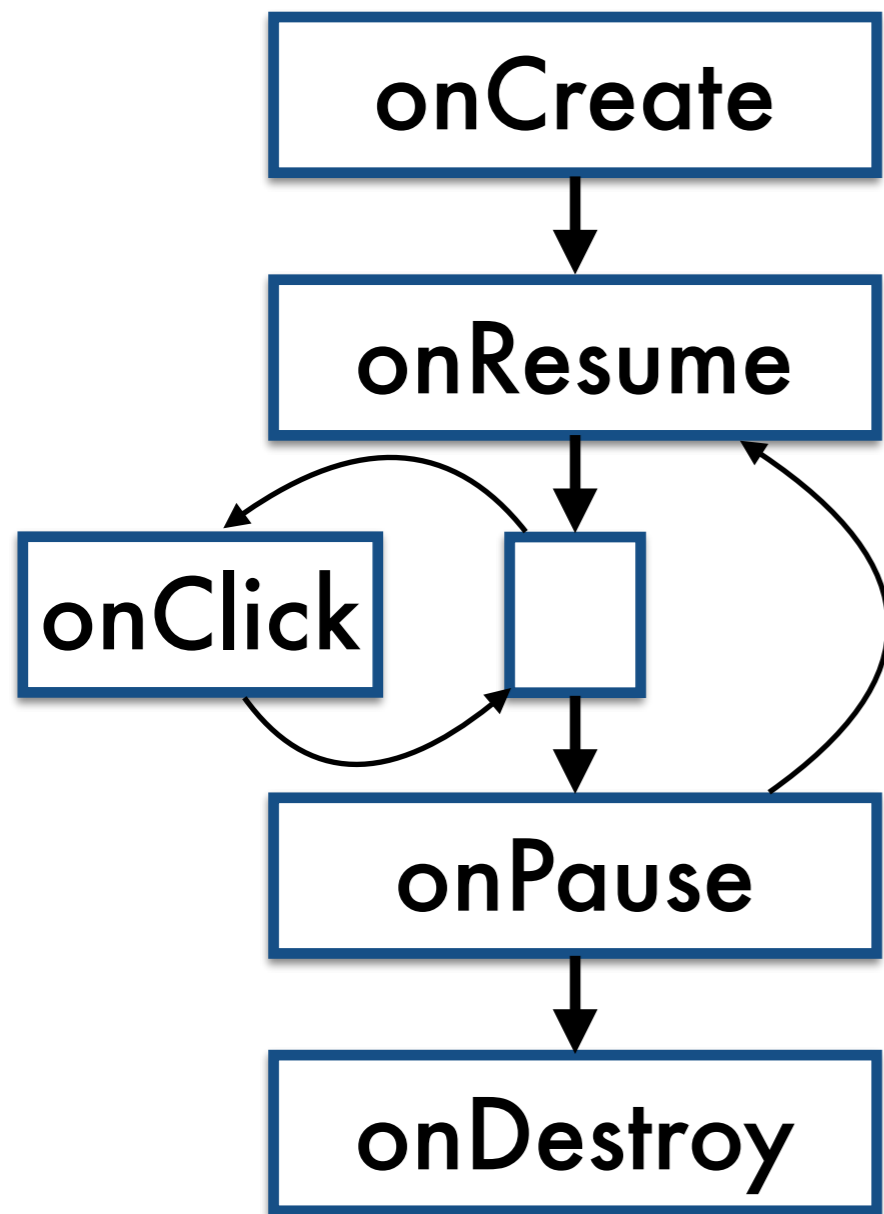
Applied to Android events

Formalizing lifecycle graphs



Formalizing lifecycle graphs

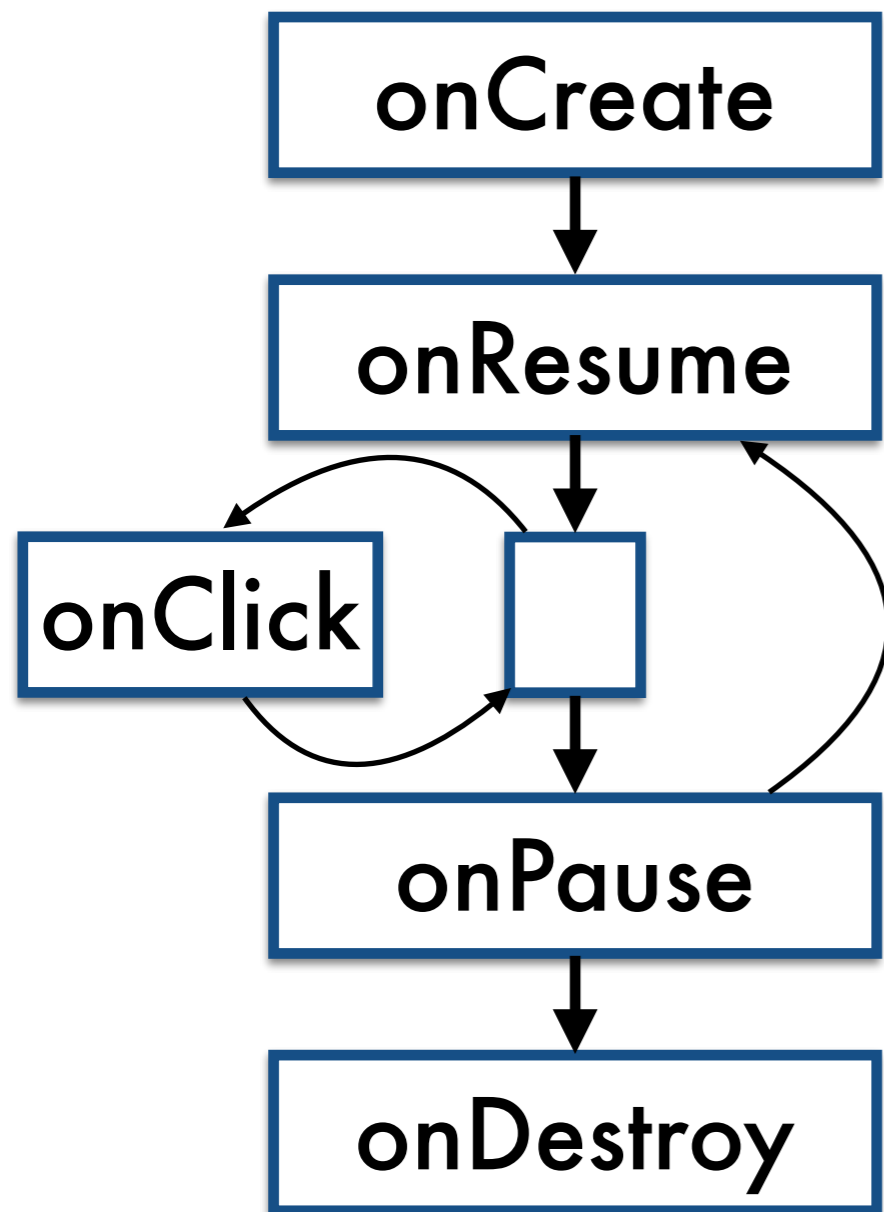
2



“Lifecycle **automata**”
accepts the concrete
trace of callbacks
projected onto its
transitions

Formalizing lifecycle graphs

2



“Lifecycle **automata**”
accepts the concrete
trace of callbacks
projected onto its
transitions

Further scalability challenge:
Lifecycle spec per class but
analysis applies per object

A jumping policy for Android analysis



A jumping policy for Android analysis



**Within an event-callback (*intra-event*),
follow predecessor transitions**

no jumping, path/context-sensitive

A jumping policy for Android analysis



Within an event-callback (**intra-event**),
follow predecessor transitions

no jumping, path/context-sensitive

Between event-callbacks (**inter-event**), jump
using lifecycle graphs for control-feasibility
filtering

**Hypothesis: Jumping is an
effective approach to
path-sensitive, inter-event
analysis**

Setup: Proving dereferences safe



Setup: Proving dereferences safe

10 open source Android apps

3,000 to 55,000 lines of code

10 to 100 components

120 to 1,320 callbacks



Setup: Proving dereferences safe

10 open source Android apps

3,000 to 55,000 lines of code

10 to 100 components

120 to 1,320 callbacks



Event product graph would
have 10^{10} to 10^{11} nodes
(with one instance per class)

Setup: Proving dereferences safe

10 open source Android apps

3,000 to 55,000 lines of code

10 to 100 components

120 to 1,320 callbacks



Setup: Proving dereferences safe



10 open source Android apps

3,000 to 55,000 lines of code

10 to 100 components

120 to 1,320 callbacks

Compared 3 analyses

Nit: flow-insensitive

Thresher: no jumping

Hopper: jumping

Is jumping effective?

	KLOC	Deref	Nit	Thr	Hop (Impr %)
drupaleditor	3	928	679	179	72 (60)
npr	5	829	617	181	51 (72)
duckduckgo	11	1969	1341	518	143 (72)
lastfm	13	4840	3528	954	477 (50)
github	19	3603	2520	601	290 (52)
seriesguide	32	8184	5438	986	625 (37)
connectbot	33	2190	1562	316	74 (77)
textsecure	38	5921	3643	698	330 (53)
k-9	55	19032	11968	3104	1988 (36)
wordpress	57	15066	9775	2431	1362 (44)
Total	266	62562	41071	9968	5412 (54)

Is jumping effective?

	KLOC	Deref	Nit	Thr Hop	(Impr %)
drupaleditor	3	928	679	179	72 (60)
npr	5	829	617	181	51 (72)
duckduckgo	11	1969	1341	518	143 (72)
lastfm	13	4840	3528	954	477 (50)
github	19	3603	2520	601	290 (52)
seriesguide	32	8184	5438	986	625 (37)
connectbot	33	2190	1562	316	74 (77)
trac	38	5921	3643	698	330 (53)
wordpress	57	19032	11968	3104	1988 (36)
Total	266	62562	41071	9968	5412 (54)

Many derefs

Is jumping effective?

unproven derefs

	KLOC	Deref	Nit	Thr	Hop (Impr %)
drupaleditor	3	928	679	179	72 (60)
npr	5	829	617	181	51 (72)
duckduckgo	11	1969	1341	518	143 (72)
lastfm	13	4840	3528	954	477 (50)
github	19	3603	2520	601	290 (52)
seriesguide	32	8184	5438	986	625 (37)
connectbot	33	2190	1562	316	74 (77)
textsecure	38	5921	3643	698	330 (53)
k-9	55	19032	11968	3104	1988 (36)
wordpress	57	15066	9775	2431	1362 (44)
Total	266	62562	41071	9968	5412 (54)

Is jumping effective?

Flow-insensitive

No jumping

Jumping

unproven derefs

	KLOC	Deref	Nit	Thr Hop (Impr %)	
drupaleditor	3	928	679	179	72 (60)
npr	5	829	617	181	51 (72)
duckduckgo	11	1969	1341	518	143 (72)
lastfm	13	4840	3528	954	477 (50)
github	19	3603	2520	601	290 (52)
seriesguide	32	8184	5438	986	625 (37)
connectbot	33	2190	1562	316	74 (77)
textsecure	38	5921	3643	698	330 (53)
k-9	55	19032	11968	3104	1988 (36)
wordpress	57	15066	9775	2431	1362 (44)
Total	266	62562	41071	9968	5412 (54)

Is jumping effective?

unproven derefs

	KLOC	Deref	Nit	Thr	Hop (Impr %)
drupaleditor	3	928	679	179	72 (60)
npr	5	829	617	181	51 (72)
duckduckgo	11	1969	1341	518	143 (72)
lastfm	13	4840	3528	954	477 (50)
github	19	3603	2520	601	290 (52)
seriesguide	32	8184	5438	986	625 (37)
connectbot	33	2190	1562	316	74 (77)
textsecure	38	5921	3643	698	330 (53)
k-9	55	19032	11968	3104	1988 (36)
wordpress	57	15066	9775	2431	1362 (44)
Total	266	62562	41071	9968	5412 (54)

Is jumping effective?

unproven derefs

	KLOC	Deref	Nit	Thr	Hop (Impr %)
drupaleditor	3	928	679	179	72 (60)
npr	5	829	617	181	51 (72)
duckduckgo	11	1969	1341	518	143 (72)
lastfm	13	4840	3528	954	477 (50)
github	19	3603	2520	601	290 (52)
seriesguide	32	8184	5438	986	625 (37)
connectbot	33	2190	1562	316	74 (77)
textsecure	38	5921	3643	698	330 (53)
k-9	55	19032	11968	3104	1988 (36)
wordpress	57	15066	9775	2431	1362 (44)
Total	266	62562	41071	9968	5412 (54)

Thr and Hop: 10 second budget per deref

Is jumping effective?

unproven derefs

	KLOC	Deref	Nit	Thr Hop	(Impr %)
drupaleditor	3	928	679	179	72 (60)
npr	5	829	617	181	51 (72)
duckduckgo	11	1969	1341	518	143 (72)
lastfm	13	4840	3528	954	477 (50)
github	19	3603	2520	601	290 (52)
seriesguide	32	8184	5438	986	625 (37)
connectbot	33	2190	1562	316	74 (77)
textsecure	38	5921	3643	698	330 (53)
k-9	55	19032	11968	3104	1988 (36)
wordpress	57	15066	9775	2431	1362 (44)
Total	266	62562	41071	9968	5412 (54)

Is jumping effective?

unproven derefs

	KLOC	Deref	Nit	Thr Hop (Impr %)
drupaleditor	3	928	679	179 72 (60)
npr	5	829	617	181 51 (72)
duckduckgo	11	1969	1341	518 143 (72)
lastfm	13	4840	3528	954 477 (50)
github	19	3603	2520	601 290 (52)
seriesguide	32	8184	5438	986 625 (37)
connectbot	33	2190	1562	316 74 (77)
textsecure	38	5921	3643	698 330 (53)
k-9	55	19032	11968	3104 1988 (36)
wordpress	57	15066	9775	2431 1362 (44)
Total	266	62562	41071	9968 5412 (54)

Thresher has 74% fewer than Nit

Is jumping effective?

unproven derefs

	KLOC	Deref	Nit	Thr	Hop (Impr %)
drupaleditor	3	928	679	179	72 (60)
npr	5	829	617	181	51 (72)
duckduckgo	11	1969	1341	518	143 (72)
lastfm	13	4840	3528	954	477 (50)
github	19	3603	2520	601	290 (52)
seriesguide	32	8184	5438	986	625 (37)
connectbot	33	2190	1562	316	74 (77)
textsecure	38	5921	3643	698	330 (53)
k-9	55	19032	11968	3104	1988 (36)
wordpress	57	15066	9775	2431	1362 (44)
Total	266	62562	41071	9968	5412 (54)

Hopper has 54% fewer than Thresher

Is jumping effective?

unproven derefs

	KLOC	Deref	Nit	Thr Hop (Impr %)	% Proven	
drupaleditor	3	928	679	179	72 (60)	92
npr	5	829	617	181	51 (72)	94
duckduckgo	11	1969	1341	518	143 (72)	93
lastfm	13	4840	3528	954	477 (50)	90
github	19	3603	2520	601	290 (52)	92
seriesguide	32	8184	5438	986	625 (37)	92
connectbot	33	2190	1562	316	74 (77)	97
textsecure	38	5921	3643	698	330 (53)	94
k-9	55	19032	11968	3104	1988 (36)	90
wordpress	57	15066	9775	2431	1362 (44)	91
Total	266	62562	41071	9968	5412 (54)	92

Is jumping effective?

unproven derefs

	KLOC	Deref	Nit	Thr Hop (Impr %)	% Proven	
drupaleditor	3	928	679	179	72 (60)	92
npr	5	829	617	181	51 (72)	94
duckduckgo	11	1969	1341	518	143 (72)	93
lastfm	13	4840	3528	954	477 (50)	90
github						92
seriesguide						92
connectbot						97
textsecure						94
k-9						90
wordpress	57	15066	9775	2431	1362 (44)	91
Total	266	62562	41071	9968	5412 (54)	92

Compare with state-of-the art NPE checking work that reports 84-91% proven on normal Java programs!

Triaging alarms to find bugs

Triaging alarms to find bugs

Triaged 200 alarms (from Hopper), 189 false

Triaging alarms to find bugs

Triaged 200 alarms (from Hopper), 189 false

Reasons: insufficient Android modeling, imprecise container and string domains

Triaging alarms to find bugs

Triaged 200 alarms (from Hopper), 189 false

Reasons: insufficient Android modeling, imprecise container and string domains

Only 17 false alarms due to timeouts

Triaging alarms to find bugs

Triaged 200 alarms (from Hopper), 189 false

Reasons: insufficient Android modeling, imprecise container and string domains

Only 17 false alarms due to timeouts

Found 11 bugs in 4 apps
(lastfm, seriesguide, connectbot, wordpress)

5 bugs due to **bad ordering assumptions**

10/11 patches accepted

Triaging alarms to find bugs

Triaging alarms to find bugs

Fixing possible NPE in OverviewActivity.launchSearch #449

Merged Fixing possible NPEs when CheckInDialogFragment.newInstance returns null #450

Merged Fix possible NPE in HostListActivity.onCreateContextMenu #60

Only 17 false Merged kruton merged 1 commit into connectbot:master from sblackshear:service_npe_fix on Mar 28

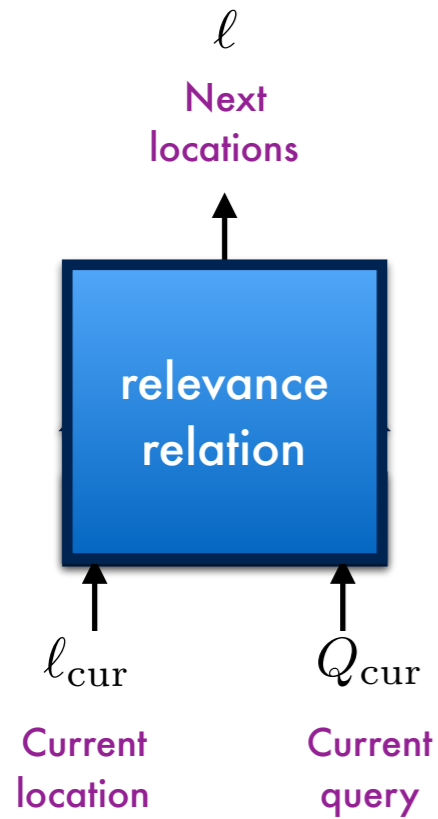
Found 11 bugs in 4 apps
(lastfm, seriesguide, connectbot, wordpress)

5 bugs due to bad ordering assumptions

10/11 patches accepted

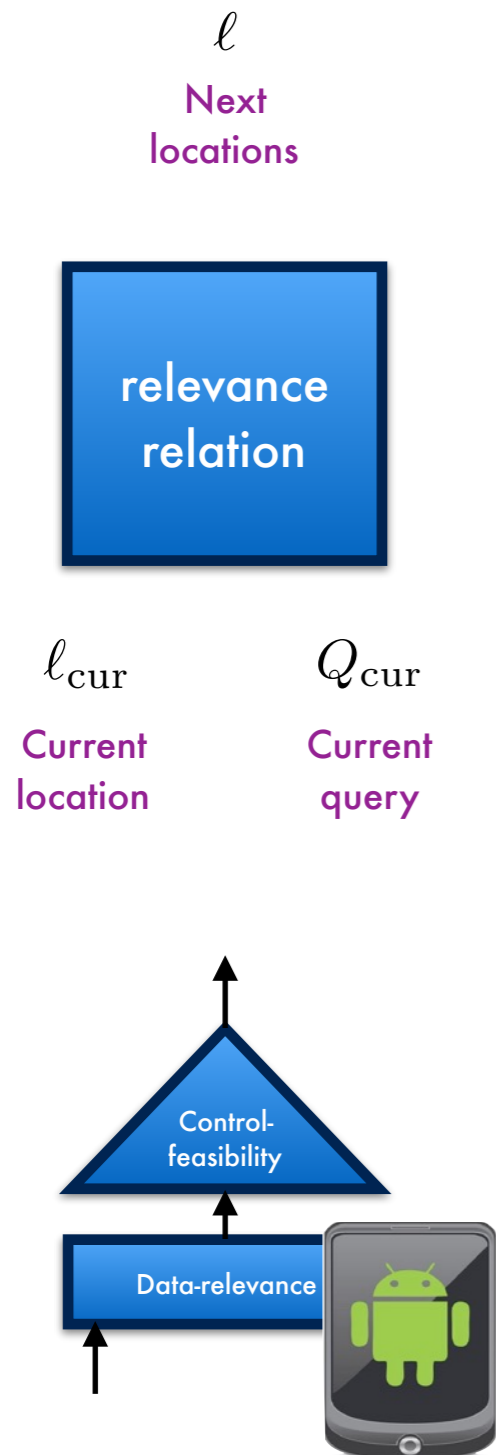
Summary

Summary



**Selective control-flow abstraction
via a sound relevance relation**

Summary



**Selective control-flow abstraction
via a sound relevance relation**

**Effective inter-event ordering-
sensitive reasoning via data-
relevance and control-feasibility**

www.cs.colorado.edu/~bec
pl.cs.colorado.edu



Cerny



Chang



Hammer



Sankaranaryananan



Somenzi

Goal-directed:

abstraction-refinement versus **abstraction-coarsening**

**staged abstraction refinement
(e.g., CEGAR)**

**on-the-fly abstraction
coarsening**

Goal-directed:

abstraction-refinement versus **abstraction-coarsening**

staged abstraction refinement
(e.g., CEGAR)

Staged

Run analysis multiple times,
abstraction changes only
between runs.

on-the-fly abstraction
coarsening

On-the-fly

Run analysis once,
abstraction changes
during analysis.

Goal-directed:

abstraction-refinement versus **abstraction-coarsening**

staged abstraction refinement
(e.g., CEGAR)

Staged

Run analysis multiple times,
abstraction changes only
between runs.

Refinement

Start with imprecise
abstraction,
become more precise.

on-the-fly abstraction
coarsening

On-the-fly

Run analysis once,
abstraction changes
during analysis.

Coarsening

Start with precise
abstraction,
become less precise.